# A Gaussian Leftover Hash Lemma for Modules over Number Fields

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**Abstract.** Given a Gaussian matrix  $\mathbf{X}$ , a Gaussian Leftover Hash Lemma (LHL) states that  $\mathbf{X} \cdot \mathbf{v}$  for a Gaussian  $\mathbf{v}$  is an essentially independent Gaussian sample. It has seen numerous applications in cryptography for hiding sensitive distributions of  $\mathbf{v}$ . We generalise the Gaussian LHL initially stated over  $\mathbb{Z}$  by Agrawal, Gentry, Halevi, and Sahai (2013) to modules over number fields. Our results have a sub-linear dependency on the degree of the number field and require only polynomial norm growth:  $\|\mathbf{v}\|/\|\mathbf{X}\|$ . To this end, we also prove when  $\mathbf{X}$  is surjective (assuming the Generalised Riemann Hypothesis) and give bounds on the smoothing parameter of the kernel of  $\mathbf{X}$ . We also establish when the resulting distribution is independent of the geometry of  $\mathbf{X}$  and establish the hardness of the k-SIS and k-LWE problems over modules (k-M-SIS/k-M-LWE) based on the hardness of SIS and LWE over modules (M-SIS/M-LWE) respectively, which was assumed without proof in prior works.

# 1 Introduction

The classic Leftover Hash Lemma (LHL) is an argument about randomness extraction. It states that if  $\mathbb{H}$  is a distribution over a set of universal hash functions and X is a random variable with some guaranteed entropy then the distributions ( $\mathbb{H}$ ,  $\mathcal{U}$ ) and ( $\mathcal{H} \leftarrow \mathbb{H}$ ,  $\mathcal{H}(X)$ ) are statistically close (where  $\mathcal{U}$  is the uniform distribution). In lattice-based cryptography usage of the LHL appears in a context when the image  $\mathbb{Z}^r$  (or  $\mathcal{O}_{\mathcal{K}}^r$ ) of  $\mathcal{H}(\cdot)$  is taken modulo some q, for example in LWE or SIS based cryptography. In that case, the domain  $\mathcal{O}_{\mathcal{K}}^r/q$  being finite allows to use variants of the classical LHL, with some modifications based on e.g. on how q splits in  $\mathcal{O}_{\mathcal{K}}$ . An overview of the LHL in the context of structured lattice based cryptography was given in [BL25]. A limitation of the classical LHL and its variants mentioned above is that it only applies when the image of  $\mathcal{H}(\cdot)$  is finite. Critically, those statements do not cover the case

when H(X) is close to a discrete Gaussian distribution over a lattice – an infinite domain.

In [AGHS13], Agrawal, Gentry, Halevi, and Sahai prove an LHL-type statement that holds over lattices rather than in finite groups. For  $\mathbf{X} \in \mathbb{Z}^{r \times m}$  with columns sampled independently from  $\mathcal{D}_{\mathbb{Z}^r,\varsigma}$  – a discrete Gaussian over  $\mathbb{Z}^r$  – and  $\mathbf{v}$  sampled from  $\mathcal{D}_{\mathbb{Z}^m,\varsigma'}$ , they upper bound the statistical distance between  $(\mathbf{X},\mathbf{X}\cdot\mathbf{v})$  and  $(\mathbf{X},\mathcal{D}_{\mathbb{Z}^r,\varsigma'},\sqrt{\mathbf{X}\mathbf{X}^r})$  for any  $m \geq r \ln(r\varsigma)$  (constant factors are omitted for this introduction). In follow-up works, Aggrawal and Regev [AR16] improve this result with an alternative proof. In 2020, Kirshanova, Nguyen, Stehlé and Wallet [KNSW20] further improve the result with a better lower bound  $m \geq r \log(\varsigma)$ . Hereafter, we refer to statements of this type as  $Gaussian\ LHLs$ .

As stated the output distribution of  $\mathbf{X} \cdot \mathbf{v}$  depends on the matrix  $\mathbf{X}$  and might be far from a spherical Gaussian. In order to have a spherical output, the solution is usually to sample  $\mathbf{X}$  and then to sample  $\mathbf{v}$  from the Gaussian distribution with parameter  $\varsigma' \cdot \widetilde{\mathbf{X}}$  where  $\widetilde{\mathbf{X}}$  is a pseudo-inverse of  $\mathbf{X}$ . However, in some applications this choice of the distribution of  $\mathbf{v}$  is not available. The 'sphericity' of the distribution  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\sqrt{\mathbf{X}\mathbf{X}^r}}$  is controlled by the singular values of  $\mathbf{X}$ . In [AGHS13, Lemma 8], the authors argue that with high probability for a wide enough Gaussian matrix  $\mathbf{X} \in \mathbb{Z}^{r \times m}$  of parameter  $\varsigma$ , its singular values are within a constant factor of  $\sqrt{m}\varsigma$ . Together, this establishes when  $\mathbf{X} \cdot \mathbf{v}$  follows a distribution close to independent from both of its inputs, enabling various applications discussed below.

A limitation of existing Gaussian LHLs is that they do not take advantage of the algebraic structure of the lattice, while most practical primitives based on lattices currently use module lattices: if  $\mathcal{M}$  is a rank-r module over the ring of integers  $\mathcal{O}_{\mathcal{K}}$  of a degree-d number field  $\mathcal{K}$ , it is a rank- $d \cdot r$  lattice when considered through, e.g. the canonical embedding of  $\mathcal{K}$ . In [KNSW20], the authors give a bound of  $m \geq dr \ln(dr)$  independent vectors required for the Gaussian LHL result to work over  $\mathcal{O}_{\mathcal{K}}^r$ . However, a better bound for the number of independent vectors is to be expected, since any single vector  $\mathbf{v} \in \mathcal{O}_{\mathcal{K}}^r$  gives d-1 other vectors (over  $\mathbb{Z}$ ) 'for free' (think  $X^i \cdot \mathbf{v}$  in the cyclotomic case). We would then expect a lower bound on m that grows linearly in r (the rank of the module) and sublinear in d (the degree of the number field).

**Contributions.** In this work, we use the structure of  $\mathcal{O}_{\mathcal{K}}$  to prove that only  $m \geq r \cdot (\log(dr_{\zeta}))^{1+o(1)}$  vectors are necessary to get a Gaussian LHL in the context of module lattices. In §5, we prove that if  $m \geq r \cdot (\log(dr_{\zeta}))^{1+o(1)}$ ,  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  is a Gaussian matrix and  $\mathbf{v} \in \mathcal{O}_{\mathcal{K}}^{m}$  is a Gaussian vector with parameter  $\sqrt{\Sigma}$  then the distributions  $(\mathbf{X}, \mathbf{X} \cdot \mathbf{v})$  and  $(\mathbf{X}, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r}}, \sqrt{\mathbf{X} \cdot \mathbf{\Sigma} \cdot \mathbf{X}^{\mathsf{T}}})$  have negligible statistical distance. In particular, in §3 we establish when  $\mathbf{X}$  generates  $\mathcal{O}_{\mathcal{K}}^{r}$  and in §4 we upper bound  $\eta_{\varepsilon}(\Lambda^{\perp}(\mathbf{X}))$ . The composition of these two results establishes the claim, but these results may be of independent interest.

Moreover, in §6 we establish a bound when we may expect  $\mathbf{X} \cdot \mathbf{v}$  to be close to spherical for any ring of integers of a number field: if  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  is a Gaussian matrix with  $m \geq (rd)^3$  with parameter  $\varsigma$ , then the Rényi divergence between the

distribution  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\sqrt{\mathbf{X}\mathbf{X}^{\mathsf{T}}}}$  and a spherical Gaussian of parameter  $\approx \varsigma\sqrt{m}$  is O(1). This result permits comparing the distributions  $(\mathbf{X},\mathbf{X}\cdot\mathbf{v})$  and  $(\mathbf{X},\mathbf{v}')$ , where  $\mathbf{v}'$  is sampled from a spherical Gaussian depending only on r,m and  $\varsigma$ .

Finally, as a first explicit application, in  $\S 7$  we give a reduction from M-SIS to k-M-SIS that supports a linear number of hints and discuss how a similar reduction from M-LWE to k-M-LWE can be obtained.

# 1.1 Technical overview

Let  $\mathbf{X} = [\mathbf{x}_1, \dots, \mathbf{x}_m] \in \mathbb{Z}^{r \times m}$  be a matrix whose entries are sampled i.i.d. from a Gaussian distribution over  $\mathbb{Z}$ . As mentioned before, the Gaussian LHL for  $\mathbb{Z}$ -lattices [AGHS13,AR16,KNSW20] states that the distribution  $\mathbf{X} \cdot \mathbf{v}$  when  $\mathbf{v}$  is a Gaussian vector in  $\mathbb{Z}^m$  is close to a Gaussian distribution, even given  $\mathbf{X}$ . A central step in the proof of this result is to give an upper bound for the smoothing parameter of the kernel lattice  $\Lambda^{\perp}(\mathbf{X}) = \{\mathbf{x} \in \mathbb{Z}^m, \mathbf{X} \cdot \mathbf{v} = \mathbf{0}\}$  of the matrix  $\mathbf{X}$ . Two different techniques were introduced to do it. In [AGHS13,KNSW20], the lattice  $\Lambda_q(\mathbf{X}) \subset \mathbb{Z}^m$ , generated by the rows of  $\mathbf{X}$  and scaled standard unit vectors  $q \cdot \mathbf{e}_i$  is considered for some well-chosen prime q. It is then proven that with overwhelming probability, all short vectors of  $\Lambda_q(\mathbf{X})$  are of the form  $\mathbf{X}^T \cdot \mathbf{v}$ , and then the lattice minima of  $\Lambda_q(\mathbf{X})$  are linked to the last minimum of  $\Lambda^\perp(\mathbf{X})$ , by transference theorems in [AGHS13], and by a more involved argument in [KNSW20]. The method used in [KNSW20] gives the best known parameters for Gaussian LHL over  $\mathbb{Z}$  in terms of number of vectors and of  $\varsigma'$ . However, we did not manage to leverage the algebraic structure of  $\mathcal{O}_{\mathcal{K}}$  to improve it.

In [AR16], another approach was used. We describe it here in more detail, since we use it as our starting point. The authors first prove the existence of a short preimage  $\mathbf{U} \in \mathbb{Z}^{m \times r}$  of  $\mathbf{I}_r$  in  $\mathbf{X} \cdot \mathbb{Z}^m$ . This matrix is then used to build a short basis of  $\Lambda^{\perp}(\mathbf{X})$ , using the fact that if  $\mathbf{X} \cdot \mathbf{U} = \mathbf{I}_r$  then all the columns of  $\mathbf{U} \cdot \mathbf{X} - \mathbf{I}_m$  are in the kernel of  $\mathbf{X}$ . The existence of  $\mathbf{U}$  is proven as follows. Let S be the set of formal 0,1 combinations of columns of  $\mathbf{X}$ :  $S = \{\sum_{i=1}^m \{0,1\} \cdot \mathbf{x}_i\}$ . The formal set S has size  $2^m$  by construction, and the vectors inside it take values in the ball of radius  $m\sqrt{r} \cdot \varsigma$  with overwhelming probability. By the pigeonhole principle, this set must contain a collision when  $m \geq r \log(nm\varsigma)$  (again, constant factors are omitted for this introduction). The collision implies that some  $\mathbf{x}_j$  must be a  $0, \pm 1$  combination of columns of  $\mathbf{X}$ , and the authors then argue that this implies with very high probability that the set  $S \cap (S + \mathbf{e}_1)$  is non-empty, leading to a preimage of  $\mathbf{e}_1$  that has coefficients in  $\{0, \pm 1, \pm 2\}$ . The same argument is then repeated with  $\mathbf{e}_j$  for  $2 \leq j \leq n$ , leading to the existence of  $\mathbf{U}$  with coefficients in  $\{0, \pm 1, \pm 2\}$ .

We generalise this proof strategy. Assume that  $\mathbf{X} = [\mathbf{x}_1, \dots, \mathbf{x}_m]$  is a matrix in  $\mathcal{O}_{\mathcal{K}}^{r \times m}$  whose coefficients are sampled i.i.d. from a discrete Gaussian distribution over  $\mathcal{O}_{\mathcal{K}}$ . We aim to prove that for a Gaussian vector  $\mathbf{v} \in \mathcal{O}_{\mathcal{K}}^m$ , the output distribution  $\mathbf{X} \cdot \mathbf{v}$  is close to the Gaussian distribution of variance parameter  $\sqrt{\mathbf{X} \cdot \mathbf{X}^T}$ . Our proof, following the same blueprint as [AR16], requires two intermediary steps.

- 1. Proving that the matrix **X** is surjective in  $\mathcal{O}_{\mathcal{K}}^r$  with high probability.
- 2. Proving an upper bound on the smoothing parameter of the kernel lattice  $\Lambda^{\perp}(\mathbf{X})$  that holds with high probability.

Surjective Gaussian matrices. To our knowledge, a general result about the surjectivity of discrete Gaussian matrices over  $\mathcal{O}_{\mathcal{K}}$  has not been established in prior work. Our strategy relies on methods developed within the theory of random integral matrices. In §3, we generalise the proof of [NP20], which focuses on matrices over  $\mathbb{Z}$  and only gives an asymptotic estimate for the probability. The surjectivity proof over Z consists of looking at the random matrix modulo every prime number p and showing that with high probability X is non-singular modulo p. If X is non-singular modulo every prime p, it implies that the determinant of the lattice spanned by X is not divisible by any primes, and hence that X is surjective. In the number field case, we have to consider prime ideals instead of prime numbers: when a matrix  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  is surjective, in particular its reduction modulo any prime ideal  $\mathfrak{p}$  of  $\mathcal{O}_{\mathcal{K}}$  is surjective. We prove that the converse also holds: **X** is surjective over  $\mathcal{O}_{\mathcal{K}}^r$  if and only if its reduction modulo any ideal  $\mathfrak{p}$ is surjective in  $\mathcal{O}_{\mathcal{K}}/\mathfrak{p}$ . We then prove that a random **X** is surjective with high probability if its coefficients are independent and sampled from a distribution with enough min-entropy. Finally, using a lemma from [JLWG25], we prove that this min-entropy condition holds for Gaussian matrices modulo ideals. Then since  $\mathcal{O}_{\mathcal{K}}$  has an infinite number of prime ideals, we give an upper bound (depending on the Extended Riemann Hypothesis) on the number of prime ideals that need to be considered to prove X's surjectivity. We do so by considering the first  $r \times r$ submatrix of X, proving that it has a non-zero determinant with high probability, and then that only the primes dividing this determinant need to be considered.

The probability of  $\mathbf{X}$  to be surjective highly depends on the number field considered (which was not an issue in the rational integer setting since the field is fixed to  $\mathbb{Q}$ ). We prove that the probability of  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}/\mathfrak{p}^{r \times m}$  to be surjective in  $(\mathcal{O}_{\mathcal{K}}/\mathfrak{p})^r$  is  $\geq 1 - r \cdot \mathcal{N}(\mathfrak{p})^{m-r+1}$  (Lemma 3.1). In particular, if  $\mathcal{O}_{\mathcal{K}}$  has many ideals of small norm<sup>5</sup>, then  $\mathbf{X}$  has to 'be surjective in many different finite fields at the same time' in order to be surjective over  $\mathcal{O}_{\mathcal{K}}^r$ . The formula we obtain depends on the so-called Prime Zeta function of the number field  $P_{\mathcal{K}}(s) = \sum_{\mathfrak{p} \subset \mathcal{O}_{\mathcal{K}}} \mathcal{N}(\mathfrak{p})^{-s}$  which 'counts' the small ideals of  $\mathcal{O}_{\mathcal{K}}$ . We also give a script estimating the growth of this function for cyclotomic fields in Appendix C.2. We present the overall result, both in the general form and in the case of prime-power cyclotomics in §3.5. Informally it states:

Theorem 1.1 (Assuming extended Riemann hypothesis, Informal). If K is a number field of degree d and  $N_K$  is the norm of its smallest ideal then for any  $r \geq 1$ ,  $\varepsilon > 0$ ,  $\varsigma$  large enough and  $m \geq 2r + \log(1/\varepsilon)/\log(N_K)$ , it holds that

$$\Pr_{\mathbf{X} \leftarrow \mathcal{D}_{\mathcal{O}_{K,S}}^{r \times m}} \left( \mathbf{X} \text{ is surjective} \right) \gtrsim 1 - \varepsilon.$$

 $<sup>^{5}</sup>$  For example, one can construct a degree d multiquadratic number field that has d different prime ideals of norm 2.

Short basis of  $\Lambda^{\perp}(\mathbf{X})$  and smoothing parameter. The second step follows and generalises the proof of [AR16]. In that work, the authors first prove existence of a short block-diagonal preimage of  $\mathbf{I}_r$ , and then use it to build a short basis of  $\Lambda^{\perp}(\mathbf{X})$ . To prove that a short preimage exists they define a formal set  $S = \{\sum_{i=1}^{m} \{0,1\} \cdot \mathbf{x}_i\}$  (where the  $\mathbf{x}_i$  are the columns of  $\mathbf{X}$ ). This set has size  $2^m$ , which must contain a collision when  $m \geq r \log(r\varsigma)$  by the pigeonhole principle. This collision is then used to build an element in the intersection  $S \cap (S + \mathbf{e}_i)$  which implies a short preimage of  $\mathbf{e}_i$  for 1 < i < r.

In order to take advantage of the structure of  $\mathcal{O}_{\mathcal{K}}$ , we now consider the set  $S = \{\sum_{i=1}^m a_i \cdot \mathbf{x}_i\}$ , where the  $a_i$  are small elements of a subset  $A \subset \mathcal{O}_{\mathcal{K}}$  (taking  $A = \{0,1\}$  recovers the proof of [AR16]). If the set A is chosen<sup>6</sup> to have size  $2^d$  then the set S has size  $2^{dr}$ , and by running a similar pigeonhole principle argument we obtain a collision for  $m \geq r \log(dr_{\varsigma})$  (instead of  $m \geq dr \log(dr_{\varsigma})$  when taking binary combinations).

However, changing the set A also changes the result. Instead of a preimage for  $\mathbf{e}_i$  we only obtain a preimage for  $a \cdot \mathbf{e}_i$  for  $a \in \{a_1 \cdot a_2, a_i \in A\}$  (Lemma 4.2). Then, if this element  $a \in \mathcal{O}_{\mathcal{K}}$  is not invertible, this does not result in an integral preimage of  $\mathbf{e}_i$  we were looking for. In order to solve this problem, a solution could be to pick  $A \subset \mathcal{O}_{\mathcal{K}}^{\times}$  such that all elements of A and their inverses have small  $\ell_2$  norm. We did not pursue this approach since the best existing subset of  $\mathcal{O}_{\mathcal{K}}^{\times}$  of size  $2^d$ in the literature (to our knowledge) contains elements of norm  $2^{O(\sqrt{d})}$  [CDPR16, Section 6.2. The next option is to run the argument again to get a preimage for  $a' \cdot \mathbf{e}_i$ , with a new a'. Then prove that with high probability a' is coprime with a and run an effective version of Bezout's identity (Lemma 2.3) to get a preimage of  $(au + a'v) \cdot \mathbf{e}_i = \mathbf{e}_i$ . However, the worst-case nature of the pigeonhole principle argument forbids us to make this argument probabilistic. Thus, instead, we have chosen to construct the set  $A' = \{a' \in \mathcal{O}_{\mathcal{K}}, a' \text{ coprime with } a, \|a'\| \leq R \}$  for some R, and to run the argument with this A'. Running the counting argument demands A' to be of exponential size. We prove in Lemma 4.3 that it is the case for large enough  $R = \mathsf{poly} \Big( d, \Delta_{\mathcal{K}}^{1/d} \Big)$ . We think that this coprime-counting lemma might have applications outside the scope of this work.

As an additional contribution, we improve the proof by considering the intersections  $S + \zeta_f^j \cdot \mathbf{e}_i \cap S + \zeta_f^k \cdot \mathbf{e}_i$  when  $\mathcal{K}$  is (or contains) the cyclotomic field of conductor f, allowing to reduce the lower bound of m by a factor of  $\log(f)$ . This yields a matrix  $\mathbf{U} \in \mathcal{O}_{\mathcal{K}}^{m \times r}$  whose norm is bounded by a polynomial in d,  $\Delta_K^{1/d}$  and m such that  $\mathbf{X} \cdot \mathbf{U} = f \cdot \mathbf{I}_r$ . As in [AR16], we then argue that  $f \cdot \mathbf{I}_m - \mathbf{U} \cdot \mathbf{X}$  is a "short" independent set of vectors of  $\Lambda^{\perp}(\mathbf{X})$ , which leads to a 'short' basis and a polynomial upper bound on the smoothing parameter of  $\Lambda^{\perp}(\mathbf{X})$ . This result is presented in Theorem 4.1, and can be summarised as follows:

**Theorem 1.2 (Informal).** If K is a number field of degree d such that  $\mathbb{Q}(\zeta_f) \subseteq K$  then for any  $r \geq 1$ ,  $\varepsilon > 0$ ,  $\varsigma$  large enough and  $m \geq r \cdot (\log(dr\varsigma))^{1+o(1)} + c$ 

<sup>&</sup>lt;sup>6</sup> The reader used to cyclotomic fields can think about  $A = \left\{\sum_{i=1}^{d} \{0,1\} \cdot \zeta^i\right\}$ .

 $\log(1/\varepsilon)/\log(f)$ , there exists an absolute polynomial P such that

$$\Pr_{\mathbf{X} \leftarrow \mathcal{D}_{\mathcal{O}_{K},\varsigma}^{r \times m}} \left( \eta_{\varepsilon} \left( \Lambda^{\perp}(\mathbf{X}) \right) \leq P(d, \Delta_{K}^{1/d}, m) \right) \geq 1 - \varepsilon.$$

We stress that the polynomial in the above theorem has a large degree in d and in  $\delta_{\mathcal{K}}$  (the size of an  $\mathcal{O}_{\mathcal{K}}$  basis as a  $\mathbb{Z}$ -module), leading to a large upper bound on the smoothing parameter of  $\Lambda^{\perp}(\mathbf{X})$ . We believe this massive size is a proof artefact coming from the worst-case nature of our existence argument and support that claim by computing small basis of random kernel lattices of various ranks for cyclotomic fields of various conductors in Appendix A.

To summarise, in §5, we give an explicit statement of the Gaussian Leftover Hash Lemma for structured Gaussian matrices in Theorem 5.1. Let  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  be a Gaussian matrix and  $\mathbf{v} \leftarrow \mathcal{D}_{\mathcal{O}_{\kappa}^{m},\varsigma}$ .

Almost-spherical Gaussians. A limitation of the Gaussian LHL is that the distribution of  $\mathbf{X} \cdot \mathbf{v}$ , even if it is close to a Gaussian, still depends on  $\mathbf{X}$  through its center and its covariance matrix, which is equal to  $\mathbf{X} \mathbf{\Sigma} \mathbf{X}^T$  if  $\mathbf{v}$  is sampled from  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^m,\sqrt{\Sigma}}$ . In §6, we give two options for removing the dependency in the covariance (the centre can usually be set to  $\mathbf{0}$ ). One option is to adapt the covariance of  $\mathbf{v}$  and follows from previous work.

When the first option is not available, e.g.  $\mathbf{v}$  is received from another party or algorithm, we can still extract an almost spherical Gaussian. To overcome this difficulty, we note that if X is of sufficient width, then it is already almost spherical with very high probability. This fact was already mentioned in [AGHS13, Lemma 8 for the  $\mathbb{Z}$  setting, which we refine and generalise to the ring of integers setting. We first prove that the geometry of a discrete Gaussian matrix  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$ is close to the one of a continuous Gaussian matrix  $\widetilde{\mathbf{X}} \leftarrow \mathcal{D}_{\varsigma}^{r \times m}$ , which can be studied using tools from the field of random matrices. We give bounds on the singular values of X following the proof of [Sil85], in which the author was only interested in their asymptotic behaviour. In Lemma 6.1, we prove that its singular values can be described as sums and products of independent  $\chi^2$  random variables. We then use concentration inequalities of the  $\chi^2$  distribution to show that these singular values are all close to  $\sqrt{m} \cdot \varsigma$  with overwhelming probability. Finally, we show that this closeness implies that the Rényi divergence between the distribution of covariance  $\mathbf{X} \mathbf{\Sigma} \mathbf{X}^{\mathsf{T}}$  and  $\sqrt{m} \cdot \varsigma$  is O(1) with overwhelming probability. This result is presented in Theorem 6.1 and can be summarised as follows:

**Theorem 1.3 (Informal).** Let K be a number field of degree d,  $\mathcal{O}_K$  its ring of integers. Let  $\mathbf{X} \in \mathcal{O}_K^{r \times m}$  sampled from  $\mathcal{D}_{\mathcal{O}_K^r,\varsigma}^m$  with  $m \geq (rd)^3$ . Then

$$\mathrm{RD}(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\varsigma};\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\sqrt{\mathbf{X}\cdot\mathbf{X}^{\mathtt{T}}}}) \leq 81.$$

# 1.2 Applications

A direct application of our Gaussian Leftover Hash Lemma over modules are reductions from M-SIS and M-LWE to k-M-SIS and k-M-LWE respectively that support

 $k=\mathcal{O}(m)$  following the blueprint from [LPSS14].<sup>7</sup> In this work, we only give the reduction for M-SIS to k-M-SIS since the M-LWE to k-M-LWE reduction proceeds analogously. This allows us to lift constructions based on k-SIS and k-LWE to the module setting, such as linearly homomorphic signature schemes [BF11], traitor tracing [LPSS14] and partial GPV-style trapdoors [ALLW25]. Moreover, k-SIS may be considered as the oldest and most well-founded (due to the reduction from SIS) member of a family of "SIS with hints" assumptions that have been proposed recently. Indeed, works such as [ACL+22] prove the existence of hard instances of their newly introduced assumption by proving these instances as hard as k-M-SIS. However, in [ACL+22,ALLW25] the hardness of k-M-SIS is only assumed and not established.

Our generalisation of the Gaussian Leftover Hash Lemma to modules also enables the translation of various applications where [AGHS13,AR16] was used to rerandomise ciphertexts to achieve a form of confidentiality, such as the rerandomisation of GSW ciphertexts [BdPMW16] or the speculative suggestion in [ACL $^+$ 22] to use similar techniques to make their construction zero-knowledge. Similarly, the special  $1 \times 2$  case of a Gaussian LHL over rings had been established and used in [LSS14] to prove that a given distribution does not leak secret values. More generally, our result implies a generalisation of Gaussian convolution theorems such as [GMPW20, Theorem 3] to the module setting. Such tools have been used to give tighter reductions for SIS and LWE [MP13].

Moreover, the  $1 \times 2$  case was generalised to more rings and used in [PS21] to establish that NTRU instances can be rerandomised. This in turn establishes a reduction to the decision NTRU problem. Our work is a building block to generalise such results to larger ranks r > 1. Similarly, in [SS13] the primitivity of the matrix of the special case  $1 \times 2$  was considered to prove that NTRU signing keys can be efficiently sampled. Our work is a building block to generalising such reductions to a module setting.

Finally, in LWE-based encryption schemes, showing correctness often involves upper bounding the norm of the product  $\mathbf{X} \cdot \mathbf{e}$  between a short matrix  $\mathbf{X}$  from the decryption key and a short vector  $\mathbf{e}$  from the encryption randomness. A naive upper bound is  $\|\mathbf{X}\| \cdot \|\mathbf{e}\|$ . This can be further upper bounded by  $\sqrt{r} \cdot d \cdot m \cdot s^2$  if both  $\mathbf{X}$  follow a discrete Gaussian distribution  $\mathcal{D}^{r \times m}_{\mathcal{O}_{\mathcal{K}},s}$  and  $\mathbf{e}$  follows  $\mathcal{D}^m_{\mathcal{O}_{\mathcal{K}},s}$ . A better approach is to consider the product for a fixed  $\mathbf{X}$  with  $\mathbf{X} \cdot \mathcal{O}^m_{\mathcal{K}} = \mathcal{O}^r_{\mathcal{K}}$  and  $s \geq \eta_{\epsilon}(\Lambda^{\perp}(\mathbf{X}))$ . In this case,  $\mathbf{X} \cdot \mathbf{e}$  is statistically close to a sample from  $\mathcal{D}^r_{\mathcal{O}_{\mathcal{K}},s\cdot\sqrt{\mathbf{X}^{\mathsf{T}\cdot\mathbf{X}}}}$ . Further upper bounding  $\|\mathbf{X}\|$  and  $\|\mathcal{D}_{\mathcal{O}^r_{\mathcal{K}},\sqrt{\mathbf{X}^{\mathsf{T}\cdot\mathbf{X}}\cdot s}}\|$  then yields a tighter bound of  $\sqrt{r \cdot d \cdot m} \cdot s^2$  at the cost of additional correctness error from the Gaussian LHL. In the literature, the second approach seems to have been applied in the setting where  $\mathcal{O}_{\mathcal{K}} = \mathbb{Z}$  or when r = 1 but not in more general settings, probably due to the difficulty of upper bounding  $\eta_{\epsilon}(\Lambda^{\perp}(\mathbf{X}))$ , which we take first steps to address in §4.

<sup>&</sup>lt;sup>7</sup> The reduction from SIS to k-SIS in [BF11] which supports  $k = \mathcal{O}(1)$  translates without our Gaussian LHL generalisation.

# 1.3 Open question and directions

Even if our Gaussian LHL gives a polynomially bounded norm growth between the norm of the matrix  $\mathbf{X}$  and the norm of the vector  $\mathbf{v}$ , the degree of the norm growth is large (see Theorem 5.1). This blowup comes from the fact that the upper bound on the smoothing parameter of  $\Lambda^{\perp}(\mathbf{X})$  is very loose. We think that this is a proof artefact, and give experiments backing this claim in Appendix A for cyclotomic number fields.

The looseness of this bound comes from the fact that the proof of Lemma 4.4 is fundamentally "worst-case". A first way to improve it would be to use more probabilistic arguments instead of pigeonhole-like ones (see §4.3). Another blowup factor in the proof of Lemma 4.4 is that when constructing a set of size  $2^d$  of elements coprime to some  $a \in \mathcal{O}_K$ , Lemma 4.3 give element of size  $\simeq ||a||^2$ . We think that restricting Lemma 4.3 to specific number fields (e.g. prime-power cyclotomics, where the prime ideal distribution is better known) will make our result more usable in practice. One last way of improving the bound of Lemma 4.4 would be to find a trade-off relation between the size of the smoothing parameter of  $\Lambda^{\perp}(\mathbf{X})$  and the width of  $\mathbf{X}$ . Currently, our proof shows a threshold effect: if the width of  $\mathbf{X}$  is less than a certain  $m_0$ , then the smoothing parameter is large, if it is above  $m_0$ , we have an upper bound on it which does not depend much on m. Finding more subtle relations between the two would allow to improve the bound, at the cost of increasing the width of  $\mathbf{X}$ .

# 2 Preliminaries

For  $i, j \in \mathbb{Z}$  denote  $[i, j] = \{i, \dots, j-1\}$ . We denote by  $\mathbb{Z}$  the set of integers,  $\mathbb{Q}$  the rationals,  $\mathbb{R}$  the real numbers and  $\mathbb{C}$  the complex numbers. For any positive function  $f: \mathbb{R}^m \to \mathbb{R}$  and discrete set  $S \subset \mathbb{R}^m$ , we write  $f(S) = \sum_{x \in S} f(x)$ .

We use the convention that  $\log = \log_2$  and  $\ln = \log_e$ . Vectors are denoted by bold-lower-case letters, such as  $\mathbf{v}$ . Matrices are denoted by bold-upper-case letters, such as  $\mathbf{M}$ . When  $\mathbf{A}, \mathbf{B}$  are two matrices with compatible sizes, we let  $[\mathbf{A} \mid \mathbf{B}]$  denote the horizontally stacked block matrix and  $[\mathbf{A} \mid \mathbf{B}] := [\mathbf{A}^T \mid \mathbf{B}^T]^T$  the vertically stacked one. When  $\mathbf{A}$  is a matrix in  $\mathbb{C}^{r \times m}$ , we denote  $\|\mathbf{A}\|$  the maximum of the  $\ell_2$  norm of its columns. In particular, for any  $\mathbf{x} \in \mathbb{C}^m$ , it holds that  $\|\mathbf{A} \cdot \mathbf{x}\| \le \sqrt{m} \cdot \|\mathbf{A}\| \cdot \|\mathbf{x}\|$ .

# 2.1 Lattices, number theory and modules

A lattice  $\Lambda \subset \mathbb{R}^m$  is a discrete additive subset of  $\mathbb{R}^m$ . It is generated by a (non-unique) basis  $\mathbf{b}_1, \dots, \mathbf{b}_r$  (often denoted in matrix form  $\mathbf{B} = [\mathbf{b}_1, \dots, \mathbf{b}_r] \in \mathbb{R}^{m \times r}$ ) of  $\mathbb{R}$ -linearly independent vectors. When m = r, the lattice is said to be full-rank. The volume of a lattice  $\Lambda = \mathcal{L}(\mathbf{B})$  is  $\operatorname{Vol}(\Lambda) := \det(\mathbf{B}^\mathsf{T} \cdot \mathbf{B})$ , and is independent of the choice of basis  $\mathbf{B}$ . Let  $\Lambda$  be a lattice of rank r. We denote by  $\lambda_i(\Lambda)$  its ith minima for  $1 \le i \le r$ :  $\lambda_i(\Lambda) = \max(\min\{(\|\mathbf{v}_1\|, \dots, \|\mathbf{v}_i\|), \mathbf{v}_j \in \Lambda, \operatorname{span}((\mathbf{v}_j)_j) = i\})$ .

For a more in-deepth introduction to number fields and number theory in general, we refer the reader to [Coh93,Neu13]. A number field is a finitedimentional extension of  $\mathbb{Q}$ , its degree being the degree of the extension. We let  $\mathcal{O}_{\mathcal{K}}$  denote its ring of integers and  $\Delta_{\mathcal{K}}$  its discriminant. A fractional ideal (we often omit saying "fractional" in this work) is a discrete subring of K that is stable under multiplication by any element of  $\mathcal{O}_{\mathcal{K}}$ . We say that an ideal is integral if it is included in  $\mathcal{O}_{\mathcal{K}}$ . The norm of an integral ideal  $\mathfrak{a}$  is  $\mathcal{N}(\mathfrak{a}) = |\mathcal{O}_{\mathcal{K}}/\mathfrak{a}|$ . The norm of a fractional ideal I is  $\mathcal{N}(I) = \mathcal{N}(N \cdot I)/N$ , where N is an integer such that  $N \cdot I$  is integral (such an integer always exists, and it can be shown that the norm does not depend on the choice of N). In this work, we denote integral ideals by gothic letters  $(\mathfrak{a},\mathfrak{p}...)$  and general ideal by uppercase letters  $(I,J,\ldots)$ . The set of ideals has a group structure, where the product of two ideals I, J is the ideal  $I \cdot J = \{a \cdot b, a \in I, b \in J\}$ , and the inverse of an ideal Iis the ideal  $I^{-1} = \{x \in K, x \cdot I \subseteq \mathcal{O}_{\mathcal{K}}\}$ . Ideals have a unique prime factorisation property. An ideal  $\mathfrak{p} \subset \mathcal{O}_{\mathcal{K}}$  is said to be prime if  $\mathcal{O}_{\mathcal{K}}/\mathfrak{p}$  is a field, and for every ideal  $I \subset \mathcal{K}$ , there exist  $\mathfrak{p}_1, \ldots, \mathfrak{p}_k$  prime ideals and  $e_1, \ldots, e_k \in \mathbb{Z}$  such that  $I = \prod_{i=1}^k \mathfrak{p}_i^{e_i}$ . Up to reordering  $(\mathfrak{p}_i)_i$  and  $(e_i)_i$  are unique. Some of our results rely on the Generalised Riemann Hypothesis (GRH) [BS96, Conjecture 8.7.3], which implies in particular that for any number field K and any  $\varepsilon > 0$ , the prime counting function  $\pi_{\mathcal{K}}(x) = |\{\mathfrak{p} \text{ prime ideal}, \mathcal{N}(\mathfrak{p}) \leq x\}| \text{ satisfies,}$ 

$$\pi_{\mathcal{K}}(x) = \int_{2}^{x} 1/\ln(t)dt + O(x^{1/2+\varepsilon}).$$

A number field  $\mathcal K$  of degree d can be embedded in  $\mathbb C$  via d different maps, named its canonical embedding. We let  $d_{\mathbb R}$  be the number of those embeddings that only take real values, and  $d_{\mathbb C}$  be the number of pairs of complex-valued embeddings. It holds that  $d=d_{\mathbb R}+2d_{\mathbb C}$ . Those maps define an injective map from  $\mathcal K$  to its completion  $\mathcal K_{\mathbb R}:=\mathcal K\otimes\mathbb R$  called the Minkowski embedding  $\Phi:\mathcal K\to\mathcal K_{\mathbb R}$ . The set  $\mathcal K_{\mathbb R}$  is isomorphic as a Hermitian space to the set  $\left\{\mathbf x\in\mathbb C^d, x_{d_{\mathbb R}+i}=\overline{x_{d_{\mathbb R}+d_{\mathbb C}+i}}\right\}$ . This set is in turn isomorphic to  $\mathbb R^d$  as an inner-product space. The Minkowski embedding allows us to endow  $\mathcal K$  with a geometric structure: for any  $x\in\mathcal K$ , we denote by  $\|x\|$  (respectively  $\|x\|_{\infty}$ ) the  $\ell_2$  (respectively  $\ell_{\infty}$ ) norm of its Minkowski embedding (in  $\mathbb C^d$ ). Note that in particular, for any  $a,b\in\mathcal K$ , it holds that  $\|a\cdot b\|\leq \|a\|_{\infty}\cdot\|b\|$ . In this geometric setting, any ideal I of  $\mathcal K$  (including its ring of integer) is a lattice in  $\mathcal K_{\mathbb R}$  of volume  $\mathcal N(I)\cdot\sqrt{\Delta_{\mathcal K}}$ .

Throughout this work, we assume that we know a basis  $\mathbf{B}^{\mathcal{O}_{\mathcal{K}}} \in \mathbb{R}^{d \times d}$  of  $\Phi(\mathcal{O}_{\mathcal{K}})$  as a  $\mathbb{Z}$ -module. We denote by  $\delta_K$  the maximal  $\ell_{\infty}$  norm of its columns (in particular  $\delta_K = 1$  for cyclotomics fields with the power basis). The results of this work are existential instead of constructive, so we use the 'best' basis of  $\mathcal{O}_{\mathcal{K}}$ . According to [MG02, Corollary 7.2] there always exists a basis of  $\mathcal{O}_{\mathcal{K}}$  such that  $1 \leq \delta_K \leq \sqrt{d} \cdot \lambda_d(\mathcal{O}_{\mathcal{K}})/2$ .

An  $\mathcal{O}_{\mathcal{K}}$ -module  $M \subset \mathcal{K}^m$  is a discrete subset of  $\mathcal{K}^m$  that is additive and stable under multiplication by any element of  $\mathcal{O}_{\mathcal{K}}$  (fractional ideals of  $\mathcal{O}_{\mathcal{K}}$  are rank-1 modules). Any  $\mathcal{O}_{\mathcal{K}}$ -module M can be represented as a (non-unique) pseudo basis  $((I_i, \mathbf{b}_i))_{i=1,\dots,r}$ , where  $I_i$  are ideals of  $\mathcal{K}$  and  $\mathbf{b}_i \in \mathcal{K}^m$  are  $\mathcal{K}$ -linearly independent.

When r = m, we say that M is full-rank. By the Minkowski embedding, every  $\mathcal{O}_{\mathcal{K}}$ -module of rank r in  $\mathcal{K}^m$  can be seen as a lattice of rank  $d \cdot r$  in  $\mathbb{R}^{d \cdot m}$ .

Let  $M \subset \mathcal{K}^r$  be a full-rank module given by a pseudo-basis  $((\mathbf{b}_i, I_i))_{1 \leq i \leq r}$ . Its determinantal ideal is the fractional ideal  $\det(M) := \det_K(\mathbf{B}) \cdot \prod_i I_i$ , and it is independent of the choice of the pseudo-basis. For any submodule  $M' \subseteq M$ , it holds that  $\det(M') | \det(M)$  with equality if and only if M = M'.

Lemma 2.1 (Adapted from [MG02, Lemma 7.1]). Let  $M \subset \mathcal{K}^r$  be a full-rank module, and  $\mathbf{v_1}, \ldots, \mathbf{v_r} \in M$  be a  $\mathcal{K}$ -free set of vectors. Then there exists a matrix  $\mathbf{B} \in \mathbb{R}^{dr \times dr}$  that generates  $\Phi(M)$  as a  $\mathbb{Z}$ -lattice satisfying

$$\|\mathbf{B}\| \leq \sqrt{rd} \cdot \delta_{\mathcal{K}} \cdot \max \|\mathbf{v}_i\|.$$

**Lemma 2.2.** Let  $M \subset \mathcal{K}^r$  be a full-rank module of,  $\mathbf{b}_1, \dots, \mathbf{b}_r \in M$ , let  $\mathbf{B} = [\mathbf{b}_1, \dots, \mathbf{b}_r]$  and  $\mathfrak{p}$  a prime ideal of  $\mathcal{O}_{\mathcal{K}}$  with residue field  $F = \mathcal{O}_{\mathcal{K}}/\mathfrak{p}$ . Then  $\det_K(\mathbf{B}) \in \det(M)$ , and furthermore it holds that  $\det_F(\mathbf{B} \bmod \mathfrak{p}) = \det_K(\mathbf{B}) \bmod (\det(M) \cdot \mathfrak{p})$ .

Proof. The fact that  $\det_K(\mathbf{B}) \in \det(M)$  comes from the fact that  $\mathbf{B} \cdot \mathcal{O}_K^r \subseteq M$  and that  $\det_K(\mathbf{B}) \cdot \mathcal{O}_K = \det(\mathbf{B} \cdot \mathcal{O}_K^r)$ . This implies that  $\det(\mathbf{B}) \mod \det(M) \cdot \mathfrak{p}$  is well-defined. Now let us define the function  $\phi(\mathbf{b}_1, \dots, \mathbf{b}_r) := \det_K(\mathbf{B}) \mod \det(M) \cdot \mathfrak{p}$ . This function is a multilinear map satisfying that for any  $\mathbf{m} \in M$  and  $p \in \mathfrak{p}$ ,  $\phi(\mathbf{b}_1 + p \cdot \mathbf{m}, b_2, \dots, b_r) = \phi(\mathbf{B}) + p \cdot \det_K(\mathbf{m}, \mathbf{b}_2, \dots, \mathbf{b}_r) = \phi(\mathbf{B}) \mod \mathfrak{p} \det(M)$ , so it can be lifted as a multilinear map  $\phi' : (M/\mathfrak{p}M)^r \to F$ . The multilinear map  $\phi'$  is alternating and clearly satisfies  $\phi'(\mathbf{I}_r) = 1$ , it is then equal to the determinant over  $(M/\mathfrak{p}M)^r$ .

**Lemma 2.3.** Let  $a, b \in \mathcal{O}_{\mathcal{K}}$  be two coprime elements. Then there exists  $u, v \in \mathcal{O}_{\mathcal{K}}$  with  $\|u\|, \|v\| \le \max(\|a\|_{\infty}, \|b\|_{\infty}) \cdot \sqrt{d} \cdot \lambda_d(\mathcal{O}_{\mathcal{K}})$  such that  $a \cdot u + b \cdot v = 1$ .

Proof. Without loss of generality, assume that  $\|a\|_{\infty} \leq \|b\|_{\infty}$ . Let  $b_1, \ldots, b_d$  be a  $\mathbb{Z}$ -basis of  $\mathcal{O}_{\mathcal{K}}$  satisfying  $\|b_i\| \leq \max(1, \sqrt{i}/2) \cdot \lambda_i(\mathcal{O}_{\mathcal{K}})$ . Let M be the rank-1 module in  $\mathcal{K}^2$  generated by  $(b_i \cdot (-b, a)^T)_{1 \leq i \leq d}$ . The norm of the elements of its basis are bounded by  $\lambda_d(\mathcal{O}_{\mathcal{K}}) \cdot \|a\|_{\infty}$ , hence the covering radius of M is at most  $\sqrt{d} \cdot \lambda_d(\mathcal{O}_{\mathcal{K}}) \cdot \|a\|_{\infty}$ . Let  $u_0, v_0 \in \mathcal{O}_{\mathcal{K}}$  such that  $a \cdot u_0 + b \cdot v_0 = 1$  and let (u, v) be the closest element of M to  $(u_0, v_0)$ , then it holds that  $a \cdot (u_0 - u) + b \cdot (v_0 - v) = 1$  and that  $\|(u_0 - u, v_0 - v)\| \leq \sqrt{d} \cdot \lambda_d(\mathcal{O}_{\mathcal{K}}) \cdot \|a\|_{\infty}$ , hence the result.  $\square$ 

When  $\mathbf{X} \in \mathcal{K}^{r \times m}$  is a matrix, we define its kernel module  $\Lambda^{\perp}(\mathbf{X}) := \{\mathbf{x} \in \mathcal{O}_{\mathcal{K}}^m, \mathbf{X} \cdot \mathbf{x} = \mathbf{0}\}$ . When  $\mathbf{X}$  has  $\mathcal{K}$ -rank l, it is a  $\mathcal{O}_{\mathcal{K}}$ -module of rank m - l.

### 2.2 Matrices and Singular Value Decomposition

Let  $\mathbf{M} \in \mathbb{R}^{n \times n}$  be a symmetric matrix. We say that  $\mathbf{M}$  is positive definite if for any  $\mathbf{x} \in \mathbb{R}^n \setminus \{\mathbf{0}\}$ , it holds that  $\mathbf{x}^T \cdot \mathbf{M} \cdot \mathbf{x} > 0$ . For any positive definite matrix  $\mathbf{M}$ , there exists a symmetric matrix  $\mathbf{S}$  such that  $\mathbf{M} = \mathbf{S}^T \mathbf{S}$ . We call this matrix a square root of  $\mathbf{M}$ , it is not unique.

Let  $\Sigma_1, \Sigma_2$  be two positive definite matrices. We say that  $\Sigma_1 \leq \Sigma_2$  if  $\Sigma_2 - \Sigma_1$  is positive definite. Let  $\varsigma > 0$ , we abuse the notation and write  $\Sigma \geq \varsigma^2$  for  $\Sigma \geq \varsigma^2 \mathbf{I}_{dr}$ . For any two invertible matrices  $\mathbf{S}, \mathbf{T}$ , we write  $\mathbf{S} \leq \mathbf{T}$  if  $\mathbf{S} \cdot \mathbf{S}^{\mathsf{T}} \leq \mathbf{T} \cdot \mathbf{T}^{\mathsf{T}}$ . In particular, for any invertible  $\mathbf{A} \geq \mathbf{B} \in \mathbb{R}^{n \times m}$ , it holds that  $\|\mathbf{A}\mathbf{x}\| \geq \|\mathbf{B}\mathbf{x}\|$  for any  $\mathbf{x} \in \mathbb{R}^m$ .

For any matrix  $\mathbf{M} \in \mathbb{R}^{n \times m}$  of rank r, there exist real values  $0 < s_r(\mathbf{M}) \le \ldots \le s_1(\mathbf{M})$ , and two orthogonal matrices  $\mathbf{U}_M \in \mathbb{R}^{n \times n}, \mathbf{V}_M \in \mathbb{R}^{m \times m}$  such that

$$\mathbf{M} = \mathbf{U}_M \cdot \begin{pmatrix} \operatorname{diag}(\{s_i(\mathbf{M})\}_{i=1}^n) \ \mathbf{0} \\ \mathbf{0} & \mathbf{0} \end{pmatrix} \cdot \mathbf{V}_M.$$

The  $s_i(\mathbf{M})$  are called the singular values of  $\mathbf{M} \in \mathbb{R}^{n \times m}$ . We denote  $s_{\max}(\mathbf{M}) := s_1(\mathbf{M})$  and  $s_{\min}(\mathbf{M}) := s_r(\mathbf{M})$ . For any  $\mathbf{x} \in \mathbb{R}^m$ , it holds that  $\|\mathbf{M} \cdot \mathbf{x}\| \le s_{\max}(\mathbf{M}) \cdot \|\mathbf{x}\|$ . If  $\mathbf{M}$  is square and full rank, it also holds that  $s_{\min}(\mathbf{M}) \cdot \|\mathbf{x}\| \le \|\mathbf{M} \cdot \mathbf{x}\|$ .

**Lemma 2.4.** Let  $\mathbf{X}, \mathbf{Y} \in \mathbb{R}^{n \times m}$  with  $0 < n \le m$ . If  $s_{\max}(\mathbf{Y}) \le \delta \cdot s_{\min}(\mathbf{X})$  for some  $\delta < 1$  then

$$s_{\max}(\mathbf{X} \pm \mathbf{Y}) \in [1 - \delta, 1 + \delta] \cdot s_{\max}(\mathbf{X}) ,$$
  
 $s_{\min}(\mathbf{X} \pm \mathbf{Y}) \in [1 - \delta, 1 + \delta] \cdot s_{\min}(\mathbf{X}) .$ 

Lemma 2.5 ([LPSS14, Adapted Lemma 16 (eprint)]). Let  $1 \le n \le m$  be integers and  $\varsigma > 0$ . Let  $\mathbf{M} \in \mathbb{R}^{n \times m}$  be a full-rank matrix. There is a polynomial-time algorithm in the size of its input constructing a matrix  $\mathbf{S} \in \mathbb{R}^{m \times m}$  with  $s_1(\mathbf{S}) \le \varsigma/s_n(\mathbf{M})$  such that

$$\varsigma^2 \cdot \mathbf{I}_n = \mathbf{M} \cdot \mathbf{S} \cdot \mathbf{S}^{\mathsf{T}} \cdot \mathbf{M}^{\mathsf{T}}.$$

*Proof.* Assume that  $\mathbf{M} = \mathbf{V}_M \cdot \operatorname{diag}(\{s_i(\mathbf{M})\}_{i=1}^n) \cdot \mathbf{U}_M$  is a singular value decomposition of  $\mathbf{M}$  with  $\mathbf{V}_M \in \mathbb{R}^{n \times n}$ ,  $\mathbf{U}_M \in \mathbb{R}^{n \times m}$  orthogonal. Set  $\mathbf{S} = [\mathbf{U}_M^\mathsf{T} | \mathbf{U}'] \cdot \operatorname{diag}(\{s_i\}_{i=1}^m)$  where  $\mathbf{U}'$  is an arbitrary extension of  $\mathbf{U}_M^\mathsf{T}$  to an orthonormal basis of  $\mathbb{R}^m$  and  $s_i = \varsigma/s_i(\mathbf{M})$  for  $i \le n$  and  $s_i = 1$  for  $n < i \le m$ . Computing the product we obtain the statement.

# 2.3 Probability

Let  $D_1, D_2$  be two (discrete or absolutely continuous w.r.t. the Lebesgue measure over  $\mathbb{R}^n$  for some  $n \geq 1$ ) probability distribution defined over the same  $\sigma$ -algebras. If  $\Omega = \operatorname{Supp}(D_1) \cup \operatorname{Supp}(D_2)$ , then the statistical distance between  $D_1$  and  $D_2$  is  $\operatorname{SD}(D_1, D_2) := \int_{t \in \Omega} |D_1(t) - D_2(t)| dt/2$ . If  $\operatorname{Supp}(D_1) \subseteq \operatorname{Supp}(D_2)$ , then order-2 Renyi divergence  $\operatorname{RD}_2(D_1||D_2) := \int_{t \in \operatorname{Supp}(D_1)} D_1(t)^2/D_2(t) dt$ . In particular for any event  $E \subseteq \operatorname{Supp}(D_2)$  it holds that

$$D_2(E) \ge D_1(E)^2 / \text{RD}(D_1 || D_2)$$
 (1)

Let  $\mathbf{S} \in \mathbb{R}^{dr \times dr}$  be full rank and  $\mathbf{c} \in \mathcal{K}^r$  we define for any  $x \in \mathcal{K}^r$ ,

$$\rho_{\mathbf{S},\mathbf{c}}(\mathbf{x}) := \exp(-\pi \cdot \Phi(\mathbf{x} - \mathbf{c})^{\mathsf{T}} \cdot (\mathbf{S}^{\mathsf{T}} \cdot \mathbf{S})^{-1} \cdot \Phi(\mathbf{x} - \mathbf{c}))$$

If  $\mathbf{S} = \varsigma \cdot \mathbf{I}_{dr}$  we simplify to  $\rho_{\varsigma,\mathbf{c}}$ , and if  $\mathbf{c} = \mathbf{0}$ , we omit it. For any matrix  $\mathbf{S} \in \mathbb{R}^{dr \times dr}$  and center  $\mathbf{c}$ , we denote  $\mathcal{D}_{\mathbf{S},\mathbf{c}}$  the continuous Gaussian distribution of covariance  $\mathbf{S}^T \cdot \mathbf{S}$  and center  $\mathbf{c}$ . For any lattice  $\Lambda \subset \mathbb{R}^r$ , the discrete Gaussian distribution over  $\Lambda$  of parameter  $\mathbf{S}$  and centre  $\mathbf{c}$ , denoted  $\mathcal{D}_{\Lambda,S,\mathbf{c}}$  is the distribution

$$\mathcal{D}_{\Lambda,\mathbf{S},\mathbf{c}}(\mathbf{x}) = \frac{\rho_{\mathbf{S},\mathbf{c}}(\mathbf{x})}{\rho_{\mathbf{S},\mathbf{c}}(\Lambda)}.$$

For a positive definite matrix  $\Sigma$ , if a statement does not depend on the particular choice of square root, we write  $\sqrt{\Sigma}$  to denote an arbitrary choice of the square root. In particular, we write  $\mathcal{D}_{\Lambda,\sqrt{\Sigma},\mathbf{c}}$  for the Gaussian distribution.

The smoothing parameter of a lattice  $\Lambda$  is a quantity defined for any  $\varepsilon > 0$  as  $\eta_{\varepsilon}(\Lambda) = \inf_{\varsigma > 0} (\rho_{1/\varsigma}(\Lambda^*) \leq 1 + \varepsilon)$ . For any  $\varepsilon > 0$  and  $k \geq 1$ , we define  $\eta_{\varepsilon}^{(k)} = \sqrt{\ln(2k \cdot (1 + 1/\varepsilon))/\pi}$  and it holds that  $\eta_{\varepsilon}(\Lambda) \leq \lambda_k(\Lambda) \cdot \eta_{\varepsilon}^{(k)}$  for rank-k lattices  $\Lambda$  ([KNSW20, Lemma 7 (eprint)]). For any full rank matrix  $\mathbf{S} \in \mathbb{R}^{n \times n}$  and any full-rank lattice  $\Lambda \subset \mathbb{R}^n$ , we say that  $\eta_{\varepsilon}(\Lambda) \leq \mathbf{S}$  if  $\eta_{\varepsilon}(\mathbf{S}^{-1}\Lambda) \leq 1$ . Note that this only depends on  $\mathbf{S} \cdot \mathbf{S}^{\mathsf{T}}$ , and is coherent with the partial ordering of matrices we previously defined.

Lemma 2.6 (Adapted from [MP12, Lemma 2.6], [Pei08, Lemma 5.3]). Let  $\Lambda \subset \mathbb{R}^n$  be an n-dimensional lattice and s, t > 0, then

$$\Pr(\|\mathcal{D}_{\Lambda,\varsigma}\| > \varsigma\sqrt{n}) \le 2^{-n},$$
  
$$\Pr(\|\mathcal{D}_{\Lambda,\varsigma}\|_{\infty} > \varsigma \cdot t) \le 2n \cdot e^{-\pi t^2}.$$

**Lemma 2.7.** Let  $\Lambda \subset \mathbb{R}^n$  be an n-dimensional lattice and t > 0 and  $\mathbf{S} \in \mathbb{R}^{n \times n}$  a non-singular matrix, then

$$\Pr(\|\mathcal{D}_{A,\mathbf{S}}\| > s_{\min}(\mathbf{S}) \cdot t \cdot \sqrt{n}) \le \beta_n(t),$$

where  $\beta_n(t) = (\sqrt{2\pi e} \cdot \exp(-\pi \cdot t))^n$ .

*Proof.* Since **S** is square and non-singular, it holds that  $\mathbf{S}^{-T} \cdot \mathcal{D}_{\Lambda,\mathbf{S}}$  is the same distribution as  $\mathcal{D}_{\mathbf{S}^{-T}\Lambda,1}$ , where  $\mathbf{S}^{-T}$  is the inverse transpose of **S**. Now, by [Ban93, Lemma 1.5], it holds that  $\Pr(\|\mathcal{D}_{\mathbf{S}^{-T}\Lambda,1}\| > t \cdot \sqrt{n}) \leq (\sqrt{2\pi e} \exp(-\pi \cdot t))^n$ , that is to say that

$$\Pr(\|\mathbf{S}^{-T} \cdot \mathcal{D}_{\Lambda,\mathbf{S}}\| > t \cdot \sqrt{n}) \le (\sqrt{2\pi e} \exp(-\pi \cdot t))^n.$$

Now, note that for any  $\mathbf{x} \in \mathbb{R}^n$ , it holds that  $\|\mathbf{S}^{-T} \cdot \mathbf{x}\| \le s_1(\mathbf{S}^{-T}) \cdot \|\mathbf{x}\|$ . Since  $s_1(\mathbf{S}^{-T}) = 1/s_{\min}(\mathbf{S})$ , the result follows.

Lemma 2.8 ([GMPW20, Theorem 3.1 (eprint)]). Let  $\varepsilon \in [0,1)$  and  $1 \le n \le m$  be integers. Let  $\Lambda \subset \mathbb{R}^m$  be an m-dimensional lattice,  $\mathbf{c} \in \mathbb{R}^m$  and  $\Sigma \in \mathbb{R}^{m \times m}$  be positive definite. For a matrix  $\mathbf{T} \in \mathbb{R}^{n \times m}$  such that  $\operatorname{span}_{\mathbb{R}}(\Lambda \cap \ker(\mathbf{T})) = \operatorname{span}_{\mathbb{R}}(\ker(\mathbf{T}))$  and  $\eta_{\varepsilon}(\Lambda \cap \ker(\mathbf{T})) \le \sqrt{\Sigma}$ . It holds

$$\mathrm{SD}(\mathbf{T}\cdot\mathcal{D}_{\varLambda,\sqrt{\mathbf{\Sigma}},\mathbf{c}},\mathcal{D}_{\mathbf{T}\cdot\varLambda,\sqrt{\mathbf{\Sigma}'},\mathbf{T}\cdot\mathbf{c}}) \leq \frac{\varepsilon}{1-\varepsilon}$$

where  $\Sigma' = \mathbf{T} \cdot \mathbf{\Sigma} \cdot \mathbf{T}^{\mathsf{T}}$ .

**Lemma 2.9 ([Reg09, Claim 3.9]).** Let  $\Lambda$  be an n-dimensional lattice,  $\varepsilon \in (0, 1/2)$ , and r, s > 0 such that  $\frac{r \cdot s}{\sqrt{r^2 + s^2}} \ge \eta_{\varepsilon}(\Lambda)$ . Then

$$SD(\mathcal{D}_{\Lambda,r} + \mathcal{D}_s^n, \mathcal{D}_{\sqrt{r^2 + s^2}}^n) \le 4 \varepsilon.$$

**Lemma 2.10** ([BL00, Lemma 1 (p. 1325)]). Let  $\chi_m^2$  be a Chi-squared random variable of degree m > 0 then

$$\Pr(\chi_m^2 \ge m + 2\sqrt{mk} + 2k) \le \exp(-k) ,$$
  
$$\Pr(\chi_m^2 \le m - 2\sqrt{mk}) \le \exp(-k) .$$

**Lemma 2.11.** Let  $S \subset \mathcal{O}_{\mathcal{K}}^r$  be a set such that S = -S and  $\mathbf{S}$  an invertible matrix. Then for any  $\mathbf{c} \in \mathcal{O}_{\mathcal{K}}^r$ 

$$\rho_{\mathbf{S}}(\mathbf{c} + S) \ge \rho_{\mathbf{S}}(\mathbf{c}) \cdot \rho_{\mathbf{S}}(S)$$

Proof. For any  $\mathbf{s} \in S$ , it holds that  $\rho_{\mathbf{S}}(\mathbf{c}+\mathbf{s}) + \rho_{\mathbf{S}}(\mathbf{c}-\mathbf{s}) = \rho_{\mathbf{S}}(\mathbf{s}) \cdot \rho_{\mathbf{S}}(\mathbf{c}) \cdot (\rho_{\mathbf{S}}(-2\mathbf{s} \cdot \mathbf{c})) + \rho_{\mathbf{S}}(2\mathbf{s} \cdot \mathbf{c})) = 2\rho_{\mathbf{S}}(\mathbf{s}) \cdot \rho_{\mathbf{S}}(\mathbf{c}) \cdot \cosh(\pi \cdot \mathbf{s} \cdot \mathbf{c}/\varsigma^2) = (\rho_{\varsigma}(\mathbf{s}) + \rho_{\varsigma}(-\mathbf{s})) \cdot \rho_{\varsigma}(\mathbf{c}) \cdot \cosh(\pi \cdot (\mathbf{S}^{-1} \cdot \mathbf{s}) \cdot (\mathbf{S}^{-1} \cdot \mathbf{c})) \geq (\rho_{\varsigma}(\mathbf{s}) + \rho_{\varsigma}(-\mathbf{s})) \cdot \rho_{\varsigma}(\mathbf{c})$  and then by adding over all pairs  $-\mathbf{s}$ ,  $\mathbf{s}$  in S, the result follows.

Lemma 2.12 (Adapted [MR07], [GMPW20, Lem 2.6 eprint]). Let  $\Lambda$  be an n-dimensional lattice and  $\sqrt{\Sigma} \geq \eta_{\varepsilon}(\Lambda)$ . Then

$$\rho_{\sqrt{\Sigma}}(\Lambda) \in [1 - \varepsilon, 1 + \varepsilon] \cdot \frac{\det(\sqrt{\Sigma})}{\det(\Lambda)}$$

*Proof.* If  $\sqrt{\Sigma} \geq \eta_{\varepsilon}(\Lambda)$  then  $1 \geq \eta_{\varepsilon}(\sqrt{\Sigma}^{-1} \cdot \Lambda)$  and the statement follows.

**Lemma 2.13.** Let  $\Lambda \subset \mathbb{R}^n$  be a full rank lattice. Let  $R \geq \sqrt{n} \cdot \eta_{1/2}(\Lambda)$  be a real. Then

$$|\Lambda \cap B(0,R)| \le 1.5 \cdot \exp(\pi \cdot n) \cdot \left(\frac{R}{\sqrt{n}}\right)^n \cdot \frac{1}{\det(\Lambda)}.$$

*Proof.* Let  $\varsigma = R/\sqrt{n}$ . By Lemma 2.12, it holds that

$$\rho_{\varsigma}(\Lambda \cap B(0,R)) \le \rho_{\varsigma}(\Lambda) \le 1.5 \frac{\varsigma^n}{\det \Lambda},$$

the result follows from the fact that  $\rho_{\varsigma}(\Lambda \cap B(0,R))$  is a sum of  $|\Lambda \cap B(0,R)|$  terms, all greater than  $\rho_{\varsigma}(R) = \exp(-\pi n)$ .

Lemma 2.14 (simplified [ALLW25, Theorem 9 (eprint)]). Let g, k, m, n, q be positive integers,  $a \ge 1$  be a real number, and  $\Sigma$  be positive semi-definite in  $\mathbb{R}^{dm \times dm}$  satisfying

$$-k, n \leq m, d \cdot (m-k) \geq \Omega(\lambda),$$

- q be an unramified prime s.t.  $\langle q \rangle = \prod_{j=1}^g \mathfrak{q}_j$  with norm  $\mathcal{N}(\mathfrak{q}_j) = q^{d/g}$  in  $\mathcal{O}_{\mathcal{K}}$ ,
- $-ng/q^{d(m-n+1)/g} \leq \text{negl}(\lambda)$ , and there exists  $\epsilon \leq \text{negl}(\lambda)$  so that

$$\begin{split} & \max\left\{\eta_A, 2\sqrt{d}\cdot (a^kq^n)^{1/(m-k)}\right\} \leq s_{\min}\left(\sqrt{\Sigma}\right), \\ & s_{\max}\left(\sqrt{\Sigma}\right) \leq \min\left\{q^{1/g}/\sqrt{m}, a\cdot s_{\min}\left(\sqrt{\Sigma}\right)\right\}, \end{split}$$

for some  $\eta_A > 8 d\sqrt{m} \cdot q^{n/m+2/(d \cdot m)}$ .

Then the following distributions are statistically close in  $\lambda$ :

$$\left\{ \left(\mathbf{A},\mathbf{U}\right) \, \left| \begin{array}{l} \mathbf{A} \leftarrow (\mathcal{O}_{\mathcal{K}}/q)^{n \times m} \\ \mathbf{X} \leftarrow (\mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}),\sqrt{\Sigma}})^k \end{array} \right\} \approx_s \left\{ \left(\mathbf{A},\mathbf{U}\right) \, \left| \begin{array}{l} \mathbf{X} \leftarrow (\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^m},\sqrt{\Sigma}})^k \\ \mathbf{A} \leftarrow (\mathcal{O}_{\mathcal{K}}/q)^{n \times m} : \mathbf{A} \cdot \mathbf{X} \equiv \mathbf{0} \end{array} \right\}.$$

**Lemma 2.15 (simplified [ALLW25, Lemma 26 (eprint)]).** Let g, k, m, n, q be positive integers,  $a \ge 1$  be a real number,  $\mathbf{A} \in (\mathcal{O}_{\mathcal{K}}/q)^{n \times m}$  be primitive, and  $\mathbf{\Sigma} \in \mathbb{R}^{dm \times dm}$  be positive semi-definite satisfying

- $-k, n \leq m, d \cdot (m-k) \geq \Omega(\lambda),$
- q is an unramified prime s.t.  $\langle q \rangle = \prod_{j=1}^g \mathfrak{q}_j$  with  $\mathcal{N}(\mathfrak{q}_j) = q^{d/g}$ , and
- there exists  $\epsilon \leq \text{negl}(\lambda)$  that

$$\max \left\{ \eta_{\epsilon}(\Lambda_{q}^{\perp}(\mathbf{A})), 2\sqrt{d} \cdot (a^{k}q^{n})^{1/(m-k)} \right\} \leq s_{\min}\left(\sqrt{\Sigma_{i}}\right),$$

$$s_{\max}\left(\sqrt{\Sigma_{i}}\right) \leq \min \left\{ q^{1/g}/\sqrt{m}, a \cdot s_{\min}\left(\sqrt{\Sigma_{i}}\right) \right\}.$$

Then the columns of  $\mathbf{X} \leftarrow (\mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}),\sqrt{\Sigma}})^k$  are  $(\mathcal{O}_{\mathcal{K}}/q)$ -linearly independent with overwhelming probability in  $\lambda$ .

Lemma 2.16 ([ALLW25, Lemma 3 (eprint)]; Generalisation of [LPR13, Theorem 4.1]). Let n, m, q be positive integers with  $n \leq m \leq \operatorname{poly}(d)$ . For  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times m}$ , with probability  $1 - 2^{-\Omega(dm)}$  we have  $\eta_{2^{-\Omega(dm)}}(\Lambda_q^{\perp}(\mathbf{A})) \leq 8d\sqrt{m} \cdot q^{n/m + 2/(d \cdot m)}$ .

**Lemma 2.17** ([LSS14, Lemma 4.2]). For any m-dimensional lattice  $\Lambda \subseteq \mathbb{R}^m$  and rank m matrix  $\sqrt{\Sigma} \in \mathbb{R}^{m \times m}$ , let  $P = \mathcal{D}_{\Lambda, \sqrt{\Sigma}, \mathbf{w}}$  and  $Q = \mathcal{D}_{\Lambda, \sqrt{\Sigma}, \mathbf{z}}$  for some fixed  $\mathbf{w}, \mathbf{z} \in \mathbb{R}^m$ . If  $\mathbf{w}, \mathbf{z} \in \Lambda$ , let  $\epsilon = 0$ . Otherwise, fix  $\epsilon \in (0, 1)$  and assume that  $s_{\min}(\sqrt{\Sigma}) \geq \eta_{\epsilon}(\Lambda)$ . Then  $\text{RD}(P||Q) \leq (\frac{1+\epsilon}{1+\epsilon})^2 \cdot \exp(2\pi ||\mathbf{w} - \mathbf{z}||^2/s_{\min}(\sqrt{\Sigma})^2)$ .

**Lemma 2.18.** Let K be a number field of degree d. Let  $n, m, k \in \mathbb{N}$ . Let  $\mathbf{A} \in \mathcal{R}_q^{n \times m}$  be such that  $\Lambda_q^{\perp}(\mathbf{A})$  is full-rank, and let  $\sqrt{\Sigma} \in \mathbb{R}^{dm \times dm}$  satisfy  $s_{\min}(\sqrt{\Sigma}) \geq \eta_{\epsilon}(\Lambda_q^{\perp}(\mathbf{A}))$  for some  $0 < \epsilon \leq 2^{-2k}$ . Let  $\mathbf{e}_i \in \mathcal{R}^m$  be the i-th unit vector for  $i \in [k]$ . The Rényi divergence between the distributions  $P := \prod_{i \in [k]} \mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}),\sqrt{\Sigma}}$  and  $Q := \prod_{i \in [k]} \mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}),\sqrt{\Sigma},\mathbf{e}_i}$  is upper-bounded by

$$\operatorname{RD}(P||Q) \le 5 \exp\left(\frac{2\pi \cdot dk}{s_{\min}(\sqrt{\Sigma})^2}\right).$$

*Proof.* Denote by  $1 + \delta := (1 + \epsilon)/(1 - \epsilon)$ , equivalently  $\delta = 2\epsilon/(1 - \epsilon)$ . Using the multiplicative property of Rényi divergence, we have

$$\operatorname{RD}(P\|Q) \leq \prod_{i \in [k]} (1+\delta)^2 \cdot \exp\left(2\pi \|\mathbf{e}_i\|^2 / s_{\min}(\sqrt{\Sigma})^2\right) \\
\leq (1+(2^{2k}-1)\cdot \delta) \cdot \exp\left(\frac{2\pi \cdot dk}{s_{\min}(\sqrt{\Sigma})^2}\right) \leq 5 \cdot \exp\left(\frac{2\pi \cdot dk}{s_{\min}(\sqrt{\Sigma})^2}\right),$$

where the first inequality is by Lemma 2.17, the second uses  $(1+x)^r \le 1+(2^r-1)x$  for any  $x \in [0,1]$  and  $r \in \mathbb{R} \setminus (0,1)$ , and the last is by  $2^{2k} \cdot \delta = 2^{2k} \cdot 2\epsilon/(1-\epsilon) \le 4$ .

# 2.4 Cryptographic Assumptions

# Definition 2.1 (SIS Problem).

Let the parameters params =  $(\mathcal{O}_{\mathcal{K}}, q, n, m, \beta)$  be parametrised by  $\lambda$ , where  $\mathcal{O}_{\mathcal{K}}$  is the ring of integers of a number field  $\mathcal{K}$ , q, n, m are positive integers and  $\beta > 0$ . Write  $\mathcal{R}_q := \mathcal{O}_{\mathcal{K}}/(q \cdot \mathcal{O}_{\mathcal{K}})$ . A SIS<sub>params</sub> problem asks the following: Given a uniform matrix  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times m}$ , find a vector  $\mathbf{u} \in \mathcal{O}_{\mathcal{K}}^m$  such that

(1) 
$$\mathbf{A} \cdot \mathbf{u} = \mathbf{0} \mod q$$
 and (2)  $0 < \|\mathbf{u}\| < \beta$ .

For any PPT A, its advantage  $Adv_A^{SIS_{params}}(\lambda)$  against this problem is equal to the probability that it solves the problem.

# Definition 2.2 (LWE Problem).

Let the parameters params =  $(\mathcal{O}_{\mathcal{K}}, q, n, m, \chi)$  be parametrised by  $\lambda$ , where  $\mathcal{O}_{\mathcal{K}}$  is the ring of integers of a number field  $\mathcal{K}$ , q, n, m are positive integers and  $\chi$  is a distribution over  $\mathcal{O}_{\mathcal{K}}$ . Write  $\mathcal{R}_q := \mathcal{O}_{\mathcal{K}}/(q \cdot \mathcal{O}_{\mathcal{K}})$ . An LWE<sub>params</sub> problem asks the following: Given a uniform matrix  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times m}$ , and a vector  $\mathbf{b} \in \mathcal{R}_q^m$  which is either

(1) 
$$\mathbf{b}^{\mathsf{T}} = \mathbf{s}^{\mathsf{T}} \mathbf{A} + \mathbf{e}^{\mathsf{T}} \bmod q \text{ where } \mathbf{s} \leftarrow \mathcal{R}_q^n, \ \mathbf{e} \leftarrow \chi^m,$$

or (2) 
$$\mathbf{b} \leftarrow \mathcal{R}_{q}^{m}$$
 is uniformly random,

distinguish which one is the case. For any PPT  $\mathcal{A}$ , its advantage  $Adv_{\mathcal{A}}^{LWE_{params}}(\lambda)$  against this problem is equal to the absolute difference in the probability of its response under these two cases.

# Definition 2.3 (k-SIS Problem).

Let the parameters params =  $(\mathcal{O}_{\mathcal{K}}, q, n, m, k, \Sigma, \beta)$  be parametrised by  $\lambda$ , where  $\mathcal{O}_{\mathcal{K}}$  is the ring of integers of a number field  $\mathcal{K}$  of degree d, q, n, m, k are positive integers with k < m,  $\Sigma \in \mathbb{R}^{dm \times dm}$  and  $\beta > 0$ . Write  $\mathcal{R}_q := \mathcal{O}_{\mathcal{K}}/(q \cdot \mathcal{O}_{\mathcal{K}})$ . A k-SIS<sub>params</sub> problem asks the following: Given a uniform matrix  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times m}$  and k vectors  $\mathbf{X} = (\mathbf{x}_1, \dots, \mathbf{x}_k) \in \mathcal{O}_{\mathcal{K}}^{m \times k}$  where  $\mathbf{x}_i \leftarrow \mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}), \sqrt{\Sigma}}$  for each  $i \in [k]$ , find a vector  $\mathbf{u} \in \mathcal{O}_{\mathcal{K}}^m$  such that

(1) 
$$\mathbf{A} \cdot \mathbf{u} = \mathbf{0} \mod q$$
, (2)  $0 < \|\mathbf{u}\| < \beta$ , and (3)  $\mathbf{u} \notin \mathcal{K}\text{-span}(\mathbf{X})$ .

For any PPT A, its advantage  $Adv_A^{k-SIS_{params}}(\lambda)$  against this problem is equal to the probability that it solves the problem.

### Definition 2.4 (k-LWE Problem).

Let the parameters params =  $(\mathcal{O}_{\mathcal{K}}, q, n, m, k, \Sigma, \chi)$  be parametrised by  $\lambda$ , where  $\mathcal{O}_{\mathcal{K}}$  is the ring of integers of a number field  $\mathcal{K}$  of degree d, q, n, m, k are positive integers with k < m,  $\Sigma \in \mathbb{R}^{dm \times dm}$  and  $\chi$  is a distribution over  $\mathcal{O}_{\mathcal{K}}$ . Write  $\mathcal{R}_q \coloneqq \mathcal{O}_{\mathcal{K}}/(q \cdot \mathcal{O}_{\mathcal{K}})$ . An LWE<sub>params</sub> problem asks the following: Given a uniform matrix  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times m}$ , k vectors  $\mathbf{X} = (\mathbf{x}_1, \dots, \mathbf{x}_k) \in \mathcal{O}_{\mathcal{K}}^{m \times k}$  where  $\mathbf{x}_i \leftarrow \mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}), \sqrt{\Sigma}}$  for each  $i \in [k]$ , and a vector  $\mathbf{b} \in \mathcal{R}_q^m$  which is either

(1) 
$$\mathbf{b}^{\mathsf{T}} = \mathbf{s}^{\mathsf{T}} \mathbf{A} + \mathbf{e}^{\mathsf{T}} \mod q \text{ where } \mathbf{s} \leftarrow \mathcal{R}_q^n, \ \mathbf{e} \leftarrow \chi^m,$$

or (2) 
$$\mathbf{b}^{\mathsf{T}} = \mathbf{d} + \mathbf{e}^{\mathsf{T}} \bmod q \text{ where } \mathbf{d} \leftarrow \{\mathbf{d} \in \mathcal{R}^m : \mathbf{d}^{\mathsf{T}} \mathbf{X} = \mathbf{0} \bmod q\},$$

distinguish which one is the case. For any PPT  $\mathcal{A}$ , its advantage  $\mathrm{Adv}_{\mathcal{A}}^{k-\mathrm{LWE}_{\mathsf{params}}}(\lambda)$  against this problem is equal to the absolute difference in the probability of its response under these two cases.

# 3 Generating module lattices with Gaussian vectors

In this section, we study the surjectivity of Gaussian matrices. For a matrix where each column is sampled from an  $\mathcal{O}_{\mathcal{K}}$ -module  $\mathcal{M}$ , we say that a matrix is surjective if  $\mathcal{O}_{\mathcal{K}}$ -linear combinations of its columns generate the module  $\mathcal{M}$ .

#### 3.1 Surjective matrices over finite fields

We first establish some facts about random matrices over finite fields  $\mathbb{F}$  to be used in subsequent proofs. First, extending the above notion of surjectivity to  $\mathbb{F}$ -vector spaces, we show that a random matrix over  $\mathbb{F}$  is surjective with overwhelming probability in its dimensions if each entry of the matrix is independently distributed and has constant min-entropy.

**Lemma 3.1.** Let  $r, u \ge 0$  be integers. Let  $\mathbb{F}$  be a finite field, and  $(D_{i,j})_{1 \le i \le r, 1 \le j \le m}$  be distributions over  $\mathbb{F}$  such that there exists  $\alpha \in (0,1)$  such that for any i,j,  $\max_{x \in \mathbb{F}} D_{i,j}(x) \le \alpha$ . If  $\mathbf{M} \in \mathbb{F}^{r \times m}$  is a random matrix such that

- all the entries of  $\mathbf M$  are independent random variables, and
- the (i,j)-th entry of **M** is sampled from  $D_{i,j}$ ,

then it holds that

$$\Pr(\mathbf{M} \cdot \mathbb{F}^m = \mathbb{F}^r) \ge 1 - r \cdot \alpha^{m-r+1}.$$

*Proof.* Let H be a subspace of  $\mathbb{F}^m$  of dimension k and  $i = 1, \ldots, d$ , it holds that the probability of a vector in  $\mathbb{F}^m$  sampled from  $D_{i,1} \times \ldots \times D_{i,m}$  belonging to this subspace is upper bounded by  $\alpha^{m-k}$  [NP20, Generalisation of Lemma 2.3]. As  $\mathbb{F}$  is a field, the matrix  $\mathbf{M}$  is surjective if and only if  $\mathbf{M}^T$  is injective. We now

lower-bound the probability of  $\mathbf{M}^T$  being injective when  $\mathbf{M}$  is built row by row. Let  $\mathbf{M}_i$  be the matrix composed of the first i rows of  $\mathbf{M}$ . For  $i=1,\ldots,r$ , the matrix  $\mathbf{M}_i^T$  is injective if and only if  $\mathbf{M}_{i-1}^T$  is injective and the i-th row of  $\mathbf{M}$  is not in the span of the i-1 previous ones, which has dimension i-1. The probability of  $\mathbf{M}^T$  being injective is then lower-bounded by

$$\prod_{i=1}^{r} (1 - \alpha^{m-i+1}) = \prod_{i=m-r+1}^{m} (1 - \alpha^{i}) \ge (1 - \alpha^{m-r+1})^{r} \ge 1 - r \cdot \alpha^{m-r+1}. \quad \Box$$

The next lemma shows the equivalence between a matrix  $\mathbf{B}$  with columns in  $\mathcal{M}$  being surjective and the simpler condition of  $(\mathbf{B} \bmod \mathfrak{p})$  being surjective for prime ideals  $\mathfrak{p}$  of  $\mathcal{O}_{\mathcal{K}}$ .

**Lemma 3.2.** For any integers  $r, u \ge 1$ , a matrix  $\mathbf{B} \in \mathcal{M}^m$  satisfies  $\mathbf{B} \cdot \mathcal{O}_{\mathcal{K}}^m = \mathcal{M}$  if and only if  $(\mathbf{B} \bmod \mathfrak{p})$  is surjective in  $\mathcal{M}/\mathfrak{p}\mathcal{M}$  for any prime ideal  $\mathfrak{p}$  of  $\mathcal{O}_{\mathcal{K}}$ .

Proof (Adaptation of the proof of [NP20, Lemma 2.1]). If **B** is surjective in  $\mathcal{M}$ , then it is clearly surjective modulo any prime  $\mathfrak{p}$ . Now we assume that **B** mod  $\mathfrak{p}$  is surjective modulo every  $\mathfrak{p}$ . Let  $\mathcal{M}' = \mathbf{B} \cdot \mathcal{O}_{\mathcal{K}}^m$ . Let  $\mathfrak{p}$  be any ideal, we will show that  $\mathfrak{p}$  does not divide  $\det(\mathcal{M}')/\det(\mathcal{M})$ . Since the matrix **B** mod  $\mathfrak{p}$  is surjective, there exists a square submatrix  $\mathbf{B}'$  of  $\mathbf{B}$  which is invertible modulo  $\mathfrak{p}$ , and hence full-rank over  $\mathcal{K}$ . Let  $\mathcal{M}'' = \mathbf{B}' \cdot \mathcal{O}_{\mathcal{K}}^r \subset \mathcal{M}'$ . Assume that the ideal  $\mathfrak{p}$  divides  $\det(\mathcal{M}')/\det(\mathcal{M})$ , then it also divides  $\det(\mathcal{M}'')/\det(\mathcal{M})$  by inclusion, which is equivalent to that  $\det_{\mathcal{K}}(\mathbf{B}') \in \det(\mathcal{M}) \cdot \mathfrak{p}$ , which contradicts the fact that  $\mathbf{B}'$  is invertible in  $\mathcal{M}/\mathfrak{p}\mathcal{M}$  by Lemma 2.2. Since  $\mathfrak{p}$  can be any ideal, it holds that  $\det(\mathcal{M}') = \det(\mathcal{M})$ , and hence that  $\mathbf{B}$  is surjective in  $\mathcal{M}$ .

#### 3.2 Surjective matrices over free modules

In this subsection, we prove that Gaussian matrices of size  $r \times m$  over some full rank modules  $\mathcal{M}$  are surjective over  $\mathcal{M}$  with high probability when  $m \geq 2r$ . The main idea of the proof is to consider the matrix modulo all the primes  $\mathfrak{p}$  of the number field, showing that with high probability (using Lemmas 3.1 and 3.2), the matrix modulo  $\mathfrak{p}$  is surjective over  $\mathcal{O}_{\mathcal{K}}/\mathfrak{p}$ . We first prove it in the case of free modules in Lemma 3.4, where the main contribution of this subsection is, and then generalise it in Theorem 3.1. We first prove that the Gaussian distribution over some ideal I modulo the ideal  $I \cdot \mathfrak{p}$  satisfies the distribution hypothesis of Lemma 3.1.

Lemma 3.3 (Generalised from [JLWG25, Corollary 7]). Let  $\varepsilon \in (0, 1/\sqrt{2})$ , I be a (possibly fractional) ideal of  $\mathcal{O}_{\mathcal{K}}$ ,  $B \geq 1$ ,  $c \in \mathcal{K}_{\mathbb{R}}$  be a centre and  $\varsigma \geq B \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \mathcal{N}(I)^{1/d} \cdot \eta_{\varepsilon}^{(d)}$ . For any integral ideal  $\mathfrak{a} \subset \mathcal{O}_{\mathcal{K}}$ , we define  $\mathcal{D}_{I,\varsigma,c}^{(\mathfrak{a})}$  the distribution  $\mathcal{D}_{I,\varsigma,c}$  modulo  $\mathfrak{a} \cdot I$ . Then for any  $a \in I/(I \cdot \mathfrak{a})$ , it holds that

$$\mathcal{D}_{I,\varsigma,c}^{(\mathfrak{a})}(a) \leq (1+4\varepsilon) \cdot \begin{cases} \frac{1}{\mathcal{N}(\mathfrak{a})} & \text{if } \mathcal{N}(\mathfrak{a}) \leq B^d \\ \frac{1}{B^d} & \text{else} \end{cases}.$$

**Lemma 3.4 (Assuming GRH).** Let  $r \geq 1$  and  $m \geq 2r$  be integers. Let  $\mathbf{C} = [\mathbf{c}_1, \dots, \mathbf{c}_m] \in \mathbb{R}^{rd \times rm}$  a centre matrix and  $\varsigma > 0$ . Let  $\mathbf{S} \geq \sqrt{2}\varsigma \mathbf{I}$  or equal to  $\varsigma \mathbf{I}$ . Let  $\mathcal{M} \subset \mathcal{K}^r$  be a free module of rank r. Let  $\varepsilon \in (0, 1/4)$ . There exists an absolute constant c > 1 such that for any  $B > \log(\Delta_{\mathcal{K}})^{c/d}$ , assuming that the following two inequalities hold

$$B^d \ge r \cdot d^2 \ln(dr \cdot \varsigma)$$
 and  $\varsigma \ge B \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M}),$ 

and  $\mathbf{X} = [\mathbf{x}_1, \dots, \mathbf{x}_m] \in \mathcal{K}^{r \times m}$  such that  $\mathbf{x}_i \leftarrow \mathcal{D}_{\mathcal{M}, \mathbf{S}, \mathbf{c}_i}$ , then

$$\Pr_{\mathbf{X}}[\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^{m} \neq \mathcal{M}] \leq \frac{m \cdot \varepsilon}{1 - \varepsilon} + r \cdot \left(\frac{B}{2}\right)^{-(r+1)d} + r \cdot 2^{-dr} + (1 + 4\varepsilon)^{m-2r+1} \cdot r \cdot P_{\mathcal{K}}(m - 2r + 1),$$

where  $P_{\mathcal{K}}(s) = \sum_{\mathfrak{p}} \mathcal{N}(\mathfrak{p})^{-s}$  is the prime zeta function of  $\mathcal{K}$ .

*Proof.* We abuse the notation and write the distribution of **X** as  $\mathcal{D}_{\mathcal{M}^m, \mathbf{S}, \mathbf{C}}$ .

From free modules to  $\mathcal{O}_{\mathcal{K}}^r$ . Let  $\mathbf{B} \in \mathcal{K}^{r \times r}$  generating  $\mathcal{M}$ . It holds that the distribution of the columns of  $\mathbf{X}$  is the same as the distribution  $\mathbf{B} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r \times m}, \mathbf{S}', \mathbf{C}'}$  where  $\mathbf{S}' = \mathbf{B}^{-1} \cdot \mathbf{S}$  and  $\mathbf{C}' = \mathbf{B}^{-1} \cdot \mathbf{C}$ . Note that the condition on  $\mathbf{S}$  implies that  $\mathbf{S}' \geq B \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}}^r)$ . We can then assume without loss of generality than  $\mathcal{M} = \mathcal{O}_{\mathcal{K}}^r$ .

From elliptic to circular Gaussian. Let  $\Sigma = \mathbf{S}\mathbf{S}^T - \varsigma^2\mathbf{I}$ , note that  $\Sigma$  is definite positive, let  $\sqrt{\Sigma}$  be a fixed square root of  $\Sigma$ . Now, let  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r \times m}, \mathbf{S}, \mathbf{C}}$  be the distribution  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r \times m}, \varsigma, \mathbf{C} - \mathbf{B}}$  where  $\mathbf{B} \leftarrow \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r \times m}, \sqrt{\Sigma}}$ . By [GMPW20, Theorem 4.5], the condition on  $\mathbf{S}$  and Lemma B.2, it holds that the statistical distance between  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}, \mathbf{C}}$  and  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}, \mathbf{C}}$  is less than  $m \cdot \varepsilon/(1 - \varepsilon)$ . It then holds that

$$\begin{split} \Pr_{\mathbf{X}}[\mathbf{X}\cdot\mathcal{O}_{\mathcal{K}}^{m} \neq \mathcal{O}_{\mathcal{K}}^{r}] &\leq m \cdot \frac{\varepsilon}{1-\varepsilon} \\ &+ \sum_{\mathbf{B}\in\mathcal{O}_{\mathcal{K}}^{r\times m}} \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r\times m},\sqrt{\mathbf{\Sigma}'}}(\mathbf{B}) \Pr_{\mathbf{A}\leftarrow\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r\times m},\varsigma,\mathbf{C}-\mathbf{B}}}(\mathbf{A}\cdot\mathcal{O}_{\mathcal{K}}^{m} \neq \mathcal{O}_{\mathcal{K}}^{r}). \end{split}$$

Bound for circular Gaussian. We now assume that  $\mathbf{S} = \varsigma \cdot \mathbf{I}$  with  $\varsigma \geq B \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}}^r)$ . Note that in that case, the coefficients of  $\mathbf{X}$  are all independents and sampled from  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}},\varsigma,c}$  for some  $c \in \mathcal{K}$ , so by Lemma 3.3, Lemma 3.1 holds for any submatrix of  $\mathbf{X}$ . Write  $\mathbf{X} = [\mathbf{X}_1 | \mathbf{X}_2]$  with  $\mathbf{X}_1 \in \mathcal{O}_{\mathcal{K}}^{r \times r}$ , and let  $\mathfrak{p}$  be a prime ideal of norm  $\geq B^d$ . Then by Lemma 3.2,  $\mathbf{X}_1$  is invertible modulo  $\mathfrak{p}$  with probability  $\geq 1 - r(B/2)^{-(r+1)d}$  and in that case, in particular  $\det_K(\mathbf{X}_1) \neq 0$ , we assume this is the case for the rest of the proof. By [Ban93, Lemma 1.5], with probability  $1 - r \cdot 2^{dr}$ , the norm of every column of  $\mathbf{X}_1$  is bounded by  $\sqrt{rd} \cdot \varsigma$ , and hence its determinant is bounded by  $(\sqrt{rd}\varsigma)^{rd}$ . Let  $\mathcal{P}$  be the set of prime ideals dividing  $\det_K(\mathbf{X}_1)$ , by Lemma B.1 it holds that  $|\mathcal{P}| \leq rd^2 \cdot \log(\sqrt{rd}\varsigma) =: l$ . We now prove

that **X** is invertible modulo every prime ideal. For any prime  $\mathfrak p$  not in  $\mathcal P$ ,  $\mathbf X_1$  is invertible modulo  $\mathfrak p$ , and so is **X**. For all  $\mathfrak p \in \mathcal P$ , by Lemma 3.1 it holds that  $\mathbf X_2$  is invertible mod  $\mathfrak p$  with probability  $\geq 1 - r(1 - 4\varepsilon)^{m-2r+1} \cdot \alpha$  for  $\alpha = B^{-(m-2r+1)d}$  if  $\mathcal N(\mathfrak p) \geq B^d$  and  $\alpha = \mathcal N(\mathfrak p)^{-(m-2r+1)}$  else.

Let us order the primes of  $\mathcal{O}_{\mathcal{K}}$  as  $\{\mathfrak{p}_1,\mathfrak{p}_2,\ldots\}$  and  $\mathcal{P} = \{\mathfrak{q}_1,\ldots,\mathfrak{q}_l\}$ , both in increasing norm order. It holds that  $\mathcal{N}(\mathfrak{p}_i) \leq \mathcal{N}(\mathfrak{q}_i)$  for any  $i=1,\ldots,l$ . We let B be a real such that the first l prime ideals of  $\mathcal{O}_{\mathcal{K}}$  have their norm  $\leq B^d$ . The existence of B will be proven at the end of the proof. Under the condition that  $\mathbf{X}_1$  is invertible and has bounded columns, by the union bound it holds that

$$\Pr(\mathbf{X} \text{ not surjective}) \leq \sum_{i=1}^{l} \Pr(\mathbf{X}_{2} \text{ not surjective mod } \mathfrak{q}_{i})$$

$$\leq r \cdot (1 + 4\varepsilon)^{m-2r+1} \cdot \left(\sum_{i=1}^{a} \frac{1}{\mathcal{N}(\mathfrak{q}_{i})^{m-2r+1}} + \sum_{i=a+1}^{l} \frac{1}{B^{d(m-2r+1)}}\right)$$

$$\leq r \cdot (1 + 4\varepsilon)^{m-2r+1} \cdot \sum_{i=1}^{l} \frac{1}{\mathcal{N}(\mathfrak{p}_{i})^{m-2r+1}} \leq r \cdot (1 + 4\varepsilon)^{m-2r+1} \cdot P_{\mathcal{K}}(m-2r+1).$$

Now let us prove that B satisfying our condition exists and give a lower bound on it. The condition is that B is such that the first l primes of  $\mathcal{O}_{\mathcal{K}}$  are of norm less than  $B^d$ . By [BS96, Theorem 8.7.4], the GRH implies that for  $X \geq \log(\Delta_{\mathcal{K}})^{O(1)}$ , the number of prime ideals  $\pi_{\mathcal{K}}(X)$  of norm less than X is less than  $1.1 \cdot X / \ln(X)$ . This implies that as long as  $B^d \geq \log(\Delta_{\mathcal{K}})^{O(1)}$ , then  $\pi_{\mathcal{K}}(B^d) \leq B^d$ , then as long as

$$l = r \cdot d^2 \ln(dr \cdot \varsigma) \le B^d,$$

the norm of the first l prime ideals is less than  $B^d$ .

# 3.3 Generating non-free modules

From now on, we fix a full-rank (not necessarily free) module  $\mathcal{M} \subset \mathcal{K}^r$  given by a pseudo-basis  $\mathcal{M} = \sum_{i=1}^r \mathbf{b}_i \cdot \mathcal{I}_i$  where  $\mathcal{I}_i$  are fixed fractional ideals of  $\mathcal{O}_{\mathcal{K}}$  and  $\mathbf{b}_i \in \mathcal{K}^r$ . Our goal is to generalise Lemma 3.4 for non-free  $\mathcal{M}$ . For  $\mathcal{M} = \mathcal{O}_{\mathcal{K}}^r$ , we proved that if sufficiently many vectors are sampled from  $\mathcal{O}_{\mathcal{K}}^r$ , then the matrix they form is surjective over  $(\mathcal{O}_{\mathcal{K}}/\mathfrak{p})^r$  for any ideal  $\mathfrak{p}$ , and then surjective over  $\mathcal{O}_{\mathcal{K}}^r$ . We will prove the same result with  $\mathcal{M}/\mathfrak{p}\mathcal{M}$  for not necessarily free  $\mathcal{M}$ .

Let  $\mathfrak{p}$  be a prime ideal with residue field  $\mathbb{F} = \mathcal{O}_{\mathcal{K}}/\mathfrak{p}$ . The set  $\mathcal{M}/\mathfrak{p}\mathcal{M}$  has a natural structure of  $\mathbb{F}$ -vector space of dimension r. For any matrix  $\mathbf{B} \in \mathcal{M}^m$ , we define  $(\mathbf{B} \bmod \mathfrak{p})$  the associated matrix on  $(\mathcal{M}/\mathfrak{p}\mathcal{M})^m$ .

**Theorem 3.1 (Assuming GRH).** Let  $r \geq 1$  and  $m \geq 2r$  be integers. Let  $\mathbf{C} = [\mathbf{c}_1, \dots, \mathbf{c}_m] \in \mathbb{R}^{rd \times rm}$  a centre matrix and  $\varsigma > 0$ . Let  $\mathbf{S} \geq \sqrt{2}\varsigma \mathbf{I}$  or equal to  $\varsigma \mathbf{I}$ . Let  $\mathcal{M} \subset \mathcal{K}^r$  be a module of rank r. Let  $\varepsilon \in (0, 1/4)$ . There exists an absolute constant c > 1 such that for any  $B > \log(\Delta_{\mathcal{K}})^{c/d}$ , assuming that the following two inequalities hold

$$B^d \ge r \cdot d^2 \ln(dr \cdot \varsigma)$$
 and  $\varsigma \ge B \cdot \Delta_{\kappa}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M}),$ 

and  $\mathbf{X} = [\mathbf{x}_1, \dots, \mathbf{x}_m] \in \mathcal{K}^{r \times m}$  such that  $\mathbf{x}_i \leftarrow \mathcal{D}_{\mathcal{M}, \mathbf{S}, \mathbf{c}_i}$ , then

$$\Pr_{\mathbf{X}}[\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^{m} \neq \mathcal{M}] \leq \frac{m \cdot \varepsilon}{1 - \varepsilon} + r \cdot \left(\frac{B}{2}\right)^{-(r+1)d} + r \cdot 2^{-dr} + (1 + 4\varepsilon)^{m-2r+1} \cdot r \cdot P_{\mathcal{K}}(m - 2r + 1),$$

where  $P_{\mathcal{K}}(s) = \sum_{\mathfrak{p}} \mathcal{N}(\mathfrak{p})^{-s}$  is the prime zeta function of  $\mathcal{K}$ .

Proof. By the same arguments than the proof of Lemma 3.4, we can assume without loss of generality that  $\mathcal{M} = \bigoplus_{i=1}^r I_i$  with  $I_1, \ldots, I_r$  ideals, and  $\mathbf{S} = \varsigma \mathbf{I}_{rd}$  with  $\varsigma \geq B \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \eta_{\varepsilon}(\bigoplus_{1 \leq i \leq d} I_i)$ . In particular, it holds that  $\varsigma \geq B \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \eta_{\varepsilon}(I_i)$  for any  $1 \leq i \leq r$ , in that case, Lemma 3.3 implies that Lemma 3.1 holds for any submatrix of  $\mathbf{X}$ , we can then run the exact same proof as the one of Lemma 3.4.

Corollary 3.1 (Prime-Power Cyclotomics, Assuming GRH). Let  $p^k > 2$  a prime power, and  $\mathcal{K} = \mathbb{Q}(\zeta_{p^k})$  be the  $p^k$ -cyclotomic number field of degree  $d = (p-1) \cdot p^{k-1}$ . Let  $0 < \varepsilon < 2^{-6}$ . Let  $\mathcal{M} \subset \mathcal{K}^r$  be a full rank module. Then for any  $\mathbf{S} \in \mathbb{R}^{dr \times dr}$ ,  $m \geq 1$  and centre  $\mathbf{C} \in \mathcal{K}^{r \times m}_{\mathbb{R}}$  such that

$$\mathbf{S} \geq 9\sqrt{2} \cdot r^{\frac{1}{d}} \cdot \log\left(rd \cdot \eta_{\varepsilon}(\mathcal{M})\right)^{\frac{1}{d}} \cdot d \cdot \eta_{\varepsilon}(\mathcal{M}) \ \ and \ \ m \geq 2r + \frac{\log(1/\varepsilon) + 1 + \log(2d)}{\log(p/1.4)},$$

it holds that

$$\Pr_{\mathbf{X} \leftarrow \mathcal{D}_{\mathcal{K}}^m \mid \mathbf{S} \mid \mathbf{C}} (\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^m \neq \mathcal{M}) \leq m \cdot (\varepsilon + 2^{-dr}).$$

*Proof.* Follows from Theorem 3.2 and Lemmas B.3 and 3.6.

# 3.4 Asymptotic behavior of $P_{\mathcal{K}}$

The upper bounds in Lemma 3.3 and Theorem 3.1 are described in terms of the prime zeta function  $P_{\mathcal{K}}$ . In the following, we upper bound  $P_{\mathcal{K}}$  in terms of the smallest prime ideal norm  $N_{\mathcal{K}}$ , for which explicit the explicit value is known when  $\mathcal{K}$  is a cyclotomic field of prime-power conductor.

**Lemma 3.5.** For any number field K such that there is exactly  $a_K$  prime ideals of norm  $N_K$ , let  $\kappa_K \in (0,1)$  be the ratio between the norm of the smallest ideal and the norm of the second-smallest. Then for any  $x \geq 2$ , it holds that

$$P_{\mathcal{K}}(x) \le \frac{a_{\mathcal{K}}}{N_{\mathcal{K}}^{x}} \cdot \left(1 + \frac{2d}{a_{\mathcal{K}}} \cdot N_{\mathcal{K}} \cdot \kappa_{\mathcal{K}}^{x-1}\right) \le \frac{3d}{N_{\mathcal{K}}^{x-1}}.$$

*Proof.* Let N be the norm of the second-smallest prime ideal of  $\mathcal{O}_{\mathcal{K}}$ , i.e.  $\kappa_{\mathcal{K}} = N_{\mathcal{K}}/N$ . For any  $n \geq N$ , let  $a_n$  be the number of prime ideals of norm N. For

any  $x \ge 1$ , let  $\pi_{\mathcal{K}}(x)$  be the number of prime ideals of norm  $\le x$ . We use Abel's summation formula and the fact that  $\pi_{\mathcal{K}}(x) \le d \cdot x$  for any  $x \ge 2$ :

$$P_{\mathcal{K}}(x) - a_{\mathcal{K}} \cdot N_{\mathcal{K}}^{-x} = \sum_{n \ge N} \frac{a_n}{n^x} = \lim_{u \to \infty} \left( \frac{\pi_{\mathcal{K}}(u)}{u^x} \right) + \int_N^{\infty} \pi_{\mathcal{K}}(u) \cdot \frac{x \cdot du}{u^{x+1}}$$
$$\leq d \cdot x \cdot \int_N^{\infty} u^{-x} du = d \cdot \frac{x}{x-1} \cdot N^{-x+1} \leq 2d \cdot N^{-x+1}.$$

Substituting  $N = N_{\mathcal{K}}/\kappa_{\mathcal{K}}$  and rearranging gives the first inequality

$$P_{\mathcal{K}}(x) \le \frac{a_{\mathcal{K}}}{N_{\mathcal{K}}^{x}} \cdot \left(1 + \frac{2d}{a_{\mathcal{K}}} \cdot N_{\mathcal{K}} \cdot \kappa_{\mathcal{K}}^{x-1}\right) .$$

The second inequality follows from  $a_{\mathcal{K}} \leq d$  and  $\kappa_{\mathcal{K}} \leq 1$ .

**Lemma 3.6.** Let  $q = p^l$  be a power of a prime and let  $K_q$  be the q-th cyclotomic field. Then  $N_{K_q} = p$  and the norm of the second-smallest ideal is at least q + 1.

Proof. First, it holds that p totally ramifies in  $\mathcal{K}_q$ , so an ideal of norm p exists in  $\mathcal{K}_q$ , now let us prove that there is no smaller ideal in this field. Let  $N \neq p$  be a prime and  $\mathfrak{p}$  a prime of  $\mathcal{K}_q$  above N. The norm of  $\mathfrak{p}$  is  $N^l$  where l is the smallest integer such that  $N^l = 1 \mod q$ , in particular, since  $N^l \neq 1$ , it holds that  $N^l \geq q+1$ .

In general, for any number field  $\mathcal{K}$  there exists  $b_{\mathcal{K}}$  such that  $P_{\mathcal{K}}(x) \sim_{x \to \infty} b_{\mathcal{K}} \cdot N_{\mathcal{K}}^{-x}$ , we provide a sagemath script computing approximation of  $b_{\mathcal{K}}$  for any cyclotomic number field. Note that  $\kappa_{\mathcal{K}}$  can be arbitrary close to 1: if  $p = 2^l - 1$  is a Mersenne prime, then one can prove that  $\kappa_{\mathcal{K}} = (p-1)/p$ .

In the ideal case (r=1), Theorem 3.1 can be compared to [SS11, Lemma 4.4], where the probability of generating  $\mathcal{O}_{\mathcal{K}}$  with two elements is proven close to  $1/\zeta_{\mathcal{K}}(2)$  for power-of-two cyclotomic number fields (note that for m=2, r=1, our upper bound is equal to  $\infty$ ). A similar proof can be used to prove that for any number fields, for suitable parameters, the probability of generating any ideal  $\mathfrak{a}$  with m elements of  $\mathfrak{a}$  is close to  $1/\zeta_{\mathcal{K}}(m)$ . One can note that when m goes to infinity, it holds that  $1/\zeta_{\mathcal{K}}(m) \sim P_{\mathcal{K}}(m) \sim 1-N_{\mathcal{K}}^{-m}$ .

# 3.5 Specialised bounds

We summarise by giving bounds for Gaussian matrices being surjective for general number fields and for the special case of prime-power cyclotomic fields.

Theorem 3.2 (General Number Field, Assuming GRH). Let K be a number field of degree  $d \geq 2$  and of smallest ideal norm  $N_K$ . Let  $0 < \varepsilon < 2^{-6}$ . Let  $\mathcal{M} \subset K^r$  be a full rank module. Then for any  $\mathbf{S} \in \mathbb{R}^{dr \times dr}$ ,  $m \geq 1$  and centre  $\mathbf{C} \in K_{\mathbb{R}}^{r \times m}$  such that

$$\mathbf{S} \ge 9\sqrt{2} \cdot r^{\frac{1}{d}} \cdot \log\left(r \cdot \Delta_K^{\frac{1}{d}} \cdot \eta_{\varepsilon}(\mathcal{M})\right)^{\frac{1}{d}} \cdot \Delta_K^{\frac{1}{d}} \cdot \eta_{\varepsilon}(\mathcal{M}) \qquad and$$

$$m \ge 2r + \frac{\log(1/\varepsilon) + 1 + \log(2d)}{\log(N_K/1.4)},$$

it holds that

$$\Pr_{\mathbf{X} \leftarrow \mathcal{D}_{\mathcal{M}, \mathbf{S}, \mathbf{C}}^{m}} (\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^{m} \neq \mathcal{M}) \leq m \cdot (\varepsilon + 2^{-dr}).$$

*Proof.* We show below that Theorem 3.1 holds for some  $B \ge 4$ . The values of the parameters and Lemma 3.5 give

$$\frac{m \cdot \varepsilon}{1 - \varepsilon} + r \left(\frac{B}{2}\right)^{-(r+1)d} + r2^{-dr} + (1 + 4\varepsilon)^{m-2r+1} \cdot r \cdot P_{\mathcal{K}}(m - 2r + 1)$$

$$\leq 1.2 \cdot m \cdot \varepsilon + r \cdot 2^{-d} \cdot 2^{-dr} + r2^{-dr} + r(1.4)^{m-2r+1} \cdot 3d \cdot N_{\mathcal{K}}^{-(m-2r)}$$

$$\leq 1.2 \cdot m \cdot \varepsilon + r \cdot (2^{-d} \cdot 2^{-dr} + 2^{-dr} + 1.4 \cdot \varepsilon)$$

$$\leq \varepsilon \cdot (1.2m + 1.4r) + 2r \cdot 2^{-dr} < m \cdot (\varepsilon + 2^{-dr}).$$

Now we argue for the existence of B. Without loss of generatity we assume  $\Sigma = \varsigma \cdot \mathbf{I}$ . The hypothesis of Theorem 3.1 are satisfied if  $B^d \geq r \cdot d^2 \ln(dr \cdot \varsigma)$  and  $\varsigma \geq B \cdot \Delta_K^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M})$ , that is to say if B satisfies

$$B^{d} \ge rd^{2} \cdot \log(dr \cdot \Delta_{\mathcal{K}} \cdot \eta_{\varepsilon}(\mathcal{M}) \cdot B) = rd \log(B^{d}) + rd^{2} \log(dr \cdot \Delta_{\mathcal{K}} \cdot \eta_{\varepsilon}(\mathcal{M})).$$

By Lemma B.5, this inequality holds for any B satisfying

$$B^d \ge 2 \left( rd^2 \log(dr \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M})) + rd \log(rd) \right).$$

It holds that

$$\left(2\left(rd^{2}\log(dr \cdot \Delta_{K}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M})\right) + rd\log(rd)\right)\right)^{1/d}$$

$$\leq \left(2r \cdot \left((d^{2} + d)\log(d) + d^{2}\log(r\Delta_{K}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M})) + d\log(r)\right)\right)^{1/d}$$

$$\leq (2r)^{1/d} \cdot \left(\left((d^{2} + d)\log(d)\right)^{1/d} + \left(d^{2}\log(r\Delta_{K}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M}))\right)^{1/d} + (d\log(r))^{1/d}\right)$$

$$\leq (2r)^{1/d} \cdot \left(2.5 + 2\log(r\Delta_{K}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M}))^{1/d} + \sqrt{2}\log(r)\right)$$

$$\leq 9 \cdot r^{1/d} \cdot \log\left(r\Delta_{K}^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M})\right)^{\frac{1}{d}},$$

hence taking  $B = 9 \cdot r^{1/d} \cdot \log(r\Delta_K^{1/d} \cdot \eta_{\varepsilon}(\mathcal{M}))^{1/d}$  gives us the desired result.  $\square$ 

# 4 Smoothing parameter of kernel lattices

In this section we upper bound  $\eta_{\varepsilon}(\Lambda^{\perp}(\mathbf{X}))$  for  $\mathbf{X} = [\mathbf{x}_1 \mid \cdots \mid \mathbf{x}_m] \in \mathcal{O}_{\mathcal{K}}^{r \times m}$ , where  $\mathbf{x}_i \leftarrow \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\varsigma}$ . We emphasise that the lattice  $\Lambda^{\perp}(\mathbf{X})$  of interest here is the rank-deficient  $\mathcal{O}_{\mathcal{K}}$ -kernel lattice of  $\mathbf{X}$  but not the full-rank q-ary lattice  $\Lambda^{\perp}_q(\mathbf{X}) := {\mathbf{x} \in \mathcal{O}_{\mathcal{K}}^m, \mathbf{X} \cdot \mathbf{x} = \mathbf{0} \bmod q}$ .

#### 4.1 Baseline: Exponential bound for m > r

We first write down an approach which gives an upper bound of  $\eta_{\varepsilon}(\Lambda^{\perp}(\mathbf{X}))$  which is exponential in r. It uses the short vectors in  $\mathbf{X}$  to construct m-r short linearly independent kernel vectors. Via the real basis of  $\mathcal{O}_{\mathcal{K}}$  we transform them into  $d \cdot (m-r)$  short linearly independent kernel vectors in  $\mathbb{R}^{dm}$ . This way we obtain a bound on  $\lambda_{d(m-r)}(\Lambda^{\perp}(\mathbf{X}))$  and the smoothing parameter is exponential in r. This can be seen as extending the approach of [LSS14] to r > 1, m > 2. The proof also relies on ideas from [BF11, Thm 4.2 (eprint)].

**Lemma 4.1.** Let  $\varsigma, t > 0$  and 0 < r < m be integers. Let  $\mathbf{X} = [\mathbf{x}_1 \mid \cdots \mid \mathbf{x}_m] \in \mathcal{O}_{\mathcal{K}}^{r \times m}$ , where  $\mathbf{x}_i \leftarrow \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \varsigma}$ . We denote  $\mathbf{X}_1 := [\mathbf{x}_1 \mid \cdots \mid \mathbf{x}_r]$ ,  $\mathbf{X}_2 := [\mathbf{x}_r \mid \cdots \mid \mathbf{x}_m]$  and assume  $\mathbf{X}_1 \in GL_r(\mathcal{K})$ . Then

$$\Pr_{\mathbf{X}} \left( \eta_{\varepsilon}(\boldsymbol{\Lambda}^{\perp}(\mathbf{X})) \leq \delta_{\mathcal{K}} \cdot \sqrt{rd} \cdot (\sqrt{r} \cdot \varsigma \cdot t)^{r} \cdot \eta_{\varepsilon}^{(d \cdot (m-r))} \right) \geq 1 - 2m \cdot rd \cdot e^{-\pi t^{2}}.$$

*Proof.* The inverse of  $\mathbf{X}_1$  can be written as  $\mathbf{X}_1^{-1} = \frac{1}{\det(\mathbf{X}_1)} \cdot \operatorname{adj}(\mathbf{X}_1)$  where  $\operatorname{adj}(\mathbf{X}_1)$  has entries in  $\mathcal{O}_{\mathcal{K}}$ . Then  $\det(\mathbf{X}_1) \cdot \mathbf{X}_1^{-1}$  is a matrix over the ring and

$$[\mathbf{X}_1 \mid \mathbf{X}_2] \cdot \begin{bmatrix} \det(\mathbf{X}_1) \cdot \mathbf{X}_1^{-1} & \mathbf{0} \\ \mathbf{0} & \mathbf{I}_{m-r} \end{bmatrix} = [\det(\mathbf{X}_1) \cdot \mathbf{I}_r \mid \mathbf{X}_2].$$

The kernel of  $[\det(\mathbf{X}_1) \cdot \mathbf{I}_r \mid \mathbf{X}_2]$  contains  $\mathcal{K}$  linearly independent vectors  $(\mathbf{x}_{r+i} \mid \det(\mathbf{X}_1) \cdot \mathbf{e}_i)$  for  $i \in [m-r]$ . Hence, vectors  $\det(\mathbf{X}_1) \cdot (\mathbf{X}_1^{-1}\mathbf{x}_{r+i} \mid \mid \mathbf{e}_i) \in \mathcal{O}_{\mathcal{K}}^m$  for  $i \in [m-r]$  are also linearly independent and in the kernel of  $\mathbf{X}$ . It remains to bound the norm of these vectors. Denote  $\mathbf{y}_i \coloneqq \det(\mathbf{X}_1) \cdot \mathbf{X}_1^{-1}\mathbf{x}_{r+i} \in \mathcal{O}_{\mathcal{K}}^r$  and let  $y_{ij}$  be the j-th coefficient of  $\mathbf{y}_i$ . Then  $\mathbf{X}_1 \cdot \mathbf{y}_i = \det(\mathbf{X}_1) \cdot \mathbf{x}_{r+i}$ . By the Cramer's rule the only solution to this linear equation over  $\mathcal{K}$  can be expressed as  $y_{ij} = \det(\tilde{\mathbf{X}}_j)/\det(\mathbf{X}_1)$ . Here  $\tilde{\mathbf{X}}_j$  stands for the matrix  $\mathbf{X}_1$  with its j-th column replaced by  $\det(\mathbf{X}_1) \cdot \mathbf{x}_{r+i}$ . Cancelling the  $\det(\mathbf{X}_1)$  factor we get  $y_{ij} = \det[\mathbf{x}_1 \mid \dots \mathbf{x}_{r+i} \dots \mid \mathbf{x}_r]$ . Hence,

$$\|\Phi(\mathbf{y}_{i})\|_{\infty} = \max_{j \in [r]} (\|\Phi(\det \left[\mathbf{x}_{1} \mid \dots \mathbf{x}_{r+i} \dots \mid \mathbf{x}_{r}\right])\|_{\infty})$$

$$= \max_{j \in [r]} \max_{k \in [d]} \det \left[\sigma_{k}(\mathbf{x}_{1}) \mid \dots \sigma_{k}(\mathbf{x}_{r+i}) \dots \mid \sigma_{k}(\mathbf{x}_{r})\right]$$

$$\leq \max_{j \in [r]} \max_{k \in [d]} \|\sigma_{k}(\mathbf{x}_{r+i})\| \cdot \prod_{\ell \in [r] \setminus \{j\}} \|\sigma_{k}(\mathbf{x}_{\ell})\|$$

$$\leq \max_{j \in [r]} \left(\sqrt{r} \cdot \|\mathbf{x}_{r+i}\|_{\infty} \cdot \prod_{\ell \in [r] \setminus \{j\}} \sqrt{r} \cdot \|\mathbf{x}_{\ell}\|_{\infty}\right)$$

$$\leq (\sqrt{r} \cdot \varsigma \cdot t)^{r}$$

with probability at least  $1 - r \cdot 2dr \cdot e^{-\pi t^2}$ . Here we apply the Hadamard's inequality for the first upper bound, the norm inequalities for the second and the tail bound from Lemma 2.6 for the last transition. Then for the Euclidean norm

we get  $\|\mathbf{y}_i\| \leq \sqrt{dr} \cdot (\sqrt{r} \cdot \varsigma \cdot t)^r$  for all  $i \in [m-r]$ . This holds with probability  $1 - m \cdot 2dr \cdot e^{-\pi t^2}$ , since the tail bound must be true for all  $\mathbf{x}_{\ell}$ ,  $\ell \in [m]$ .

As a result, we obtain m-r vectors of bounded norm that are linearly independent over  $\mathcal{K}$ . Multiplying by the ring basis  $\mathbf{B}_{\mathcal{O}_{\mathcal{K}}}$  we extend this set to  $d \cdot (m-r)$  linearly independent vectors in  $\Phi(\Lambda^{\perp}(\mathbf{X})) \subset \mathbb{R}^{dm}$  with norm bounded by  $\delta_{\mathcal{K}} \cdot \sqrt{rd} \cdot (\sqrt{r} \cdot \varsigma \cdot t)^r$ . Applying a classic bound on smoothing from §2.3, we get the statement.

# 4.2 Polynomial bound for $m \geq \Omega(\log(\epsilon^{-1})/\log(f))$ in extensions of f-th cyclotomic field

When  $r \geq \omega(1)$ , which is typically the case when the ring degree d is small, the strategy in §4.1 leads to a superpolynomial bound on  $\lambda_{d(m-r)}(\Lambda^{\perp}(\mathbf{X}))$  which is often not usable in applications. For a more useful bound we turn to the approach of [AR16] which focuses on the setting  $\mathcal{K} = \mathbb{Q}$ . We extend their strategy to any number field  $\mathcal{K} \supseteq \mathbb{Q}(\zeta_f)$  containing the f-th cyclotomic field where the quality of the obtained bound depends on f. For a high level description of our proof strategy, see §1.1. For readability, we chose to name all of the implied absolute constants of the theorems and lemmas of this section C. We highlight that those refer to possibly different constants.

**Theorem 4.1.** Let K be a number field of degree d containing the cyclotomic field of conductor f. Let  $\mathbf{S} \ge \max(\eta_{1/2}(\mathcal{O}_K^r), \sqrt{2\pi d})$ , and  $\mathbf{X} \in \mathcal{O}_K^{r \times m}$  be a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_K^r}$ ,  $\mathbf{S}$  for some  $m \ge 1$ . Let  $\varepsilon > 0$ , and  $T_{K,r,\mathbf{S}} = r \cdot \log(dr \cdot \Delta_K^{1/(2d)} \cdot s_{\min}(\mathbf{S}))$ . There exists an absolute constant C > 2 such that for

$$m \ge C \cdot T_{\mathcal{K},r,\mathbf{S}} \cdot \log(T_{\mathcal{K},r,\mathbf{S}}) + \frac{2\log(1/\varepsilon)}{\log(f/1.8)},$$

we have with probability  $\geq 1 - \varepsilon$  that

$$\eta_{\varepsilon}(\varLambda^{\perp}(\mathbf{X})) = O\left(\eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot d^{12} \cdot \delta_{\mathcal{K}}^{14} \cdot f^2 \cdot \varDelta_{\mathcal{K}}^{4/d} \cdot s_{\min}(\mathbf{S})\right).$$

Corollary 4.1 (Cyclotomics version of Theorem 4.1). Let f > 2 and  $\mathcal{K} = \mathbb{Q}(\zeta_{p^k})$  be the cyclotomic number field of conductor f and degree d. Let  $1 \leq r \leq m$ ,  $\varepsilon > 0$  and  $T = r \cdot \log(r \cdot d^{3/2} \cdot s_{\min}(\mathbf{S}_X))$ . Let  $\mathbf{S}_X \in \mathbb{R}^{dr \times dr}$  be positive definite, and  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  be a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}_X}$  and  $\mathbf{c} \in \mathcal{O}_{\mathcal{K}}^m$ . There exists an absolute constant C > 2 such that for

$$\mathbf{S}_X \ge \max(\eta_{1/2}(\mathcal{O}_{\mathcal{K}}^r), \sqrt{2\pi d}), \qquad m \ge C \cdot T \cdot \log(T) + \frac{2\log(1/\varepsilon)}{\log(p^k/1.8)},$$

with probability  $\geq 1 - \varepsilon$  we have

$$\eta_{\varepsilon}(\varLambda^{\perp}(\mathbf{X})) = O\left(\eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot f^2 \cdot d^{16} \cdot s_{\min}(\mathbf{S}_X)\right).$$

Note that the lower bound on m in Theorem 4.1 has two summands, with only the second one  $(2\log(1/\varepsilon)/\log(f/1.8))$  related to the failure probability of the theorem. This is because the theorem has a sharp threshold in m, due to the "pigeonhole" nature of Lemma 4.7. That is, we show that to prove the existence of short vector in  $\Lambda(\mathbf{X})$ , m must be larger than  $C \cdot T_{\mathcal{K},r,\mathbf{S}} \cdot \log(T_{\mathcal{K},r,\mathbf{S}})$ . Once this threshold is reached, we show in Lemma 4.9 that we can increase the width of  $\mathbf{X}$  to lower the failure probability. Indeed, we show that adding a column to  $\mathbf{X}$  multiplies the probability of failure of the theorem by 1-1/f, where  $\zeta_f \in \mathcal{K}$  is a root of unity. This leads to the  $\log(1/\varepsilon)/\log(f)$  factor in the lower bound on m.

Note that this improves on [AR16] lifted to modules in two ways. The second summand is superlinear in r and logarithmic in d in contrast to lifting [AR16] directly, where the lower bound would be superlinear in  $r \cdot d$ . The first summand is also smaller by a factor of  $\log f$ .

To prove Theorem 4.1, we state in Lemma 4.2 that for a large enough set  $A \subseteq \mathcal{O}_{\mathcal{K}}$  there exists a preimage of a small multiple of the *i*-th unit vector  $\mathbf{e}_i \in \mathcal{O}_{\mathcal{K}}^r$  for any *i*. This almost generalises Lemma 4.2 in [AR16] to the ring setting except that instead of obtaining a preimage of  $\mathbf{e}_i$  we only obtain a preimage of a small multiple of  $f \cdot \mathbf{e}_i$ . To not break the flow, we defer the proof of Lemma 4.2 to §4.3.

**Definition 4.1.** We say that a set  $A \subset \mathcal{O}_{\mathcal{K}}$  is B-admissible if it is symmetric (-A = A); it contains 0 and 1; its size is at least  $2^d$ ; all of  $x \in A$  satisfy  $\|x\|_{\infty} \leq B$ . If a set A is B-admissible, we define

$$A_{\Pi} = \{a \cdot b \mid a, b \in A\}, \quad A_{\Sigma\Pi} = \{a + b \mid a, b \in A_{\Pi}\}.$$

**Lemma 4.2.** Let  $K \supseteq \mathbb{Q}(\zeta_f)$  be a number field of degree d containing the cyclotomic field of conductor f. Let  $\mathbf{S} \ge \max(\eta_{1/2}(\mathcal{O}_K^r), \sqrt{2\pi d})$  and let  $\mathbf{X} \in \mathcal{O}_K^{r \times m}$  be a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_K^r}, \mathbf{s}$  for some  $m \ge 1$ .

Let  $\varepsilon > 0$ ,  $1 \le i \le r$ , and  $A \subset \mathcal{O}_K$  be B-admissible set for some B > 0. Let  $T = T_{K,r,B,\mathbf{S}} = r \cdot \log(r \cdot B \cdot s_{\min}(\mathbf{S})/\Delta_K^{1/(2d)})$ , then there exists an absolute constant C > 2 such that for any

$$m \ge C \cdot T \cdot \log(T) + \frac{2\log(1/\varepsilon)}{\log(f/1.8)},$$

there exists  $\mathbf{c} \in A^m_{\Sigma\Pi}$ ,  $u \in \mathcal{O}_{\mathcal{K}}$  with  $\|u\|_{\infty} \leq f^2/2$  and  $a' \in A_{\Pi}$  such that  $\mathbf{X} \cdot u \cdot \mathbf{c} = f \cdot a' \cdot \mathbf{e}_i$  with probability  $\geq 1 - \varepsilon$  over the randomness of  $\mathbf{X}$ .

Lemma 4.2 gives a probabilistic guarantee for the existence of a short preimage of  $f \cdot a \cdot \mathbf{e}_i$  for some short  $a \in A_{\varPi}$  defined relative to a set A. To obtain preimages of  $f \cdot \mathbf{e}_i$ , our strategy is to invoke Lemma 4.2 again on a different B'-admissible set A' whose elements are coprime with a. This gives a short preimage of  $f \cdot a' \cdot \mathbf{e}_i$  for a short  $a' \in A'_{\varPi}$ . We then use the co-primality of a and a' and an effective version of Bezout's identity (Lemma 2.3) to create a preimage of  $f \cdot \mathbf{e}_i$ . We discuss why

<sup>&</sup>lt;sup>8</sup> Reducing this part of the bound by a factor of d seems unlikely due to the rich structure of  $\mathcal{K}$ , i.e. the d  $\mathbb{Z}$ -vectors corresponding to an  $\mathcal{O}_{\mathcal{K}}$ -vector are highly correlated.

preimages of  $f \cdot \mathbf{I}_r$  are sufficient in the statement on the smoothing parameter in Lemma 4.5. We did not find a formula for the size of the set A' in the literature, so we provide one in Lemma 4.3, which we think might have application beyond this work.

**Lemma 4.3.** Let K be a number field of degree d with ring of integer  $\mathcal{O}_K$ . Let  $\mathfrak{a}$  be an integral ideal of  $\mathcal{O}_K$  with prime factorisation  $\mathfrak{a} = \prod_{i=1}^l \mathfrak{p}_i^{e_i}$  and  $\varepsilon = \min(2^{-d}, 3^{-l}/10)$ . Then for

$$R \ge \frac{\mathcal{N}(\mathfrak{a})^{1/d} \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \ln(1/\varepsilon^{1/d})}{2\sqrt{d}},$$

we have

$$\left|\left\{x \in \mathcal{O}_{\mathcal{K}}, \ \|x\| \le R, \ x \ coprime \ with \ \mathfrak{a}\right\}\right| \ge \frac{R^d}{4\sqrt{\Delta_{\mathcal{K}}} \cdot d^{d/2} \cdot 2^{d+l}}.$$

Proof. Let D denote the Gaussian distribution over  $\mathcal{O}_{\mathcal{K}}$  with parameter  $\varsigma \geq \mathcal{N}(a)^{1/d} \cdot \Delta_{\mathcal{K}}^{1/d} \cdot \ln(1/\varepsilon)/d$  tail-cut for  $\|x\| \geq R := 2 \cdot \sqrt{d} \cdot \varsigma$ . By [Ban93, Lemma 1.5] and [MR07, Lemmas 3.3 and 4.4] and some computations [FPS22, Lemma 2.7], it holds that for any ideal  $\mathfrak{a}$  such that  $\varsigma \geq \eta_{\varepsilon}(\mathfrak{a}), \, D(\mathfrak{a}) \in [1 - 5\varepsilon, 1 + 5\varepsilon] \cdot 1/\mathcal{N}(\mathfrak{a})$ . We write  $\mathfrak{a} = \prod_{i=1}^{l} \mathfrak{p}_{i}^{e_{i}}$ . An element  $x \in \mathcal{O}_{\mathcal{K}}$  is coprime with  $\mathfrak{a}$  if and only if it

We write  $\mathfrak{a} = \prod_{i=1}^{l} \mathfrak{p}_{i}^{e_{i}}$ . An element  $x \in \mathcal{O}_{\mathcal{K}}$  is coprime with  $\mathfrak{a}$  if and only if it does not belong to any of the  $\mathfrak{p}_{i}$ . We now compute the probability of sampling an element coprime to  $\mathfrak{a}$  from D. By the inclusion-exclusion principle and the fact that for any prime ideal  $\mathfrak{p}, \mathfrak{q}, \mathfrak{p} \cap \mathfrak{q} = \mathfrak{pq}$ , it holds that

$$D(\mathcal{O}_{\mathcal{K}} \setminus \bigcup_{i=1}^{l} \mathfrak{p}_{i}) = \sum_{k=0}^{l} (-1)^{k} \sum_{|I|=k} D(\prod_{i \in I} \mathfrak{p}_{i}) \ge \sum_{k=0}^{l} (-1)^{k} \sum_{|I|=k} \frac{1 - (-1)^{k} \cdot 5\varepsilon}{\mathcal{N}(\prod_{i \in I} \mathfrak{p}_{i})}$$

$$= \sum_{k=0}^{l} (-1)^{k} \sum_{|I|=k} \frac{1}{\mathcal{N}(\prod_{i \in I} \mathfrak{p}_{i})} - 5\varepsilon \sum_{k=0}^{l} \sum_{|I|=k} \frac{1}{\mathcal{N}(\prod_{i \in I} \mathfrak{p}_{i})}$$

$$= \prod_{i=1}^{l} \left(1 - \frac{1}{\mathcal{N}(\mathfrak{p}_{i})}\right) - 5\varepsilon \cdot \prod_{i=1}^{l} \left(1 + \frac{1}{\mathcal{N}(\mathfrak{p}_{i})}\right)$$

Setting  $\varepsilon = \min(2^{-d}, 0.1 \cdot \prod_i (\mathcal{N}(\mathfrak{p}_i) - 1)/(\mathcal{N}(\mathfrak{p}_i) + 1))$ , we get that  $D(\mathcal{O}_{\mathcal{K}} \setminus \bigcup_{i=1}^l \mathfrak{p}_i) \geq 0.5 \cdot \prod_i (1 - 1/\mathcal{N}(\mathfrak{p}_i))$ . To make this into a counting argument it suffices to say that the maximal value of D(x) is equal to  $D(0) \leq (1 + 5\varepsilon) \cdot \sqrt{\Delta_{\mathcal{K}}}/\varsigma^d \leq 2\sqrt{\Delta_{\mathcal{K}}}/\varsigma^d$ . Then the number of ring elements of  $\ell_2$  norm less than R coprime with  $\mathfrak{a}$  is greater than

$$\frac{D(\mathcal{O}_{\mathcal{K}} \setminus \bigcup_{i=1}^{l} \mathfrak{p}_i)}{D(0)} \geq \frac{\prod_i (1 - 1/\mathcal{N}(\mathfrak{p}_i)) \cdot \varsigma^d}{4\sqrt{\Delta_{\mathcal{K}}}} = \frac{\prod_i (1 - 1/\mathcal{N}(\mathfrak{p}_i)) \cdot R^d}{4\sqrt{\Delta_{\mathcal{K}}} \cdot d^{d/2} \cdot 2^d}.$$

Note that since  $\mathcal{N}(\mathfrak{p}_i) \geq 2$  for all i, then  $\prod_i (1 - 1/\mathcal{N}(\mathfrak{p}_i)) \geq 2^{-l}$  and  $\prod_i (\mathcal{N}(\mathfrak{p}_i) - 1)/(\mathcal{N}(\mathfrak{p}_i) + 1) \geq 3^{-l}$ , which concludes the proof.

**Lemma 4.4.** Let  $K \supseteq \mathbb{Q}(\zeta_f)$  be a number field of degree d containing the cyclotomic field of conductor f. Let  $\mathbf{S} \ge \max(\eta_{1/2}(\mathcal{O}_K^r), \sqrt{2\pi d})$  and  $\mathbf{X} \in \mathcal{O}_K^{r \times m}$  be a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_K^r, \mathbf{S}}$  for some  $m \ge 1$ .

Let  $\varepsilon > 0$ , and  $T_{\mathcal{K},r,\mathbf{S}} = r \cdot \log(dr \cdot \delta_K \cdot \Delta_K^{1/(2d)} \cdot s_{\min}(\mathbf{S}))$ . There exists an absolute constant C > 2 such that for

$$m \ge C \cdot T_{\mathcal{K},r,\mathbf{S}} \cdot \log(T_{\mathcal{K},r,\mathbf{S}}) + \frac{2\log(1/\varepsilon)}{\log(f/1.8)},$$

there exists a matrix  $\mathbf{U} \in \mathcal{O}_{\mathcal{K}}^{m \times r}$  with  $\|\mathbf{U}\| \leq O(\sqrt{md} \cdot f^2 \cdot \Delta_{\mathcal{K}}^{4/d} \cdot d^{10.5} \cdot \delta_{\mathcal{K}}^{13})$  and  $\mathbf{X} \cdot \mathbf{U} = f \cdot \mathbf{I}_r$ , with probability  $\geq 1 - \varepsilon$  over the randomness of  $\mathbf{X}$ .

Proof. Let  $A = \mathbf{B}_{\mathcal{O}_{\mathcal{K}}} \cdot \{\pm 1, 0\}^d \subset \mathcal{O}_{\mathcal{K}}$ , which is bounded in infinity norm by  $d \cdot \delta_{\mathcal{K}}$ , symmetric and contains  $3^d$  elements. As long as m is large enough by Lemma 4.2, there exists a short preimage to a short multiple of  $f \cdot \mathbf{e}_1$  with probability  $1 - \varepsilon/(2r)$ . Namely, there exists  $\mathbf{c} \in A^m_{\Sigma\Pi}, u \in \mathcal{O}_{\mathcal{K}}$  with  $\|u\|_{\infty} \leq f^2/2$ ,  $\|\mathbf{c}\|_{\infty} \leq 2d^2\delta_{\mathcal{K}}^2$  and  $a' \in A_{\Pi}$  such that  $\mathbf{X} \cdot u \cdot \mathbf{c} = f \cdot a' \cdot \mathbf{e}_1$ . Note that by construction  $\|a'\|_{\infty} \leq d^2\delta_{\mathcal{K}}^2$  and that a' belongs to a maximum of  $l = \log_2(\mathcal{N}(a')) \leq d\log(d^2\delta_{\mathcal{K}}^2)$  different prime ideals.

Let  $\widetilde{A}$  be the set of elements of  $\mathcal{O}_{\mathcal{K}}$  which are coprime with a' of norm bounded by  $R = 4 \cdot \Delta_{\mathcal{K}}^{1/d} \cdot d^{2.5} \cdot \delta_{\mathcal{K}}^3$ . One can check that R follows the hypothesis of Lemma 4.3, and that in that case  $\left|\widetilde{A}\right| \geq 2^d$ . The set  $\widetilde{A}$  is then R-admissible. Then by Lemma 4.2 there exists  $\widetilde{\mathbf{c}} \in \widetilde{A_{\Sigma\Pi}}^m$ ,  $\widetilde{u} \in \mathcal{O}_{\mathcal{K}}$  with  $\|\widetilde{u}\|_{\infty} \leq f^2/2$ ,  $\|\widetilde{\mathbf{c}}\|_{\infty} \leq 2R^2$  and  $\widetilde{a'} \in \widetilde{A_{\Pi}}$ ,  $\|\widetilde{a'}\|_{\infty} \leq R^2$  such that  $\mathbf{X} \cdot \widetilde{u} \, \widetilde{\mathbf{c}} = f\widetilde{a'} \cdot \mathbf{e}_1$  with probability  $1 - \varepsilon/(2r)$ .

By construction, a' and  $\widetilde{a'}$  are coprime, then by Lemma 2.3 there exist bounded norm  $\alpha, \beta \in \mathcal{O}_{\mathcal{K}}$  with  $\|\alpha\|_{\infty}, \|\beta\|_{\infty} \leq R^2 \cdot \sqrt{d} \cdot \lambda_d(\mathcal{O}_{\mathcal{K}})$  such that  $a'\alpha + \widetilde{a'}\beta = 1$ , and finally it holds that

$$\mathbf{X} \cdot (\alpha \cdot u \cdot \mathbf{c} + \beta \cdot \widetilde{u} \cdot \widetilde{\mathbf{c}}) = f \cdot \mathbf{e}_1.$$

We name  $\mathbf{u}_1 := \alpha \cdot u \cdot \mathbf{c} + \beta \cdot \widetilde{u} \cdot \widetilde{\mathbf{c}}$ . It holds that  $\|\mathbf{u}_1\|_{\infty} \leq 2\sqrt{d} \cdot f^2 \cdot \lambda_d(\mathcal{O}_{\mathcal{K}}) \cdot R^4$ . Running the same argument r times for all the  $\mathbf{e}_j$  and using the union bound, we get that with probability  $1 - \varepsilon$ , we can construct  $\mathbf{U} = [\mathbf{u}_1, \dots, \mathbf{u}_r]$  such that  $\mathbf{X} \cdot \mathbf{U} = f \cdot \mathbf{I}_r$  and  $\|\mathbf{U}\| \leq O(\sqrt{md} \cdot f^2 \cdot \Delta_{\mathcal{K}}^{4/d} \cdot d^{10.5} \cdot \delta_{\mathcal{K}}^{13})$ . The choice of lower bound on m follows from the parameters of Lemma 4.2 when we used it during the proof.

Corollary 4.2 (Cyclotomics version of Theorem 4.1). Let f > 2 and  $\mathcal{K} = \mathbb{Q}(\zeta_{p^k})$  be the cyclotomic number field of conductor f and degree d. Let  $1 \leq r \leq m$ ,  $\varepsilon > 0$  and  $T = r \cdot \log(r \cdot d^{3/2} \cdot s_{\min}(\mathbf{S}_X))$ . Let  $\mathbf{S}_X \in \mathbb{R}^{dr \times dr}$  be positive definite, and  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  be a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}_X}$  and  $\mathbf{c} \in \mathcal{O}_{\mathcal{K}}^m$ . There exists an absolute constant C > 2 such that for

$$\mathbf{S}_X \ge \max(\eta_{1/2}(\mathcal{O}_{\mathcal{K}}^r), \sqrt{2\pi d}), \qquad m \ge C \cdot T \cdot \log(T) + \frac{2\log(1/\varepsilon)}{\log(p^k/1.8)},$$

with probability  $\geq 1 - \varepsilon$  we have

$$\eta_{\varepsilon}(\Lambda^{\perp}(\mathbf{X})) = O\left(\eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot f^2 \cdot d^{16} \cdot s_{\min}(\mathbf{S}_X)\right).$$

We now use the short preimage computed in Lemma 4.4 to compute a short basis of  $\Lambda^{\perp}(\mathbf{X})$ , and then an upper bound on the smoothing parameter.

**Lemma 4.5.** Let  $\mathbf{U} \in \mathcal{O}_{\mathcal{K}}^{m \times r}$  such that  $\mathbf{X} \cdot \mathbf{U} = f \cdot \mathbf{I}_r$ , then there exists a basis  $\mathbf{B} \in \mathbb{R}^{dm \times d(m-r)}$  of  $\Phi(\Lambda^{\perp}(\mathbf{X}))$  such that

$$\|\mathbf{B}\| \le \sqrt{(m-r) d} \cdot \delta_{\mathcal{K}} \cdot (\|\mathbf{U}\| \cdot \|\mathbf{X}\| + f),$$

in particular, for any  $\varepsilon > 0$ , it holds that

$$\eta_{\varepsilon}(\Lambda^{\perp}) \leq \eta_{\varepsilon}^{((m-r)d)} \cdot \sqrt{(m-r)d} \cdot \delta_{\mathcal{K}} \cdot (\|\mathbf{U}\| \cdot \|\mathbf{X}\| + f).$$

Proof. One can check that the column vectors of matrix  $\mathbf{Y} = f \cdot \mathbf{I}_m - \mathbf{U} \cdot \mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{m \times m}$  belong to  $\Lambda^{\perp}(\mathbf{X})$  and that their norm is bounded by  $\|\mathbf{U}\| \cdot \|\mathbf{X}\| + f$ . Since  $\mathbf{U} \cdot \mathbf{X}$  is of rank at most r (over  $\mathcal{K}$ ), the rank of  $\mathbf{Y}$  is at least m - r. Hence,  $\mathbf{Y}$  contains a free subset of size m - r generating  $\operatorname{span}_{\mathcal{K}}(\Lambda^{\perp}(\mathbf{X}))$ . Then Lemma 2.1 implies the existence of a basis  $\mathbf{B}$  with the corresponding norm bound. Lastly, we get a bound on the smoothing parameter by noting that  $\lambda_{(m-r)d}(\Lambda^{\perp}(\mathbf{X})) \leq \|\mathbf{B}\|$ .

#### 4.3 Proof of Lemma 4.2

In this subsection, we fix  $\mathcal{K} \supseteq \mathbb{Q}(\zeta_f)$  be a number field of degree d containing the cyclotomic field of conductor f,  $\mathbf{S} \ge \max(\sqrt{2\pi d}, \eta_{1/2}(\mathcal{O}_{\mathcal{K}}^r))$ , and  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}}$  for some  $m \ge 1$ .

It remains to prove Lemma 4.2. Without loss of generality we take i = 1. We fix a B-admissible set A and define two set sequences, for  $1 \le j \le m$ :

$$S_j = \left\{ \sum_{i=1}^j a_i \cdot \mathbf{x}_i, \ a_i \in A \right\}, \quad \hat{S}_j = \left\{ \mathbf{s}/a, \ \mathbf{s} \in S_j, a \in A \setminus \{0\} \right\}$$

For  $1 \leq j \leq m$  we define the following events:

 $\operatorname{Win}_j$ : Two sets in  $\left\{\hat{S}_j, \hat{S}_j + \zeta_f^x \cdot \mathbf{e_1} \mid x \in [f]\right\}$  have non-empty intersection,  $\operatorname{Gain}_j : \mathbf{x}_{j+1} \notin \hat{S}_j$  and  $\|\mathbf{x}_{j+1}\| \leq \sqrt{dr} \cdot s_{\min}(\mathbf{S})$ .

The following Lemma 4.6 shows that the event  $Win_j$  implies the existence of a short preimage of a multiple of  $e_1$ .

**Lemma 4.6.** Let  $j \geq 1$ . If  $\operatorname{Win}_j$  is true, then there exists  $\mathbf{c} \in (A_{\Sigma\Pi})^j$ ,  $u \in \mathcal{O}_K$  with  $\|u\|_{\infty} \leq f^2/2$  and  $a' \in A_\Pi$  such that  $\mathbf{X}_{[1:j]} \cdot u \cdot \mathbf{c} = f \cdot a' \cdot \mathbf{e}_1$ , where  $A_\Pi$  and  $A_{\Sigma\Pi}$  are defined as in Lemma 4.2.

Proof. Since  $\operatorname{Win}_j$  is true, two cases can happen: (1)  $(\hat{S}_j + \zeta_f^x \cdot \mathbf{e_1}) \cap (\hat{S}_j + \zeta_f^y \cdot \mathbf{e_1}) \neq \emptyset$  for some  $x, y \in [f]$ ; (2)  $\hat{S}_j \cap (\hat{S}_j + \zeta_f^y \cdot \mathbf{e_1}) \neq \emptyset$  for some  $y \in [f]$ . We focus on Case (1) and note that Case (2) can be treated similarly. In Case (1), there exist  $\mathbf{a}_1, \mathbf{a}_2 \in A^j$  and  $b_1, b_2 \in A$  such that  $\mathbf{X} \cdot (\mathbf{a}_1/b_1 - \mathbf{a}_2/b_2) = (\zeta_f^x - \zeta_f^y) \cdot \mathbf{e_1}$ . Multiplying the previous equality by  $b_1 \cdot b_2 \cdot f/(\zeta_f^x - \zeta_f^y)$  and applying Lemma B.4 gives the result.

**Lemma 4.7** ([AR16, Clm. 4.1 adapted]). There exists an absolute constant  $c_1 > 1$  such that, for any sequence  $1 \le i_1 < \ldots < i_k$  of integers of length  $k \ge c_1 \cdot r \log(r \cdot B \cdot s_{\min}(\mathbf{S})/\Delta_K^{1/(2d)})$ , if  $\mathtt{Gain}_{i_l}$  is true for  $l \in [1, k+1]$  then  $\mathtt{Win}_{i_k}$  is true.

Proof. Assume that  $\operatorname{Win}_{i_k}$  is not true, we prove the statement by contradiction. Without loss of generality, we also remove all  $\mathbf{x}_j$  for which  $\operatorname{Gain}_j$  is false, and consider  $i_l = l$  for  $l \in [1, k+1]$ . For all  $l \in [1, k+1]$  and  $a \in A$ , we have that  $a \cdot \mathbf{x}_{l+1} \notin S_l$ , since else  $\mathbf{x}_{l+1} \in \hat{S}_l$  which is forbidden by  $\operatorname{Gain}_l$ . This implies that the size of  $S_l$  is at least  $|A|^l \geq 2^{dl}$ . Then, since the sets  $\hat{S}_j$  are non-intersecting, we have  $|S_l| + \sum_a \left| S_l + \zeta_f^a \cdot \mathbf{e}_1 \right| \geq (f+1)2^{dl}$ .

Furthermore, note that every element of  $S_l$  and  $S_l + \zeta_f^a \cdot \mathbf{e}_1$  has norm bounded by  $2 \cdot \sqrt{rd} \cdot B \cdot s_{\min}(\mathbf{S}) \cdot l$ . By Lemma 2.13 and the condition on  $\mathbf{S}$ , the set  $\mathcal{O}_K^r \cap B(0,C)$  is of size at most  $1.5 \cdot (2\exp(\pi)B \cdot l \cdot s_{\min}(\mathbf{S})/\Delta_K^{1/(2d)})^{rd} \leq C^{rd}$  for  $C = O(B \cdot l \cdot s_{\min}(\mathbf{S})/\Delta_K^{1/(2d)})$ . This implies that for k large enough we have  $(f+1)2^{dk} > C^{rd}$ , and a pair of sets in  $\left\{S_j, S_j + \zeta_f^a \cdot \mathbf{e}_1 \mid a \in [f]\right\}$  must have a non-empty intersection by the pigeonhole principle. This means that  $\operatorname{Win}_k$  is true, and we have a contradiction. Using Lemma B.5, it can be computed that there exists an absolute c' > 1 such that  $(f+1) \cdot 2^{dk} > C^{rd}$  holds for any  $k \geq c' \cdot r \log(r \cdot B \cdot s_{\min}(\mathbf{S})/\Delta_K^{1/(2d)})$ .

We now show that as long as  $Win_j$  is false, the event  $Gain_j$  happens with high probability.

**Lemma 4.8.** Let  $\mathbf{S} \geq \sqrt{2\pi d}$  and  $j \geq 1$ . It holds that  $\Pr(\operatorname{Gain}_j \mid \neg \operatorname{Win}_j) \geq 1 - 1.8/f$ .

*Proof.* For clarity, we omit the " $|\neg \mathtt{Win}_j$ " for the probabilities in this proof. Note that

$$\Pr(\neg \mathtt{Gain}_j) = \frac{\rho_{\mathbf{S}}(\hat{S}_j \cup \mathcal{O}_{\mathcal{K}}^r \setminus B(0, R))}{\rho_{\mathbf{S}}(\mathcal{O}_{\mathcal{K}}^r)} \leq \frac{\rho_{\mathbf{S}}(\hat{S}_j)}{\rho_{\mathbf{S}}(\mathcal{O}_{\mathcal{K}}^r)} + \beta_{dr}(\sqrt{2\pi d}),$$

where  $R = \sqrt{dr}s_1(\mathbf{S})$  and the inequality follows from Lemma 2.7. One can prove that it holds that  $\beta_{rd}(\sqrt{2\pi d}) \leq 1/(10d)$  for any  $d \geq 2$ ,  $r \geq 1$ . For any  $i \in [1, f]$ , by Lemma 2.11 and the bound on  $\mathbf{S}$ , we have that  $\rho_{\mathbf{S}}(\hat{S}_j + \zeta_f^i \mathbf{e}_1) \geq \rho_{\mathbf{S}}(\hat{S}_j)/\sqrt{e}$ . Now since  $\operatorname{Win}_j$  is false, it holds that  $\rho_{\mathbf{S}}(\hat{S}_j) + \sum_i \rho_{\mathbf{S}}(\hat{S}_j + \zeta_f^i \mathbf{e}_1) \leq \rho_{\mathbf{S}}(\mathcal{O}_K^r)$ , and then  $\rho_{\mathbf{S}}(\hat{S}_j)/\rho_{\mathbf{S}}(\mathcal{O}_K^r) \leq 1/(1+f/\sqrt{e}) \leq \sqrt{e}/f$ . Finally,  $\operatorname{Pr}(\neg \operatorname{Gain}_j) \leq 1.8/f$ 

**Lemma 4.9.** Let  $Y_1, Y_2,...$  be a sequence of not necessarily independent random Boolean variables such that for any  $i \geq 1$  and any  $\{y_1,...,y_{i-1}\} \in \{0,1\}^{i-1}$  it holds that

$$\Pr(Y_i = 1 \mid Y_{i-1} = y_{i-1}, \dots Y_1 = y_1) \ge 1 - \alpha \tag{2}$$

for some  $\alpha \in (0,2/3)$ . Let  $T \geq 3$  be a target value. Then exists an absolute constant  $c_2 > 2$  such that for any  $\varepsilon \in (0,1/2)$  and any

$$m \ge c_2 \cdot T \log(T) + 2 \cdot \frac{\log(1/\varepsilon)}{\log(1/\alpha)},$$

it holds that

$$\Pr(\sum_{i=1}^{m} Y_i \le T) \le \varepsilon.$$

Furthermore, if we assume  $T \geq 13$  and  $\alpha \leq 0.1$ , one can choose  $c_2 = 4$ .

**Lemma 4.10.** Let  $Y_1, Y_2, ...$  be a sequence of not necessarily independent random Boolean variables such that for any  $i \geq 1$  and any  $\{y_1, ..., y_{i-1}\} \in \{0, 1\}^{i-1}$  it holds that

$$\Pr(Y_i = 1 \mid Y_{i-1} = y_{i-1}, \dots Y_1 = y_1) \ge 1 - \alpha$$

for some  $\alpha \in (0,1)$ . Let  $\tau \in \{0,1\}^m$  and  $|\tau|$  its hamming weight. It holds that

$$\Pr((Y_1, \dots, Y_m) = \tau) \le \alpha^{m - |\tau|}.$$

Proof. It holds that

$$\Pr((Y_1, \dots, Y_m) = \tau) = \prod_{i=1}^m \Pr(Y_i = \tau_i | (Y_1, \dots, Y_{i-1}) = \tau_{[1:i-1]}),$$

we upper-bound each  $\Pr(Y_i = \tau_i | (Y_1, \dots, Y_{i-1}) = \tau_{[1:i-1]})$  by 1 if  $\tau_i = 1$ , and by  $\alpha$  if  $\tau_i = 0$ , giving the result.

Proof (Of Lemma 4.9). By Lemma 4.10, it holds that

$$\Pr\left(\sum_{i=1}^{m} Y_{i} \leq T\right) = \sum_{\substack{\tau \in \{0,1\}^{m} \\ |\tau| \leq T}} \Pr\left(\left(Y_{1}, \dots, Y_{m}\right) = \tau\right)$$

$$\leq \sum_{k=0}^{T} {m \choose k} \alpha^{m-k} \leq \alpha^{m-T} \cdot {m \choose T} \cdot (T+1)$$

$$\leq \alpha^{m-T} \cdot \frac{e^{T} \cdot (T+1) \cdot m^{T}}{T^{T}}$$

$$\leq 7 \cdot \alpha^{m-T} \cdot m^{T},$$

where we used the fact that  $T \leq m/2$ , by assumption. Now, note that

$$7 \cdot \alpha^{m-T} \cdot m^T \le \varepsilon \Leftrightarrow m \ge T + \frac{\log(1/\varepsilon) + \log(7)}{\log(1/\alpha)} + \frac{T}{\log(1/\alpha)} \cdot \log(m),$$

which, by Lemma B.5 holds for any

$$m \geq 2 \cdot \left(T + \frac{\log(1/\varepsilon) + \log(7)}{\log(1/\alpha)} + \frac{T}{\log(1/\alpha)} \cdot \log\left(\frac{T}{\log(1/\alpha)}\right)\right),$$

a proper analysis proves that for the range of parameters  $T \geq 13$  and  $\alpha \leq 0.1$ , the above bound is less than

$$4T\log(T) + \frac{2\log(1/\varepsilon)}{\log(1/\alpha)},$$

which concludes the proof.

With all of these result, we can now prove Lemma 4.2.

Proof (of Lemma 4.2). Combining Lemmas 4.6, 4.7 and 4.9 with  $\alpha = 1.8/f$  ( $f \ge 2$  since  $\zeta_2 = -1 \in K$ ) by Lemma 4.8 proves Lemma 4.2.

# 5 Gaussian LHL over rings

We combine the arguments above to obtain a Gaussian Leftover Hash Lemma.

Theorem 5.1 (Gaussian LHL, Assuming GRH). Let K be a number field of degree  $d \geq 2$  containing the cyclotomic field of conductor f and of smallest ideal norm  $N_{\mathcal{K}}$ . Let  $\mathbf{c} \in \mathcal{O}_{\mathcal{K}}^m$ ,  $\mathbf{S}_X \in \mathbb{R}^{dr \times dr}$  positive definite,  $1 \leq r \leq m$ ,  $0 < \varepsilon < 2^{-6}$ , and  $T = T_{\mathcal{K},r,\mathbf{S}_X} = r \cdot \log(dr \cdot \Delta_{\mathcal{K}}^{1/(2d)} \cdot s_{\min}(\mathbf{S}_X))$ . There exists an absolute constant C > 2 such that for

$$\mathbf{S}_{X} \geq \max\left(\underbrace{\sqrt{2\pi d}}_{for\ full\ image}, \underbrace{9 \cdot r^{\frac{1}{d}} \cdot \log\left(r \cdot \Delta_{K}^{\frac{1}{d}} \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}}^{r})\right)^{\frac{1}{d}} \cdot \Delta_{\mathcal{K}}^{\frac{1}{d}} \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}}^{r})}_{for\ full\ image}\right),$$

$$m \geq \max\left(\underbrace{2r + \frac{\log(1/\varepsilon) + 1 + \log(2d)}{\log(N_{\mathcal{K}}/1.4)}}_{for\ full\ image}, \underbrace{C \cdot T \cdot \log(T) + \frac{2\log(1/\varepsilon)}{\log(f/1.8)}}_{for\ the\ smoothing\ bound}\right),$$

$$\sqrt{\Sigma} \in \omega\left(\eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot d^{12} \cdot \delta_{\mathcal{K}}^{14} \cdot f^{2} \cdot \Delta_{\mathcal{K}}^{4/d} \cdot s_{\min}(\mathbf{S}_{X})\right),$$

it holds

$$\begin{split} &\mathrm{SD}((\mathbf{X} \leftarrow \left(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}_{\mathbf{X}}}\right)^m, \mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^m, \sqrt{\mathbf{\Sigma}}, \mathbf{c}}), (\mathbf{X} \leftarrow \left(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}_{\mathbf{X}}}\right)^m, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \sqrt{\mathbf{\Sigma}'}, \mathbf{X} \cdot \mathbf{c}})) \leq \varepsilon' \\ & where \ \mathbf{\Sigma}' = \tilde{\Phi}(\mathbf{X}) \cdot \mathbf{\Sigma} \cdot \tilde{\Phi}(\mathbf{X})^\mathsf{T}, \ \varepsilon' = \frac{\varepsilon}{1-\varepsilon} + m \cdot (\varepsilon + 2^{-dr}) + \varepsilon. \end{split}$$

*Proof.* Denote  $S \subset \mathcal{O}_{\mathcal{K}}^{r \times m}$  the set of matrices  $\mathbf{X}$  satisfying  $\sqrt{\overline{\Sigma}} \geq \eta_{\varepsilon}(\Lambda^{\perp}(\mathbf{X}))$  and  $\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^{m} = \mathcal{O}_{\mathcal{K}}^{r}$ . Consider a fixed element  $\mathbf{X} \in S$ . We verify conditions of Lemma 2.8 for the real lattice embedding of  $\mathcal{O}_{\mathcal{K}}^{m}$  and the  $\tilde{\varPhi}(\cdot)$  embedding of  $\mathbf{X}$ .

For matrix  $\tilde{\Phi}(\mathbf{X})$  as a linear map over  $\mathbb{R}$  the set  $\ker(\tilde{\Phi}(\mathbf{X}))$  is indeed equal to  $\operatorname{span}_{\mathbb{R}}(\Phi(\Lambda^{\perp}(\mathbf{X})))$  and  $\ker(\tilde{\Phi}(\mathbf{X})) \cap \Phi(\mathcal{O}_{\mathcal{K}}^m) = \Phi(\Lambda^{\perp}(\mathbf{X}))$ . Hence,

$$\mathrm{SD}(\mathbf{X}\cdot\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^m,\sqrt{\mathbf{\Sigma}},\mathbf{c}},\mathcal{D}_{\mathbf{X}\cdot\mathcal{O}_{\mathcal{K}}^m,\sqrt{\mathbf{\Sigma}'},\mathbf{X}\cdot\mathbf{c}})\leq\frac{\varepsilon}{1-\varepsilon}.$$

Lastly, we apply  $\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^m = \mathcal{O}_{\mathcal{K}}^r$  for the support. Now for the statistical distance between tuples  $E \coloneqq \mathrm{SD}((\mathbf{X}, \mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^m, \sqrt{\mathbf{\Sigma}}, \mathbf{c}}), (\mathbf{X}, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \sqrt{\mathbf{\Sigma}'}, \mathbf{X} \cdot \mathbf{c}}))$  we have

$$\begin{split} E &= \mathbb{E}_{\mathbf{X}}(\mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{m}, \sqrt{\mathbf{\Sigma}}, \mathbf{c}}, \mathcal{D}_{\mathbf{X} \cdot \mathcal{O}_{\mathcal{K}}^{m}, \sqrt{\mathbf{\Sigma}'}, \mathbf{X} \cdot \mathbf{c}}) \\ &\leq \Pr(\mathbf{X} \in \mathcal{S}) \cdot \max_{\mathbf{X} \in \mathcal{S}} (\operatorname{SD}(\mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{m}, \sqrt{\mathbf{\Sigma}}, \mathbf{c}}, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r}, \sqrt{\mathbf{\Sigma}'}, \mathbf{X} \cdot \mathbf{c}})) \\ &+ \Pr(\mathbf{X} \notin \mathcal{S}) \cdot \max_{\mathbf{X} \notin \mathcal{S}} (\operatorname{SD}(\mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{m}, \sqrt{\mathbf{\Sigma}}, \mathbf{c}}, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r}, \sqrt{\mathbf{\Sigma}'}, \mathbf{X} \cdot \mathbf{c}})) \\ &\leq 1 \cdot \frac{\varepsilon}{1 - \varepsilon} + (m \cdot (\varepsilon + 2^{-dr}) + \varepsilon) \cdot 1 \end{split}$$

where the first equality holds by the definition of statistical distance.

Corollary 5.1 (Prime-Power Cyclotomics version of Theorem 5.1). Let  $p^k > 2$  a prime power and  $\mathcal{K} = \mathbb{Q}(\zeta_{p^k})$  be the  $p^k$ -cyclotomic number field of degree  $d = (p-1) \cdot p^{k-1}$ . Let  $1 \le r \le m$ ,  $0 < \varepsilon < 2^{-6}$ , and  $T = r \cdot \log(r \cdot d^{3/2} \cdot s_{\min}(\mathbf{S}_X))$ . Let  $\mathbf{S}_X \in \mathbb{R}^{dr \times dr}$  be positive definite, and  $\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{r \times m}$  be a matrix with columns sampled from  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \mathbf{S}_X}$  and  $\mathbf{c} \in \mathcal{O}_{\mathcal{K}}^m$ . There exists an absolute constant C > 2 such that for

$$\mathbf{S}_X \ge 9 \cdot r^{1 + \frac{1}{d}} \cdot d^{3/2} \cdot \eta_{\varepsilon}^{(rd)} \cdot \log^{1/d} \left( (rd)^{3/2} \cdot \eta_{\varepsilon}^{(rd)} \right),$$

$$m \ge \max \left( 2r + \frac{\log(1/\varepsilon) + 1 + \log(2d)}{\log(p/1.4)}, C \cdot T \cdot \log(T) + \frac{2\log(1/\varepsilon)}{\log(p^k/1.8)} \right),$$

and for  $\sqrt{\Sigma} \in \omega \left( \eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot p^{18k} \cdot s_{\min}(\mathbf{S}_X) \right), \ \mathbf{\Sigma}' = \tilde{\Phi}(\mathbf{X}) \cdot \mathbf{\Sigma} \cdot \tilde{\Phi}(\mathbf{X})^{\mathsf{T}} \ it holds$ 

$$\mathrm{SD}((\mathbf{X}, \mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{m}, \sqrt{\mathbf{\Sigma}}, \mathbf{c}}), (\mathbf{X}, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r}, \sqrt{\mathbf{\Sigma}'}, \mathbf{X} \cdot \mathbf{c}})) \leq \frac{\varepsilon}{1 - \varepsilon} + m \cdot (\varepsilon + 2^{-dr}) + \varepsilon$$

# 6 Spherical covariance

One way to make the distribution close to a spherical Discrete Gaussian is to adapt the covariance parameter of the input vector as in the corollary below. Besides this, we set the centre of the distribution to  $\mathbf{c} = 0$  to remove the dependence on  $\mathbf{X}$  in the second distribution completely.

Corollary 6.1 (of Theorem 5.1). Let  $1 \leq r \leq m$ , and  $\varsigma > 0$ . Following Lemma 2.5 and using that  $\mathbf{M} = \tilde{\Phi}(\mathbf{X})$  is full rank we can efficiently construct a matrix  $\mathbf{S} \in \mathbb{R}^{dm \times dm}$  s.t.  $\varsigma^2 \cdot \mathbf{I}_{dr} = \tilde{\Phi}(\mathbf{X}) \cdot \mathbf{S} \cdot \mathbf{S}^{\mathsf{T}} \cdot \tilde{\Phi}(\mathbf{X})^{\mathsf{T}}$ . Letting  $\mathbf{\Sigma} = \mathbf{S} \cdot \mathbf{S}^{\mathsf{T}}$  and  $\mathbf{c} = \mathbf{0}$  we obtain

$$\mathrm{SD}((\mathbf{X}, \mathbf{X} \cdot \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^m, \sqrt{\mathbf{\Sigma}}}), (\mathbf{X}, \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \varsigma})) \leq \frac{\varepsilon}{1 - \varepsilon} + m \cdot (\varepsilon + 2^{-dr}) + \varepsilon$$

Furthermore,  $\varsigma/s_{\max}(\tilde{\Phi}(\mathbf{X})) \le s_{\min}(\sqrt{\Sigma}) \le s_{\max}(\sqrt{\Sigma}) \le \varsigma/s_{\min}(\tilde{\Phi}(\mathbf{X}))$ .

We also consider a different approach for removing the dependence on  $\mathbf{X}$  in the case when the covariance of the input cannot be adjusted. We prove that when the number of columns m is large enough the singular values of  $\mathbf{X} \cdot \mathbf{X}^\mathsf{T}$  are so close that the distribution has a constant Rényi Divergence to a spherical Discrete Gaussian. The ideas are inspired by the arguments in [AGHS13, Lemma 8]. First we detail an adaptation of the proof of [Sil85] for continuous Gaussian matrices. To adapt it we tighten the bounds on the convergence speed of the singular values.

Lemma 6.1 (Continuous Gaussians. Adapted from [Sil85]). Let  $\mathbf{Z} \leftarrow \mathcal{D}_{\varsigma}^{r \times m}$  for  $1 \leq r \leq m$ ,  $\varsigma > 0$  and parameters  $d, k \geq 1$  s.t  $rd \geq 4$ . If  $m \geq 49 \cdot \max(r, k) \cdot (rd)^2$  then with probability at least  $1 - 4r \cdot \exp(-k)$  we have:

$$m \cdot \left(1 - \frac{1}{rd}\right) \le s_{\min}^2(\mathbf{Z}) \le s_{\max}^2(\mathbf{Z}) \le m \cdot \left(1 + \frac{1}{rd}\right).$$

*Proof.* Denote the first row of  $\mathbf{Z}$  as  $\mathbf{z}_0 \leftarrow \mathcal{D}_{\varsigma}^m$ . Define  $\mathbf{O}_0 \in \mathbb{R}^{m \times m}$  a matrix where the first column is equal to  $\mathbf{z}_0/\|\mathbf{z}_0\|$  and the remaining columns are a non-random completion to an orthogonal matrix (e.g. taking the unit vectors and running the Gram-Shmidt orthogonalisation starting from  $\mathbf{z}_0/\|\mathbf{z}_0\|$ ). Then

$$\mathbf{Z} \cdot \mathbf{O}_0 = \begin{bmatrix} X_m & 0 & \dots & 0 \\ & \tilde{\mathbf{Z}}_0 & & & \end{bmatrix}$$

where  $X_m := ||\mathbf{z}_0|| \leftarrow \chi_m$ . The distribution of the rows in  $\mathbf{Z}_0$  remains independent Gaussian. This is because  $\mathbf{O}_0$  is orthogonal and the Gaussian Distribution is rotationally invariant. Then we define a rotation  $\mathbf{O}_1 \in \mathbb{R}^{r \times r}$  on the right in a similar way such that

$$\begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & & & \\ \vdots & \tilde{\mathbf{O}}_1 & & \\ 0 & & & \end{bmatrix} \cdot \mathbf{Z} \cdot \mathbf{O}_0 = \begin{bmatrix} X_m & 0 & \dots & 0 \\ Y_{r-1} & & & \\ \vdots & \tilde{\mathbf{Z}}_1 & & \\ 0 & & & & \end{bmatrix}$$

for  $Y_{r-1} \leftarrow \chi_{r-1}$ . Here the first row of  $\tilde{\mathbf{O}}_1$  is a normalisation of the first column in  $\tilde{\mathbf{Z}}_0$  and the rest is completed to an orthogonal matrix. We repeat the algorithm until

$$\prod_{j=0}^{r-2} \mathbf{O}_{2j+1} \cdot \mathbf{Z} \cdot \prod_{i=0}^{r-1} \mathbf{O}_{2i} = \begin{bmatrix} X_m & 0 & 0 & \dots & 0 \\ Y_{r-1} & X_{m-1} & 0 & \dots & 0 \\ \vdots & Y_{r-2} & & & \vdots \\ 0 & \vdots & & & \dots \end{bmatrix}$$

As consequence, there exists an orthogonal transformation  $\mathbf{O} \in \mathbb{R}^{r \times r}$  s.t.  $\mathbf{O} \cdot \mathbf{Z} \cdot \mathbf{Z}^{\mathsf{T}} \cdot \mathbf{O}^{\mathsf{T}} =$ 

$$\begin{bmatrix} X_m^2 & X_m Y_{r-1} & 0 & 0 \\ X_m Y_{r-1} & X_{m-1}^2 + Y_{r-1}^2 & X_{m-1} Y_{r-2} & 0 & 0 \\ & & \ddots & \ddots & \\ 0 & & & 0 & X_{m-r+2} Y_1 & X_{m-r+1}^2 + Y_1^2 \end{bmatrix}$$

where  $X_i \leftarrow \chi_i, Y_i \leftarrow \chi_i$ , and the variables denoted by different letter or index are pairwise independent. By the concentration inequalities of Lemma 2.10 and the union bound,  $\forall i \in [m-r+1,m+1], j \in [1,r]$  it holds  $\sqrt{i-2\sqrt{ik}} \leq X_i \leq \sqrt{i+2\sqrt{ik}+2k}$  and  $\sqrt{j-2\sqrt{jk}} \leq Y_j \leq \sqrt{j+2\sqrt{jk}+2k}$  with overall probability  $\geq 1-4r \cdot \exp(-k)$ . Consider r>1 and denote  $Y_r=Y_0=0$ . Then by the Gershgorin circle theorem the eigen values of  $\mathbf{C} \cdot \mathbf{Z} \cdot \mathbf{C}^{\mathsf{T}}$  (and hence the eigen values of  $\mathbf{Z} \cdot \mathbf{Z}^{\mathsf{T}}$ ) are bounded as

$$\lambda_{\min}(\mathbf{Z} \cdot \mathbf{Z}^{\mathsf{T}}) \geq \min_{i \in [0,r]} (X_{m-i}^2 + Y_{r-i}^2 - X_{m-i+1} Y_{r-i} - X_{m-i} Y_{r-i-1})$$

$$\lambda_{\max}(\mathbf{Z} \cdot \mathbf{Z}^{\mathsf{T}}) \le \max_{i \in [0,r]} (X_{m-i}^2 + Y_{r-i}^2 + X_{m-i+1}Y_{r-i} + X_{m-i}Y_{r-i-1})$$

Substituting in the concentration inequalities and simplifying the expressions we get that for  $m \ge c^2 \cdot \max(r, k) \cdot x^2$ 

$$\frac{\lambda_{\min}(\mathbf{Z} \cdot \mathbf{Z}^{\mathsf{T}})}{m} \ge 1 - \frac{2}{c \cdot x} - \frac{2}{c^2 \cdot x^2} - 2\sqrt{5} \cdot \frac{1}{c \cdot x} \cdot \sqrt{1 + \frac{2}{c \cdot x} + \frac{2}{c^2 \cdot x^2}}, \quad (3)$$

$$\frac{\lambda_{\max}(\mathbf{Z} \cdot \mathbf{Z}^{\mathsf{T}})}{m} \le 1 + \frac{2}{c \cdot x} + \frac{7}{c^2 \cdot x^2} + 2\sqrt{5} \cdot \frac{1}{c \cdot x} \cdot \sqrt{1 + \frac{2}{c \cdot x} + \frac{2}{c^2 \cdot x^2}}. \quad (4)$$

Then for c = 7, x = rd for  $rd \ge 4$  we have

$$\Pr(1 - \frac{1}{rd} \le \lambda_{\min}/m \le \lambda_{\max}/m \le 1 + \frac{1}{rd}) \ge 1 - 4r \cdot \exp(-k).$$

Lastly, we consider the special case when r=1. Then using  $\mathbf{z}^{\mathsf{T}} \cdot \mathbf{z} \sim \chi_m^2$ ,  $k/m \leq 1/(7rd)^2$  and Lemma 2.10 we get an even tighter bound

$$1 - \frac{1}{7rd} \le 1 - \sqrt{\frac{k}{m}} \le s_{\min}^2(\mathbf{z})/m = s_{\max}^2(\mathbf{z})/m \le 1 + \sqrt{\frac{k}{m}} + 2 \cdot \frac{k}{m} \le 1 + \frac{1}{7rd}$$

with probability  $\geq 1 - 2\exp(-k)$ .

As a side-result, we give a statement of [AGHS13, Lemma 8] with explicit constants, in the general case of ring of integers.

**Lemma 6.2 (Discrete Gaussians).** Let  $1 < r \le m$  be integers and  $k \ge 1$  s.t.  $m \ge 64 \cdot \max(r, k)$  let  $\varsigma_x \ge 10 \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}})$  for  $\varepsilon \in (0, 1/2)$ . Let  $\mathbf{X} \leftarrow \left(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \varsigma_x}\right)^m$  then it holds that

$$\frac{\sqrt{m} \cdot \varsigma_x}{7} \le s_{\min}(\tilde{\Phi}(\mathbf{X})) \le s_{\max}(\tilde{\Phi}(\mathbf{X})) \le 2.2 \cdot \sqrt{m} \cdot \varsigma_x,$$

each inequality holding with probability at least  $1 - 4rm \cdot \varepsilon - 4rd \cdot \exp(-k)$  over the randomness of  $\mathbf{X}$ .

<sup>&</sup>lt;sup>9</sup> See the included SageMath script or Appendix C.3 for the corresponding calculations.

Proof. We follow a proof similar as the one of Theorem 6.1. We work with diagonal by block matrices of dimension  $dr \times dm$  with blocks of dimension d. For matrix  $\mathbf{M}$  we denote  $\mathbf{M}_{ij}, i \in [r], j \in [m]$  the corresponding  $d \times d$  block. Let  $\varsigma_y := \varsigma_x/\sqrt{96}$ , and consider  $\mathbf{Z} = \tilde{\Phi}(\mathbf{X}) + \mathbf{Y}$  where  $\mathbf{X} \leftarrow (\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \varsigma_x})$  and  $\mathbf{Y}$  is diagonal by block with  $\mathbf{Y}_{ij} \leftarrow \operatorname{diag}\left(\mathcal{D}_{\varsigma_y}^d\right), i \in [r], j \in [m]$ . One can compute that  $\varsigma_x \cdot \varsigma_y/\sqrt{\varsigma_x^2 + \varsigma_y^2} \ge \varsigma_x/10$ , and then by the condition on  $\varsigma_x$  we can use Lemma 2.9 for  $\mathbf{Z}$ . We then have  $\operatorname{SD}(\mathbf{Z}_{ij}, \operatorname{diag}(\mathcal{D}_{\varsigma_x^2 + \varsigma_y^2}^d)) \le 4\varepsilon$ .

We analyse the singular values of **Z**. First we multiply by unitary matrices on the right and on the left to reorder the rows and columns and obtain a **block-diagonal** matrix. Now the on the diagonal we have matrices we denote  $\mathbf{Z}_i$  sampled independently from  $\mathbf{Z}_i \leftarrow \mathcal{D}_{\sqrt{\varsigma_x^2 + \varsigma_y^2}}^{r \times m}$  and similarly  $\mathbf{Y}_i \leftarrow \mathcal{D}_{\varsigma_y}^{r \times m}$  for  $i \in [d]$ . We can use Eqs. (3) and (4) of the proof of Lemma 6.1 with  $c \cdot x = 8$  (which matches the condition on m) to obtain, for any  $i \in [d]$ :

$$s_{\max}(\mathbf{Y}_i)^2 \le 2 \cdot m \cdot \varsigma_y^2,$$

$$s_{\min}(\mathbf{Z}_i)^2 \ge \frac{m}{12} \cdot (\varsigma_x^2 + \varsigma_y^2),$$

$$s_{\max}(\mathbf{Z}_i)^2 \le 2 \cdot m \cdot (\varsigma_x^2 + \varsigma_y^2)$$

with probability  $\geq 1 - 4rd \exp(-k)$ . Now, one can compute that with our choice of  $\varsigma_y$ , it holds that  $s_{\max}(\mathbf{Y}_i)/s_{\min}(\mathbf{Z}_i) \leq 1/2$ . Then using Lemma 2.4 and the fact that  $\tilde{\Phi}(\mathbf{X}) = \mathbf{Z} - \mathbf{Y}$ , we get that for any  $i \in [d]$ :

$$s_{\min}(\mathbf{X}_i) \ge \frac{1}{2} \cdot s_{\min}(\mathbf{Z}_i) \ge \frac{\sqrt{m}}{2\sqrt{12}} \cdot \sqrt{\varsigma_x^2 + \varsigma_y^2} \ge \frac{\sqrt{m} \cdot \varsigma_x}{7},$$
$$s_{\max}(\mathbf{X}_i) \le \frac{3}{2} \cdot s_{\max}(\mathbf{Z}_i) \le 2.2 \cdot \sqrt{m} \cdot \varsigma_x,$$

for any  $i \in [d]$  with probability  $\geq 1 - 4rm \cdot \varepsilon - 4rd \exp(-k)$ . The union bound give the desired result.

Now we provide an adaptation of [PS21, Lemma 2.3] on the Rényi Divergence between discrete Gaussian distributions with different covariance parameters. We provide the proof for completeness.

**Lemma 6.3.** Let  $\Lambda$  be an n-dimensional lattice  $\varepsilon \in (0,1)$  and  $\Sigma_0, \Sigma_1 \in \mathbb{R}^{n \times n}$  be positive definite s.t.  $\sqrt{\Sigma_i} \ge \eta_{\varepsilon}(\Lambda), i = 0, 1$  and  $s_{\max}(\sqrt{\Sigma_0}) \le s_{\min}(\sqrt{\Sigma_1})$ . Then

$$RD(\mathcal{D}_{\Lambda,\sqrt{\Sigma_0}}; \mathcal{D}_{\Lambda,\sqrt{\Sigma_1}}) \leq \frac{1+\varepsilon}{1-\varepsilon} \cdot \frac{\det(\sqrt{\Sigma_1})}{\det(\sqrt{\Sigma_0})}$$

*Proof.* By definition

$$RD(\mathcal{D}_{\Lambda,\sqrt{\Sigma}_{0}}; \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{1}}) = \sum_{\mathbf{s} \in \Lambda} \Pr(\mathbf{s} \leftarrow \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{0}}) \cdot \frac{\Pr(\mathbf{s} \leftarrow \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{0}})}{\Pr(\mathbf{s} \leftarrow \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{1}})}$$

$$= \mathbb{E}_{\mathbf{s} \leftarrow \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{0}}} \left( \frac{\rho_{\sqrt{\Sigma}_{0}}(\mathbf{s})}{\rho_{\sqrt{\Sigma}_{0}}(\Lambda)} \cdot \frac{\rho_{\sqrt{\Sigma}_{1}}(\Lambda)}{\rho_{\sqrt{\Sigma}_{1}}(\mathbf{s})} \right) = \frac{\rho_{\sqrt{\Sigma}_{1}}(\Lambda)}{\rho_{\sqrt{\Sigma}_{0}}(\Lambda)} \cdot \mathbb{E}_{\mathbf{s} \leftarrow \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{0}}} \left( \frac{\rho_{\sqrt{\Sigma}_{0}}(\mathbf{s})}{\rho_{\sqrt{\Sigma}_{1}}(\mathbf{s})} \right)$$

$$\leq \frac{1+\varepsilon}{1-\varepsilon} \cdot \frac{\det(\sqrt{\Sigma}_{1})}{\det(\sqrt{\Sigma}_{0})} \cdot \mathbb{E}_{\mathbf{s} \leftarrow \mathcal{D}_{\Lambda,\sqrt{\Sigma}_{0}}} \exp\left(\pi \cdot \left( \left\| \sqrt{\Sigma}_{1}^{-1} \cdot \mathbf{s} \right\|^{2} - \left\| \sqrt{\Sigma}_{0}^{-1} \cdot \mathbf{s} \right\|^{2} \right) \right)$$

where the last transition follows from Lemma 2.12. Denote  $\tilde{\mathbf{s}} := \sqrt{\Sigma_0^{-1}} \cdot \mathbf{s}$ , then

$$\begin{aligned} \left\| \sqrt{\overline{\Sigma}_{1}^{-1}} \cdot \mathbf{s} \right\|^{2} - \left\| \sqrt{\overline{\Sigma}_{0}^{-1}} \cdot \mathbf{s} \right\|^{2} &= \left\| \sqrt{\overline{\Sigma}_{1}^{-1}} \cdot \sqrt{\overline{\Sigma}_{0}} \cdot \tilde{\mathbf{s}} \right\|^{2} - \|\tilde{\mathbf{s}}\|^{2} \\ &\leq \left( \frac{s_{\max}^{2}(\sqrt{\overline{\Sigma}_{0}})}{s_{\min}^{2}(\sqrt{\overline{\Sigma}_{1}})} - 1 \right) \cdot \|\tilde{\mathbf{s}}\|^{2} \leq 0 \end{aligned}$$

since  $s_{\max}(\sqrt{\Sigma_0}) \leq s_{\min}(\sqrt{\Sigma_1})$ . Combining the above we get the statement.  $\square$ 

Finally, we prove the target statement. On high level, we first add continuous noise to the discrete Gaussian matrix  $\tilde{\Phi}(\mathbf{X})$  to make it approach a continuous Gaussian sample. Then we apply the bounds from Lemma 6.1 to the continuous matrices. Lastly, we obtain bounds on the singular values of their difference equal to  $\Phi(\mathbf{X})$  and use Lemma 6.3.

**Theorem 6.1.** Let  $1 < r \le m$  be integers and  $k \ge 1$  s.t.  $m \ge 49 \cdot \max(r, k) \cdot (rd)^2$ ,  $rd \ge 4$  let  $\varsigma_x \ge 2rd \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}})$  for  $\varepsilon \in (0, 1/2)$ . Let  $\mathbf{X} \leftarrow \left(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \varsigma_x}\right)^m$  then for  $\varsigma = \sqrt{m \cdot (\varsigma_x^2 + 2 \cdot \eta_{\varepsilon}^2(\mathcal{O}_{\mathcal{K}}))} \cdot \left(1 - \frac{1}{rd}\right)^{3/2}$  we have

$$\mathrm{RD}(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\varsigma};\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r,\sqrt{\tilde{\varPhi}(\mathbf{X})\cdot\tilde{\varPhi}(\mathbf{X})^T}})\leq 4\cdot\exp(3).$$

with probability at least  $1 - 4rm \cdot \varepsilon - 4rd \cdot \exp(-k)$  over the randomness of **X**.

Proof. In the proof we work with diagonal by block matrices of dimension  $dr \times dm$  with blocks of dimension d. For matrix  $\mathbf{M}$  we denote  $\mathbf{M}_{ij}$ ,  $i \in [r], j \in [m]$  the corresponding  $d \times d$  block. Consider  $\mathbf{Z} = \tilde{\boldsymbol{\Phi}}(\mathbf{X}) + \mathbf{Y}$  where  $\mathbf{X} \leftarrow (\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^r, \varsigma_x})$  and  $\mathbf{Y}$  is diagonal by block with  $\mathbf{Y}_{ij} \leftarrow \operatorname{diag}\left(\mathcal{D}_{\varsigma_y}^d\right)$ ,  $i \in [r], j \in [m]$ . By Lemma 2.9 for  $\mathbf{Z}$  we have  $\operatorname{SD}(\mathbf{Z}_{ij}, \operatorname{diag}(\mathcal{D}_{\varsigma_x^2 + \varsigma_y^2}^d)) \leq 4\varepsilon$  for  $\varsigma_y \coloneqq \sqrt{2} \cdot \eta_{\varepsilon}(\mathcal{O}_{\mathcal{K}})$ .

We analyse the singular values of **Z**. First we multiply by unitary matrices on the right and on the left to reorder the rows and columns and obtain a **block-diagonal** matrix. Now the on the diagonal we have matrices we denote  $\mathbf{Z}_i$  sampled independently from  $\mathbf{Z}_i \leftarrow \mathcal{D}_{\sqrt{\varsigma_x^2 + \varsigma_y^2}}^{r \times m}$  and similarly  $\mathbf{Y}_i \leftarrow \mathcal{D}_{\varsigma_y}^{r \times m}$  for

 $i \in [d]$ . By Lemma 6.1 for  $i \in [d]$ 

$$m \cdot (\varsigma_x^2 + \varsigma_y^2) \cdot \left(1 - \frac{1}{rd}\right) \le s_{\min}^2(\mathbf{Z}_i) \le s_{\max}^2(\mathbf{Z}_i) \le m \cdot (\varsigma_x^2 + \varsigma_y^2) \cdot \left(1 + \frac{1}{rd}\right)$$
$$m \cdot \varsigma_y^2 \cdot \left(1 - \frac{1}{rd}\right) \le s_{\min}^2(\mathbf{Y}_i) \le s_{\max}^2(\mathbf{Y}_i) \le m \cdot \varsigma_y^2 \cdot \left(1 + \frac{1}{rd}\right)$$

with probability at least  $1 - 8dr \cdot \exp(-k)$ . Note that  $\mathbf{X}_i = \mathbf{Z}_i - \mathbf{Y}_i$  and one can verify that  $s_{\max}(\mathbf{Y}_i) \leq \delta \cdot s_{\min}(\mathbf{Z}_i)$  for  $\varsigma_x \geq \varsigma_y \cdot \sqrt{2} \cdot rd$ ,  $\delta = 1/rd$  and  $rd \geq 3$ , see below.

$$\begin{split} m \cdot \varsigma_y^2 \cdot \left( 1 + \frac{1}{rd} \right) &\leq \delta^2 \cdot m \cdot (\varsigma_x^2 + \varsigma_y^2) \cdot \left( 1 - \frac{1}{rd} \right) \\ \frac{\varsigma_x^2}{2 \cdot (rd)^2} \cdot \left( 1 + \frac{1}{rd} \right) &\leq \delta^2 \cdot \varsigma_x^2 \cdot \left( 1 - \frac{1}{rd} \right) \\ \frac{1}{2 \cdot (rd)^2} \cdot \left( 1 + \frac{1}{rd} \right) &\leq \frac{1}{(rd)^2} \cdot \left( 1 - \frac{1}{rd} \right) \\ 1 + \frac{1}{rd} &\leq 2 - \frac{2}{rd} \\ \frac{3}{rd} &\leq 1 \end{split}$$

Then using Lemma 2.4 we get

$$m \cdot (\varsigma_x^2 + \varsigma_y^2) \cdot \left(1 - \frac{1}{rd}\right) \cdot (1 - \delta)^2 \le s_{\min}^2(\mathbf{X}_i)$$
$$m \cdot (\varsigma_x^2 + 2 \cdot \eta_\varepsilon^2(\mathcal{O}_K)) \cdot \left(1 - \frac{1}{rd}\right)^3 \le s_{\min}^2(\mathbf{X}_i)$$

and

$$s_{\max}^{2}(\mathbf{X}_{i}) \leq m \cdot (\varsigma_{x}^{2} + \varsigma_{y}^{2}) \cdot \left(1 + \frac{1}{rd}\right) \cdot (1 + \delta)^{2}$$
$$s_{\max}^{2}(\mathbf{X}_{i}) \leq m \cdot (\varsigma_{x}^{2} + 2 \cdot \eta_{\varepsilon}^{2}(\mathcal{O}_{K})) \cdot \left(1 + \frac{1}{rd}\right)^{3}$$

Then for  $\varsigma = \sqrt{m \cdot (\varsigma_x^2 + 2 \cdot \eta_\varepsilon^2(\mathcal{O}_K))} \cdot (1 - \frac{1}{rd})^{3/2}$  by Lemma 6.3

$$\operatorname{RD}(\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r},\varsigma}; \mathcal{D}_{\mathcal{O}_{\mathcal{K}}^{r},\sqrt{\tilde{\phi}(\mathbf{X})\cdot\tilde{\phi}(\mathbf{X})^{\mathsf{T}}}}) \leq \frac{1+\varepsilon}{1-\varepsilon} \cdot \frac{\left(1+\frac{1}{rd}\right)^{3rd/2}}{\left(1-\frac{1}{rd}\right)^{3rd/2}} \\
\leq \frac{1+\varepsilon}{1-\varepsilon} \cdot \frac{\exp(3/2)}{\exp(-3/2)\cdot\left(1-\frac{1}{rd}\right)} \leq 4\exp(3)$$

with probability at least  $1 - 4rm \cdot \varepsilon - 4rd \cdot \exp(-k)$  over the randomness of **X**. The last inequality used  $\exp(x) \cdot (1 - \frac{x^2}{n}) \le (1 + \frac{x}{n})^n \le \exp(x)$  for  $n \ge 1$ ,  $|x| \le n$  and also  $\varepsilon \le 1/2, rd \ge 4$ .

Remark 6.1. We expect an equivalent statement for  $negl(\lambda)$  statistical distance requires a number of columns exponential in rd. We leave this as an open problem.

```
\mathsf{HintTG}_{m,k}(\varsigma_1,\varsigma_2) \to (\mathbf{X} \in \mathcal{O}_{\mathcal{K}}^{k \times (m+k)}, \mathbf{U} \in \mathcal{O}_{\mathcal{K}}^{(m+k) \times (m+k)})
  1: \mathbf{X}_1 \leftarrow \mathcal{D}^m_{\mathcal{O}^k_{\mathcal{K}},\varsigma_1}
  2: \Sigma := \sqrt{\mathbf{S} \cdot \mathbf{S}^{\mathsf{T}}} \text{ s.t. } \varsigma_2^2 \cdot \mathbf{I}_{dr} = \tilde{\Phi}(\mathbf{X}_1) \cdot \mathbf{S} \cdot \mathbf{S}^{\mathsf{T}} \cdot \tilde{\Phi}(\mathbf{X}_1)^{\mathsf{T}} \text{ from Cor. 6.1}
   3: \forall : i \in [k]: \mathbf{r}_i \leftarrow \mathcal{D}_{\mathcal{O}_K^m, \sqrt{\Sigma}}. Let \mathbf{R} \coloneqq (\mathbf{r}_1, \dots, \mathbf{r}_k)
  4: \mathbf{X}_2 := \mathbf{X}_1 \cdot \mathbf{R} + \mathbf{I}_k \in \mathcal{O}_{\mathcal{K}}^{k \times k}

5: \mathbf{U} := \begin{bmatrix} -\mathbf{R} & -\mathbf{I}_m - \mathbf{R} \mathbf{X}_1 \\ \mathbf{I}_k & \mathbf{X}_1 \end{bmatrix} \in \mathcal{O}_{\mathcal{K}}^{(m+k) \times (m+k)}
   6: return ((\mathbf{X}_1, \mathbf{X}_2), \mathbf{U})
```

Fig. 1. HintTG procedure.

# Hardness of k-SIS and k-LWE problems

In this section we formalise a major application of our Gaussian LHL result: a reduction from the SIS problem to the k-SIS problem in the module setting. In order to be coherent with prior work's notation, in this section if K is a cyclotomic number field of degree d and conductor f, and we denote  $\mathcal{R} = \mathcal{O}_{\mathcal{K}}$  and  $\mathcal{R}_q = \mathcal{R}/(q \cdot \mathcal{R})$  for any  $q \geq 2$ . Based on our new Gaussian LHL (Theorem 5.1), we obtain the following generalisation and adaption of [LPSS14, Theorem 17 (eprint)], which is the core of our SIS-to-k-SIS reduction over modules.

**Lemma 7.1.** Let K be a number field of degree  $d \geq 2$  containing the cyclotomic field of conductor f. Let  $k = r \ge 1$ ,  $m \ge 64 \max(\lambda/\log(dk), k)$ ,  $0 < \varepsilon \le \operatorname{negl}(\lambda)$ ,  $d \cdot k \ge \lambda$  and  $\varsigma_1 \cdot \mathbf{I} = \mathbf{S}_X \in \mathbb{R}^{dk \times dk}$  follow constraints of Theorem 5.1 and satisfy  $\varsigma_1 \ge 10 \cdot \eta_{\epsilon}(\mathcal{O}_{\mathcal{K}})$ . Additionally, let  $\varsigma_2 \ge \varsigma_1 d\sqrt{mk} \cdot \omega \left(\eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot d^{12} \cdot \delta_{\mathcal{K}}^{14} \cdot f^2 \cdot \Delta_{\mathcal{K}}^{4/d} \cdot \varsigma_1\right)$ . Denote by  $\mathbf{e}_i \in \mathcal{O}_{\mathcal{K}}^k$  the i-th unit vector. Then there exists a PPT algorithm  $HintTG_{m,k}$  (Fig. 1) which, on input parameters  $(\varsigma_1, \varsigma_2)$ , outputs  $((\mathbf{X}_1, \mathbf{X}_2), \mathbf{U}) \in \mathcal{O}_{\mathcal{K}}^{k \times (m+k)} \times \mathcal{O}_{\mathcal{K}}^{(m+k) \times (m+k)}$  satisfying

- 1.  $\mathrm{SD}\left((\mathbf{X}_1,\mathbf{X}_2),\;((\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^k,\varsigma_1})^m,\prod_{i\in[k]}\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^k,\varsigma_2,\mathbf{e}_i})\right)\leq \mathsf{negl}(\lambda).$ 2. We have  $\det(\mathbf{U})=\underline{1},\; and\; letting\; \bar{\mathbf{U}}\in\mathcal{O}_{\mathcal{K}}^{(m+k)\times m}\;\; be\; the\; last\; m\; columns\; of$  $\mathbf{U}$ , then  $(\mathbf{X}_1, \mathbf{X}_2) \cdot \bar{\mathbf{U}} = \mathbf{0}_{k \times m}$ .
- 3. Lastly, it holds that  $\|\bar{\mathbf{U}}\| = \Omega(\varsigma_2 \cdot d \cdot \sqrt{m})$  with overwhelming probability in  $\lambda$ .

*Proof.* We generate required values using  $\mathsf{HintTG}_{m,k}$  described in Fig. 1. We show that the output from HintTG satisfies all Items 1 to 3 in the lemma.

For Item 1 applying Theorem 5.1 and Cor. 6.1 k times over all  $\mathbf{X}_1 \cdot \mathbf{r}_i$  together with the union bound, we have

$$\mathrm{SD}\left((\mathbf{X}_1,\mathbf{X}_1\cdot\mathbf{R}),\;((\mathcal{D}_{\mathcal{O}_K^k,\varsigma_1})^m,(\mathcal{D}_{\mathcal{O}_K^k,\varsigma_2})^k)\right)\leq \mathsf{negl}(\lambda).$$

The only difference between  $\mathbf{X}_1 \cdot \mathbf{R}$  and the output  $\mathbf{X}_2 = (\mathbf{x}_{2,1}, \dots, \mathbf{x}_{2,k})$  from HintTG is that each of the *i*-th column  $\mathbf{x}_{2,i}$  of  $\mathbf{X}_2$  is offset by  $\mathbf{e}_i \in \mathcal{O}_{\mathcal{K}}^k$ , the *i*-th unit vector, respectively. The claim follows by noting that for any  $\mathbf{y} \in \mathcal{O}_{\mathcal{K}}^k$  following  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}},\varsigma_2}^k$  and any fixed  $\mathbf{e}_i \in \mathcal{O}_{\mathcal{K}}^k$ , it holds that  $\mathbf{y} + \mathbf{e}_i$  follows  $\mathcal{D}_{\mathcal{O}_{\mathcal{K}}^k,\varsigma_2,\mathbf{e}_i}$ . For Item 2, observe that it holds

$$\mathbf{U} = egin{pmatrix} \mathbf{I}_m & -\mathbf{R} \\ \mathbf{I}_k \end{pmatrix} \cdot egin{pmatrix} -\mathbf{I}_m \\ \mathbf{I}_k & \mathbf{X}_1 \end{pmatrix}.$$

Noting that both matrices are triangular with diagonal entries  $\pm 1$  implies  $det(\mathbf{U}) = 1$ . For the kernel condition we note

$$(\mathbf{X}_1, \mathbf{X}_2) = (\mathbf{X}_1, \mathbf{X}_1 \mathbf{R} + \mathbf{I}_k) = (\mathbf{X}_1, \mathbf{I}_k) \cdot \begin{pmatrix} \mathbf{I}_m \ \mathbf{R} \\ \mathbf{I}_k \end{pmatrix},$$

thus

$$(\mathbf{X}_1, \mathbf{X}_2) \cdot \mathbf{U} = (\mathbf{X}_1, \mathbf{I}_k) \cdot \begin{pmatrix} -\mathbf{I}_m \\ \mathbf{I}_k & \mathbf{X}_1 \end{pmatrix} = (\mathbf{I}_k, \mathbf{0}_{k \times m}),$$

implying  $(\mathbf{X}_1, \mathbf{X}_2) \cdot \bar{\mathbf{U}} = \mathbf{0}_{k \times m}$ .

Finally, for Item 3, we bound the norm of  $\bar{\mathbf{U}}$ , obtaining

$$\|\bar{\mathbf{U}}\| = \left\| \begin{pmatrix} -\mathbf{I}_m - \mathbf{R} \mathbf{X}_1 \\ -\mathbf{X}_1 \end{pmatrix} \right\|$$

$$\leq (1 + \|\mathbf{R}\|^2 \cdot \|\mathbf{X}_1\|^2 + \|\mathbf{X}_1\|^2)^{1/2}$$

$$\leq 1 + (\|\mathbf{R}\| + 1) \cdot \|\mathbf{X}_1\|$$

$$\leq 1 + (\varsigma_2 \cdot \sqrt{d \cdot m} / s_{\min}(\tilde{\boldsymbol{\Phi}}(\mathbf{X}_1)) + 1) \cdot \varsigma_1 \cdot \sqrt{d \cdot m}$$

$$\leq 1 + (\varsigma_2 \cdot \sqrt{d \cdot m} \cdot 7 / (\sqrt{m} \cdot \varsigma_1) + 1) \cdot \varsigma_1 \cdot \sqrt{d \cdot m}$$

$$= \Omega(\varsigma_2 \cdot d \cdot \sqrt{m})$$

with probability overwhelming in  $d \cdot k \geq \lambda$ . In the above, the third  $\leq$  is by Gaussian tail bound together with the bound on  $s_{\max}(\sqrt{\Sigma})$  from Cor. 6.1, the fourth  $\leq$  by Lemma 6.2, and the last = because  $\varsigma_2 \cdot d$  dominates  $\varsigma_1 \cdot \sqrt{d}$ .

In the rest, we focus on the case where  $\mathcal{K}$  is a cyclotomic field. From Lemma 7.1, we have the following generalisation of [LPSS14, Lemma 19 (eprint)] without the restriction of  $\mathcal{O}_{\mathcal{K}} = \mathbb{Z}$ .

The original proof relies on [BF11, Lemma 4.5 and Theorem 4.3] which we replace by its generalisation to cyclotomic rings in Lemmas 2.14 and 2.15.

Lemma 7.2 (Generalisation of [LPSS14, Lemma 19 (eprint)]). Let  $\mathcal{K}$  be a cyclotomic field of degree  $d \geq 2$  with conductor f. Let  $k = r \geq 1$ ,  $m \geq 64 \max(\lambda/\log(dk), k)$ ,  $0 < \varepsilon \leq \operatorname{negl}(\lambda)$ , and  $\varsigma_1 \cdot \mathbf{I} = \mathbf{S}_X$  follow constraints of Theorem 5.1 and satisfy  $\varsigma_1 \geq 10 \cdot \eta_{\epsilon}(\mathcal{O}_{\mathcal{K}})$ . Additionally, let  $\varsigma_2 \geq \varsigma_1 d\sqrt{mk} \cdot \omega \left(\eta_{\varepsilon}^{((m-r)d)} \cdot m^{1.5} \cdot d^{12} \cdot \delta_{\mathcal{K}}^{14} \cdot f^2 \cdot \Delta_{\mathcal{K}}^{4/d} \cdot \varsigma_1\right)$ . Denote  $\sqrt{\Sigma} \coloneqq \operatorname{diag}(\varsigma_1 \mathbf{I}_m, \varsigma_2 \mathbf{I}_k)$ , and suppose the following holds  $^{10}$ :

<sup>&</sup>lt;sup>10</sup> The additional constraints are such that we can apply [ALLW25, Lemma 29 and Theorem 9 (eprint)] in the proof.

- -q is an unramified prime that splits into q ideals in  $\mathcal{R}$ ;
- $\begin{array}{l} q \text{ is all any sets } p \text{ that of spectrum } q \text{ is all any sets } p \text{ is all a$

Denote by  $\mathbf{e}_i \in \mathcal{R}^{(m+k)}$  the (m+i)-th unit vector. Let  $\mathsf{HintTG}_{m,k}$  be the PPT algorithm in Fig. 1. Then, the following distributions are statistically close in  $\lambda$ :

$$\left\{ (\mathbf{B}, \mathbf{X}^{\mathsf{T}}) \middle| \begin{array}{l} \mathbf{A} \leftarrow \mathcal{R}_q^{n \times m} \\ (\mathbf{X}, \mathbf{U}) \leftarrow \mathsf{HintTG}_{m, k}(\varsigma_1, \varsigma_2) \\ \bar{\mathbf{U}}^{\mathsf{T}} \coloneqq \mathbf{U} \cdot (\mathbf{0}_{(m+k) \times k}, \mathbf{I}_m) \\ \mathbf{B} \coloneqq \mathbf{A} \cdot \bar{\mathbf{U}} \bmod q \end{array} \right\} \approx_s \left\{ (\mathbf{B}, \mathbf{X}) \middle| \begin{array}{l} \mathbf{B} \leftarrow \mathcal{R}_q^{n \times (m+k)} \\ \mathbf{X} \leftarrow \prod_{i \in [k]} \mathcal{D}_{A_q^{\perp}(\mathbf{B}), \sqrt{\Sigma}, \mathbf{e}_i} \end{array} \right\}.$$

*Proof.* We consider a sequence of hybrid distributions.

 $H_0$ : The RHS distribution in the statement above.

 $\mathsf{H}_1$ : We first sample  $\mathbf{X} \leftarrow \prod_{i \in [k]} \mathcal{D}_{\mathcal{R}^{m+k}, \sqrt{\Sigma}, \mathbf{e}_i}$ , then sample  $\mathbf{B} \leftarrow \mathcal{R}_q^{n \times (m+k)}$ subject to  $\mathbf{B} \cdot \mathbf{X} = \mathbf{0} \mod q$ .

 $\mathsf{H}_2$ : Same as  $\mathsf{H}_1$ , except we sample  $\mathbf{X}$  via  $(\mathbf{X}^\mathsf{T},\mathbf{U}) \leftarrow \mathsf{HintTG}_{m,k}(\varsigma_1,\varsigma_2)$ .  $\mathsf{H}_3$ : Same as  $\mathsf{H}_2$ , except that we sample  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times m}$  and set  $\mathbf{B} \coloneqq \mathbf{A} \cdot \bar{\mathbf{U}} \bmod q$ .

It holds RHS  $\approx_s H_1$  by Lemma 2.14. Then it holds  $H_1 \approx_s H_2$  by Lemma 7.1 item 1. In the rest we show  $H_2 \approx_s H_3$ , which completes the proof since  $H_3 \equiv LHS$ .

By Lemma 2.15, the columns of **X** sampled in  $H_0$  are  $\mathcal{R}_q$ -linearly independent with overwhelming probability in  $\lambda$ . In both  $H_2$  and  $H_3$ , the sampling of  $(\mathbf{X}^T =$  $(\mathbf{X}_1,\mathbf{X}_2),\mathbf{U})$  is identical. Since we have shown RHS  $\approx_s \mathsf{H}_2$ , we have that columns of X from  $H_2$  and  $H_3$  are also  $\mathcal{R}_q$ -linearly independent with overwhelming probability in  $\lambda$ . Conditioned on the latter, we argue that the distribution of **B** in  $H_2$  and  $H_3$  are identical.

Consider an arbitrary sample  $(\mathbf{X}, \mathbf{U})$  where  $\mathbf{X}$  is  $\mathcal{R}_q$ -injective. In  $\mathsf{H}_2$ , the matrix  $\mathbf{B}$  is sampled uniformly from  $S \coloneqq \left\{ \mathbf{M} \in \mathcal{R}_q^{n \times (m+k)} : \mathbf{M} \cdot \mathbf{X} = \mathbf{0}_{n \times k} \bmod q \right\}$ . Since **X** is injective, we have  $|S| = q^{dn(m+k)}/q^{dnk} = q^{dnm}$ . Next, observe that for  $\bar{\mathbf{U}}$  in  $\mathsf{H}_3$  we have  $\bar{\mathbf{U}} \cdot \mathbf{X} = \mathbf{0}_{m \times k}$  by Lemma 7.1 Item 2. Therefore, for any  $\mathbf{A} \in \mathcal{R}_q^{n \times m}$  it holds  $\mathbf{A} \cdot \bar{\mathbf{U}} \cdot \mathbf{X} = \mathbf{0}_{m \times k} \mod q$ , implying  $\mathcal{R}_q^{n \times m} \cdot \bar{\mathbf{U}} \subseteq S$ . Further, by Lemma 7.1 Item 2,  $det(\mathbf{U}) = 1$ , thus (right-multiplication by)  $\mathbf{U}$  is a bijective map, and (right-multiplication by)  $\bar{\mathbf{U}}$  is injective. Thus,  $|\mathcal{R}_q^{n \times m} \cdot \bar{\mathbf{U}}| = |\mathcal{R}_q^{n \times m}| = q^{dnm}$ . Therefore,  $\mathcal{R}_q^{n\times m}\cdot \bar{\mathbf{U}}=S,$  and the two ways of sampling  $\mathbf{B}$  are the same. This implies  $\mathsf{H}_2\approx_s\mathsf{H}_3$  concluding the proof.

For a random matrix **A** and k Gaussian samples  $(\mathbf{x}_i)_{i \in [k]}$  from  $\Lambda_q^{\perp}(\mathbf{A})$ , the k-SIS problem asks to find a short non-zero **u** which is not in the K-span of  $(\mathbf{x}_i)_{i \in [k]}$ and s.t.  $\mathbf{A} \cdot \mathbf{u} = \mathbf{0} \mod q$ . See Definition 2.3 for a formal definition. The following theorem provides a reduction from M-SIS to k-M-SIS.

Theorem 7.1 (SIS-to-k-SIS reduction over cyclotomic rings). Let K be a cyclotomic field of degree  $d \geq 2$  with conductor f. Take the parameters  $k, n, m, q \in$ 

 $\mathbb{N}$ ,  $0 < \varepsilon \leq \text{negl}(\lambda)$ ,  $\varsigma_1 > 0$ ,  $\sqrt{\Sigma} \in \mathbb{R}^{d(m+k) \times d(m+k)}$  specified by Lemma 7.2. Further suppose that  $\varsigma_2, \varsigma_1 \geq 8d\sqrt{m} \cdot q^{n/m+2/(d \cdot m)}$ .

Let  $\beta_1 > 0$  and  $\beta_0 \geq \Omega(\varsigma_2 \cdot d \cdot \sqrt{m}) \cdot \beta_1$ . Denote  $\mathsf{params}_0 = (\mathcal{O}_{\mathcal{K}}, q, n, m, \beta_0)$ ,  $\mathsf{params}_1 = (\mathcal{O}_{\mathcal{K}}, q, n, m + k, k, \sqrt{\Sigma}, \beta_1)$ . For any PPT adversary  $\mathcal{B}$ , there is a PPT adversary  $\mathcal{A}$  such that

$$\mathsf{Adv}^{\mathsf{SIS}_{\mathsf{params}_0}}_{\mathcal{A}}(\lambda) \geq (\mathsf{Adv}^{k-\mathsf{SIS}_{\mathsf{params}_1}}_{\mathcal{B}}(\lambda))^2/O(1) - \mathsf{negl}(\lambda).$$

Remark 7.1. Firstly, here we consider the hints with simple covariance matrices, but this approach extends directly to  $\sqrt{\Sigma} = \text{diag}(\mathbf{S}_{X1}, \mathbf{S}_{X2})$ . Secondly, the constant 64 in  $m \geq 64 \max(\lambda/\log(dk), k)$  can be reduced depending on the parametrisation of the problem. For example, we can get a suitable singular value lower bound for  $\lambda = 128$ ,  $k = 4\lambda$ ,  $m = 10 \max(\lambda/\log(dk), k)$ . We leave more specific parameter settings for the applications.

*Proof.* Throughout the proof we denote by  $\mathcal{R} = \mathcal{O}_{\mathcal{K}}$  and  $\mathcal{R}_q = \mathcal{R}/q\mathcal{R}$ . We consider an intermediate search problem, which asks a PPT algorithm  $\mathcal{B}$  to solve Shift-k-SIS, defined below:

$$\begin{split} &\frac{\mathsf{Shift-}k\text{-}\mathsf{SIS}_{\mathcal{B}}(1^{\lambda})}{\mathbf{B} \leftarrow \mathcal{R}_{q}^{n \times (m+k)}; \ \mathbf{X} \leftarrow (\prod_{i \in [k]} \mathcal{D}_{\Lambda_{q}^{\perp}(\mathbf{B}), \sqrt{\mathbf{\Sigma}}, \mathbf{e}_{i}})}{\mathbf{u} \leftarrow \mathcal{B}(\mathbf{B}, \mathbf{X})} \\ &\mathbf{return} \ (\mathbf{B} \cdot \mathbf{u} = \mathbf{0} \bmod q) \wedge (\|\mathbf{u}\| \leq \beta_{1}) \wedge (\mathbf{u} \notin \mathcal{K}\text{-}\operatorname{span}(\mathbf{X})) \end{split}$$

The only difference between Shift-k-SIS and k-SIS<sub>params<sub>1</sub></sub> is the Gaussian parameters of **X**, where in Shift-k-SIS the i-th preimage is offset by the i-th unit vector  $\mathbf{e}_i$ . The theorem follows from combining Lemmas 7.3 and 7.4 below.

**Lemma 7.3.** Given the parameter constrains of Theorem 7.1, for any PPT algorithm  $\mathcal{B}$ , it holds  $\mathsf{Adv}^{\mathsf{Shift}\text{-}k\text{-SIS}}_{\mathcal{B}}(\lambda) \geq (\mathsf{Adv}^{k\text{-SIS}\mathsf{params}_1}_{\mathcal{B}}(\lambda))^2/O(1) - \mathsf{negl}(\lambda)$ .

*Proof.* We define a sequence of hybrids:

 $H_0$ : Identical to the k-SIS<sub>params<sub>1</sub></sub> experiment. We denote the problem instance as  $(\mathbf{A}, \mathbf{X}) \in \mathcal{R}_a^{n \times (m+k)} \times \mathcal{R}^{(m+k) \times k}$ .

 $\mathsf{H}_1$ : Same as  $\mathsf{H}_0$ , except that we sample  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times (m+k)}$  subject to  $\eta_{\epsilon}(\Lambda_q^{\perp}(\mathbf{A})) \le 8d\sqrt{m} \cdot q^{n/m+2/(d \cdot m)}$  for some  $\epsilon \le 2^{-2k}$ .

 $\mathsf{H}_2$ : Same as  $\mathsf{H}_1$ , except we shift the distribution  $\mathbf{X} \leftarrow \prod_{i \in [k]} \mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}),\sqrt{\Sigma}}$  to  $\mathbf{X} \leftarrow \prod_{i \in [k]} \mathcal{D}_{\Lambda_q^{\perp}(\mathbf{A}),\sqrt{\Sigma},\mathbf{e}_i}$ .

 $\mathsf{H}_3$ : Same as  $\mathsf{H}_2$ , except that we sample  $\mathbf{A} \leftarrow \mathcal{R}_q^{n \times (m+k)}$  without the constraint on  $\eta_{\epsilon}(\Lambda_q^{\perp}(\mathbf{A}))$ , i.e. we undo the change introduced in  $\mathsf{H}_1$ . Note that  $\mathsf{H}_3$  is identical to Shift-k-SIS.

Let  $\mathcal{B}$  be a PPT solver against  $\mathsf{H}_0 = k\text{-SIS}_{\mathsf{params}_1}$  with advantage  $\mathsf{Adv}_{\mathcal{B}}^{k\text{-SIS}}(\lambda) \in \mathsf{poly}(\lambda)$ . We claim:

- 1.  $\mathcal{B}$  solves  $\mathsf{H}_1$  with advantage  $\delta_1 \geq \mathsf{Adv}^{\mathsf{SIS}}_{\mathcal{B}}(\lambda) 2^{-\lambda}$ .
- 2.  $\mathcal{B}$  solves  $\mathsf{H}_2$  with advantage  $\delta_2 \geq \delta_1^2/O(1)$ .
- 3.  $\mathcal{B}$  solves  $H_3 = \text{Shift-}k\text{-SIS}$  with advantage  $\delta_3 \geq \delta_2 2^{-\lambda}$ .

Items 1 and 3 hold by Lemma 2.16, which says that the probability of  $\eta_{\epsilon}(\Lambda_q^{\perp}(\mathbf{A})) \leq 8d\sqrt{m} \cdot q^{n/m+2/(d \cdot m)}$  for any  $\epsilon \leq 2^{-2k}$  is at least  $1 - 2^{-\Omega(dm)} \geq 1 - 2^{-\lambda}$ . In the rest we argue that Item 2 holds, which concludes the proof.

We notice that the Rényi divergence between the distributions  $P \coloneqq k\text{-SIS}_{\mathsf{params}_1}$  and  $Q \coloneqq \mathsf{Shift}\text{-}k\text{-SIS}$  is identical to the one between the distributions of  $\mathbf{X}$  in  $\mathsf{H}_1$  and  $\mathsf{H}_2$ . Thus we use Lemma 2.18 to obtain  $R(P\|Q) \le 5 \exp\left(\frac{2\pi \cdot dk}{s_{\min}(\sqrt{\Sigma})^2}\right) \in O(1)$ . Note that the conditions required by Lemma 2.18 are satisfied since, for any  $\epsilon \le 2^{-2k}$ , we have  $s_{\min}(\sqrt{\Sigma}) \ge 8d\sqrt{m} \cdot q^{n/m+2/(d \cdot m)} \ge \eta_{\epsilon}(\Lambda_q^{\perp}(\mathbf{A}))$  in  $\mathsf{H}_2$ , where the first inequality is by design of  $\mathsf{H}_1$ . From here, Item 2 follows from Eq. (1).  $\square$ 

**Lemma 7.4.** Following the parameter constraints of Theorem 7.1, for any PPT adversary  $\mathcal{B}$ , there is a PPT adversary  $\mathcal{A}$  such that  $\mathsf{Adv}^{\mathsf{SIS}_{\mathsf{params}_0}}_{\mathcal{A}}(\lambda) \geq \mathsf{Adv}^{\mathsf{Shift-}k\text{-SIS}}_{\mathcal{B}}(\lambda) - \mathsf{negl}(\lambda)$ .

*Proof.* We construct the PPT algorithm  $\mathcal{A}$ . On input a  $\mathsf{SIS}_{\mathsf{params}_0}$  instance  $\mathbf{A}$ , let  $\mathcal{A}$  proceed as follows:

Sample  $(\mathbf{X}^{\mathsf{T}}, \mathbf{U}) \leftarrow \mathsf{HintTG}_{m,k}(\varsigma_1, \varsigma_2)$  using  $\mathsf{HintTG}_{m,k}$  from Fig. 1. Let  $\bar{\mathbf{U}}^{\mathsf{T}} \in \mathcal{R}^{(m+k) \times m}$  be the last m columns of  $\mathbf{U}$  and  $\mathbf{B} \coloneqq \mathbf{A} \cdot \bar{\mathbf{U}} \bmod q$ . Send  $(\mathbf{B}, \mathbf{X})$  to  $\mathcal{B}$ , and receive a vector  $\mathbf{u}$  from  $\mathcal{B}$ . Return  $\mathbf{u}^* \coloneqq \bar{\mathbf{U}} \cdot \mathbf{u}$ .

By Lemma 7.2, the distribution of  $(\mathbf{B}, \mathbf{X})$  is statistically close in  $\lambda$  to that from Shift-k-SIS. It remains to show that  $\mathbf{u}^*$  returned by  $\mathcal{A}$  is a SIS solution for  $\mathbf{A}$ . Suppose that  $\mathbf{u}$  returned by  $\mathcal{B}$  is a valid solution for Shift-k-SIS, meaning that  $\mathbf{B} \cdot \mathbf{u} = \mathbf{0} \mod q, \ 0 < \|\mathbf{u}\| \le \beta_1$ , and  $\mathbf{u} \notin \mathcal{K}$ -span $(\mathbf{X})$ . Then, we have

$$\mathbf{A} \cdot \mathbf{u}^* = \mathbf{A} \cdot \bar{\mathbf{U}} \cdot \mathbf{u} = \mathbf{B} \cdot \mathbf{u} = \mathbf{0} \bmod q.$$

Moreover,  $\bar{\mathbf{U}} \cdot \mathbf{X} = \mathbf{0}$  by Lemma 7.1 Item 2, and  $\mathbf{X}$  is  $\mathcal{R}_q$ -linearly independent with overwhelming probability by Lemma 2.15. Conditioned on the latter,  $\mathbf{X}$  is an  $\mathcal{R}_q$ -basis of the (right-)kernel of  $\bar{\mathbf{U}}$ , and thus  $\bar{\mathbf{U}} \cdot \mathbf{u} = \mathbf{0}$  would imply  $\mathbf{u} \in \mathcal{R}_q$ -span( $\mathbf{X}$ ), therefore also  $\mathbf{u} \in \mathcal{K}$ -span( $\mathbf{X}$ ), a contradiction. We thus conclude  $\mathbf{u}^* = \bar{\mathbf{U}} \cdot \mathbf{u} \neq \mathbf{0}$  with overwhelming probability. Finally, the norm of  $\mathbf{u}^*$  is bounded by

$$\|\mathbf{u}^*\| \le \|\bar{\mathbf{U}}\| \cdot \|\mathbf{u}\| \le \Omega(\varsigma_2 \cdot d \cdot \sqrt{m}) \cdot \beta_1 \le \beta_0$$

with overwhelming probability, where the second inequality follows from Lemma 7.1 Item 3. Putting everything together, we conclude that  $\mathbf{u}^*$  returned by  $\mathcal{A}$  is a SIS solution for  $\mathsf{SIS}_{\mathsf{params}_0}$  with overwhelming probability.

LWE to k-LWE reduction over cyclotomic rings. Analogous to the above, using Theorem 5.1, we obtain a reduction from the LWE problem to the k-LWE problem over general cyclotomic rings. The reduction is similar to the above, we provide a sketch and refer to [LPSS14] for the details in the special case of  $\mathcal{O}_{\mathcal{K}} = \mathbb{Z}$ .

To recall, in the k-LWE problem (formal definition given in Definition 2.4), an adversary is given as challenge instance  $(\mathbf{B}, \mathbf{X}, \mathbf{c})$ , where  $\mathbf{B} \leftarrow \mathcal{R}_q^{n \times (m+k)}$  is uniform,  $\mathbf{X} \in \mathcal{R}^{(m+k) \times k}$  are Gaussian preimages satisfying  $\mathbf{B}\mathbf{X} = \mathbf{0} \mod q$ , and it is asked to distinguish whether  $\mathbf{c} \in \mathcal{R}^{m+k}$  is an LWE sample, i.e.  $\mathbf{c}^T = \mathbf{s}^T \mathbf{B} + \mathbf{e}^T \mod q$  for some secret  $\mathbf{s}$  and error  $\mathbf{e}$ , or  $\mathbf{c} = \mathbf{d} + \mathbf{e} \mod q$  where  $\mathbf{d}$  is a random sample over  $\{\mathbf{d} : \mathbf{d}^T \mathbf{X} = \mathbf{0} \mod q\}$ . In the LWE to k-LWE reduction, given an LWE instance  $(\mathbf{A}, \hat{\mathbf{c}})$ , we simulate  $\mathbf{B}, \mathbf{X}$  using Lemma 7.2, in the same way as in the reduction in Theorem 7.1. It remains to simulate the LWE sample  $\mathbf{c}$ , for which we let  $\mathbf{c}^T \coloneqq \hat{\mathbf{c}}^T \cdot \bar{\mathbf{U}} + \bar{\mathbf{e}} \mod q$  for some freshly sampled error  $\bar{\mathbf{e}} \in \mathcal{R}^{m+k}$ . By picking the Gaussian parameters of  $\bar{\mathbf{e}}$  appropriately and appealing to standard Gaussian convolution lemmas, we can arrive at that, if  $\hat{\mathbf{c}}^T = \mathbf{s}^T \mathbf{A} + \hat{\mathbf{e}}^T \mod q$  then  $\mathbf{c}^T = (\mathbf{s}^T \mathbf{A} + \hat{\mathbf{e}}^T) \cdot \bar{\mathbf{U}} + \bar{\mathbf{e}} \approx \mathbf{s}^T \mathbf{B} + \mathbf{e} \mod q$  for  $\mathbf{e}$  of the desired shape (independent of  $\bar{\mathbf{U}}$ ), else if  $\hat{\mathbf{c}}$  is uniformly random then  $\mathbf{c}^T \approx \mathbf{d} + \mathbf{e} \mod q$ .

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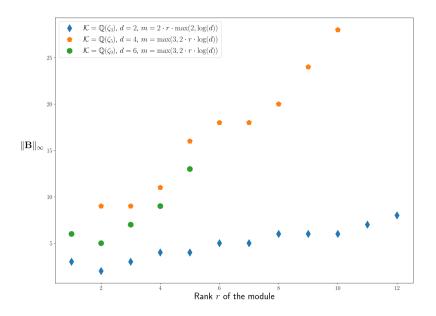
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# A Numerical experiments on Gaussian matrix's kernels



 $\bf Fig.\,2.$  Output of Appendix C.2 for constant conductor and varying rank.

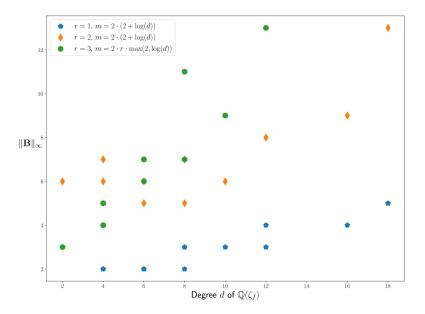


Fig. 3. Output of Appendix C.2 for constant rank and varying conductor.

In order to support our claim of the introduction that the bounds of Theorem 4.1 are very loose, we propose a sagemath program in Appendix C.2 (and as an attachment to this pdf). This program takes as input:

- -f the conductor of a cyclotomic field;
- -r a rank;
- -m a width;
- $-\beta$  a block factor;
- $-\varsigma$  a Gaussian parameter;
- num\_samples an integer.

It then do num\_samples times the following:

- 1. Sample  $\mathbf{X} \leftarrow \mathcal{D}^{r \times m}_{\mathcal{O}_{K,S}}$ ; 2. Compute  $\mathbf{X}' \in \mathbb{Z}^{rd \times md}$  the coefficient embedding of  $\mathbf{X}$  with the power-basis;
- 3. Compute  $\mathbf{B} \in \mathbb{Z}^{(m-r)d \times rd}$  a basis of  $\Lambda^{\perp}(\mathbf{X}')$ ;
- 4. Compute  $\mathbf{B}' \in \mathbb{Z}^{(m-r)d \times rd}$  the BKZ reduction of  $\mathbf{B}$  with block-size  $\beta \cdot (m-r)d \cdot m = 0$  $r) \cdot d;$
- 5. Compute the canonical embedding of each column of  ${\bf B}'$  and return the maximal infinite norm.

The algorithm then output the quantils of the num\_samples experiments (in Figs. 2 and 3 we give only the maximum). We could only test it for small values of m, r and f, but we give some plot in Figs. 2 and 3, where we take  $\beta = 0.1$ ,  $\zeta = 2\sqrt{r \cdot d}$  and num\_samples = 500. We do not claim that those figures indicate anything for values of f, m and r relevant for cryptographic purposes. We included them as an indicator that more work is required to refine the bound of Lemma 4.4.

#### B Miscellaneous results

**Lemma B.1.** For any integer  $N \ge 2$ , the number of prime ideals dividing N is at most  $d \cdot \log_2(N)$ .

*Proof.* If  $\mathfrak{p}$  is a prime ideal dividing N, then it must be above a integral prime dividing N. Since prime numbers split in at most d prime ideals, and that the number of prime integer dividing a number N is always less than  $\log_2(N)$ , the result follows.

**Lemma B.2.** Let  $\Sigma = \mathbf{S}\mathbf{S}^{\mathsf{T}} - \varsigma^2 \mathbf{I}$  with  $\mathbf{S} \in \mathbb{R}^{dr \times dr} \geq x \cdot \varsigma \mathbf{I}$  for some  $x \geq 1$ . Let  $\Sigma' = (\Sigma^{-1} + \varsigma^{-2}\mathbf{I})$ . Then  $\Sigma' \geq \varsigma^2 \cdot (1 - 1/x^2)\mathbf{I}$ .

*Proof.* By definition,  $\Sigma$  is diagonalisable with eigenvalues  $\geq (x^2 - 1) \cdot \varsigma^2$ . The eigenvalues of  $\Sigma'$  are exactly the  $(s^{-1} + \varsigma^{-2})^{-1}$  where the s are the eigenvalues of  $\Sigma$ . We then have, for every  $s \geq (x^2 - 1) \cdot \varsigma^2$ :

$$\frac{1}{\frac{1}{s} + \frac{1}{\varsigma^2}} \ge \frac{\varsigma^2}{\frac{1}{x^2 - 1} + 1} = \varsigma^2 \cdot \left(1 - \frac{1}{x^2}\right).$$

**Lemma B.3.** Let  $K = \mathbb{Q}(f)$  be the cyclotomic number field of conductor f, degree d and discriminant  $\Delta_K$ . It holds that

$$\left|\Delta_{\mathcal{K}}\right|^{1/d} \le d.$$

*Proof.* Since for any  $f = f_1 \cdot f_2$  with  $gcd(f_1, f_2) = 1$  it holds that

$$\left| \varDelta_{\mathbb{Q}(\zeta_f)} \right|^{1/\phi(f)} = \left| \varDelta_{\mathbb{Q}(\zeta_{f_1})} \right|^{1/\phi(f_1)} \cdot \left| \varDelta_{\mathbb{Q}(\zeta_{f_2})} \right|^{1/\phi(f_2)},$$

and  $\phi(f) = \phi(f_1) \cdot \phi(f_2)$ , it suffices to prove the lemma for  $f = p^k$ . In that case, it holds that  $d = p^{k-1}(p-1)$  and that  $|\Delta_{\mathcal{K}}| = p^{p^{k-1}(pk-k-1)} = p^{d \cdot (pk-k-1)/(p-1)}$ . Now we compute

$$\log_p(|\Delta_{\mathcal{K}}|^{1/d}/d) = \frac{pk - k - 1}{p - 1} - (k - 1 + \log_p(p - 1))$$
$$= \frac{p - 2}{p - 1} - \log_p(p - 1) = 1 - \frac{1}{p - 1} - \log_p(p - 1).$$

It can be computed that the function  $f(x) = 1 - 1/(x-1) - \log(x-1)/\log(x)$  is  $\leq 0$  for  $x \geq 2$ , implying the result.

**Lemma B.4 ([KLNO24, Lemma 2]).** For any  $f \geq 4$  and number field K containing a fth root of unity  $\zeta_f$ , and  $a \in \mathbb{Z}/f\mathbb{Z} \setminus \{0\}$ , it holds that  $(1 - \zeta_f^a)|f$  and  $\|f/(1 - \zeta_f^a)\|_{\infty} \leq f^2/2$ .

**Lemma B.5.** Let  $A, B \ge 0$ , then for any  $x \ge 2(A + B\log(B))$ , it holds that  $x \ge A + B \cdot \log(x)$ .

#### C Source code

The Python source code files are available as attachments to this PDF.

#### C.1 Approximate the $\zeta$ function

```
import numpy as np
from sage.all import primes, mod, exp, RR, euler_phi, ln
def find_splitting(p, n):
    Return (N, g), g being the number of prime above p in the nth cyclotomic, and N their norm https://math.stackexchange.com/a/1666456
    (Proposition 10.3 in Neukirch - Algebraic Number Theory)
    k = 0
    d = n
    while d % p == 0:
        k += 1
d /= p
    if d == 1:
         return (p, 1)
    r = mod(p, d).multiplicative_order()
    return (p**r, euler_phi(n) / (euler_phi(p**k) * r))
def find_approx_zeta_prime_cyclo(n, eps):
    Return '(a, b)' such that the Prime zeta function of the cyclotomic
    field of conductor 'n' is approximately equal to 'b·a^x'.
    :param n:
    :param eps: a precision parameter.
    EXAMPLE::
         sage: n = 2**16 + 1
         sage: a, b = find_approx_zeta_prime_cyclo(n, 1e-100)
         sage: n, euler_phi(n), a, b
(65537, 65536, 0.0000152585561077570, 1.00000011433444)
    L = [find\_splitting(p, n) for p in primes(1, n + 1)]
    L = [(N, a) \text{ for } (N, a) \text{ in } L \text{ if } N < 1 / eps]
    def approx_zeta_prime(x):
    return sum([RR(a / N**x) for (N, a) in L])
    X = list(range(2, 20)) # No particular reason for the choice of 20
    Y1 = [approx_zeta_prime(x) for x in X]
```

```
Y = [float(ln(y)) for y in Y1 if y != 0]
if len(Y) == 0:
    return (None, None)

X = X[: len(Y)]
a, b = np.polyfit(X, Y, 1)
return RR(exp(a)), RR(exp(b))
```

# C.2 Compute the basis size of the kernel of random Gaussian matrices

```
{\tt from} \ {\tt sys} \ {\tt import} \ {\tt argv}
from sage.all import vector, RDF, matrix, ZZ, CyclotomicField from sage.interfaces.r import r as r_stats
from sage.stats.distributions.discrete_gaussian_lattice import (
     {\tt DiscreteGaussianDistributionLatticeSampler\ as\ DGSL}\,,
import multiprocessing
import argparse
def nf_element_to_rd_vector(x, places):
     Compute the canonical embedding of \boldsymbol{x}.
     (separating real and imaginary part for complex embeddings)
    :param x:
     :param places:
     cc_elts = [f(x) for f in places]
     rr_elts = []
     for z in cc\_elts:
         rr_elts.append(z.real())
         rr_elts.append(z.imag())
     return vector(rr_elts)
def sample_gaussian_element(DGaussian, basis_canonical_emb, basis_roi):
     TODO describe function
     :param DGaussian:
     :param basis_canonical_emb:
     :param basis_roi:
     :returns:
    d = len(basis_roi)
     coords = vector(RDF, DGaussian() * basis_canonical_emb.inverse())
coords = vector([x.round() for x in coords])
return sum([coords[i] * basis_roi[i] for i in range(d)])
def rotation_matrix(x, OK, basis_roi):
     Return the multiplication matrix of \boldsymbol{x} in coordinate embedding.
     :param x:
     :param OK:
    :param basis_roi:
     return matrix([list(OK.coordinates(b * x)) for b in basis_roi])
```

```
def ok_matrix_to_rotation_matrix(M, OK, basis_roi):
    Return a ZZ-rotation matrix from an OK matrix.
    :param M:
    :param OK:
    :param basis_roi:
   L = \hbox{\tt [[rotation\_matrix(x, OK, basis\_roi) for x in r] for r in M.rows()]}
    return matrix.block(ZZ, L)
def infinity_norm_elt(x, places):
   Return the infinite norm in canonical embedding of \boldsymbol{x}
   :param x:
   :param places:
   return max([abs(f(x)) for f in places])
def infinity_norm_matrix(M, places):
    Return the maximum infinite norm (in canonical embedding) of all
        \hookrightarrow \text{ coefficients of } \text{M}\,.
   :param M:
   :param places:
    return max([infinity_norm_elt(x, places) for x in M.coefficients()])
def print_result(f, d, m, r, beta, out):
    Compute and print the quartiles of the output of run_batch
    :param f:
    :param d:
    :param m:
    :param r:
    :param beta:
    :param out:
    stats = dict(
       zip(
            (Q := [0, 1 / 4, 1 / 2, 3 / 4, 1]),
            r_stats.quantile(out, Q)._sage_()["DATA"],
        )
    )
    print(
        f"\{stats[0]\}, \; \{stats[1/4]\}, \; \{stats[1/2]\}, \; \{stats[3/4]\}, \; \{stats[1]\}" 
        "\n"
   )
class Worker:
```

This worker is used for parallel computing.

```
It samples a Gaussian matrix of size r*m with parameter sigma in OK and compute a beta-BKZ-reduced basis of its kernel and returns the infinity \hookrightarrow norm
of this basis.
def __init__(
    self, r, m, sigma, OK, basis_canonical_emb, basis_roi, places,
          \hookrightarrow block_size
    TODO describe function
    :param r:
     :param m:
     :param sigma:
     :param OK:
     :param basis_canonical_emb:
     :param basis_roi:
    :param places:
:param block_size:
    :returns:
    self.r = r
     self.m = m
     self.basis_canonical_emb = basis_canonical_emb
    self.basis_roi = basis_roi
self.block_size = block_size
    self.places = places
self.OK = OK
    self.sigma = sigma
def __call__(self, _):
    TODO describe function
    DGaussian = DGSL(self.basis_canonical_emb, self.sigma)
     \mbox{\# Sample} a \mbox{r*m} matrix from the Gaussian distribution
     X = matrix(
         [
                   sample_gaussian_element(
                        DGaussian, self.basis_canonical_emb, self.basis_roi
                   for _ in range(self.r)
              for _ in range(self.m)
     # Make it into a ZZ-matrix to run BKZ
     coef_X = ok_matrix_to_rotation_matrix(X, self.OK, self.basis_roi)
     basis_kernel = coef_X.kernel().matrix()
     # If there is no kernel, return 0 (this means that m should be chosen
         \hookrightarrow larger)
     if basis_kernel.nrows() == 0:
         return 0
     \# Run BKZ on the kernel basis
    small_basis_kernel = basis_kernel .BKZ(block_size=self.block_size)
return infinity_norm_matrix(small_basis_kernel, self.places)
```

def run\_batch(

```
r,
    m.
     sigma,
     beta,
    OK,
     basis_canonical_emb,
    basis_roi,
    places,
    num_samples,
    num_workers,
):
    Runs num_samples Workers in parallel using num_workers threads
    :param r:
    :param m:
     :param sigma:
    :param beta:
    :param OK:
    :param basis_canonical_emb:
    :param basis_roi:
    :param places:
     :param num_samples:
    :param num_workers:
    block_size = max(1, int((m - r) * OK.degree() * beta))
worker = Worker(r, m, sigma, OK, basis_canonical_emb, basis_roi, places,
         \hookrightarrow block_size)
     with multiprocessing.Pool(processes=num\_workers) as pool:
    results = pool.map(worker, range(1, num_samples))
return sorted(results)
def fixed_mode(f, r, m, sigma, beta, num_samples, num_workers):
    Compute stats on the size of the kernel of Gaussian matrices with fixed rank r and number field \rm QQ\,(zeta\_f) .
    :param f:
    :param r:
     :param m:
     :param sigma:
     :param beta:
     :param num_samples:
     :param num_workers:
    K, zeta = CyclotomicField(f, "zeta").objgen()
     d = K.degree()
    OK = K.ring_of_integers()
basis_roi = OK.basis()
        f"Fixed mode, parameters: f={f}, d={d}, m={m}, "
f"sigma={sigma}, beta={beta}, num_samples={num_samples}"
    m = int(eval(m))
    sigma = float(eval(sigma))
    places = K.places()
     basis_canonical_emb = matrix(
         RDF, [list(nf_element_to_rd_vector(b, places)) for b in basis_roi]
     out = run_batch(
          sigma,
```

```
beta,
         OK.
         basis_canonical_emb,
         basis_roi,
         places,
         num_samples,
         num workers.
    print_result(f, d, m, r, beta, out)
def fixed_rank_mode(r, f_min, f_max, m, sigma, beta, num_samples, num_workers
    Compute stats on the size of the kernel of Gaussian matrices with fixed
    rank r and in varying number field QQ(zeta_f).
    :param r:
    :param f_min:
    :param f_max:
    :param m:
    :param sigma:
     :param beta:
     :param num_samples:
    :param num_workers:
    print(
        f"Fixed rank mode, parameters: r={r}, f_min={f_min}, f_max={f_max}, "
         "m={m}, sigma={sigma}, beta={beta}, num_samples={num_samples}'
    for f in range(f_min, f_max):
    K, _ = CyclotomicField(f, "zeta").objgen()
    d = K.degree()
         OK = K.ring_of_integers()
basis_roi = OK.basis()
         current_m = int(eval(m))
         current_sigma = float(eval(sigma))
         places = K.places()
         basis_canonical_emb = matrix(
             RDF, [list(nf_element_to_rd_vector(b, places)) for b in basis_roi
                  \hookrightarrow ]
         )
         out = run_batch(
              current_sigma,
             beta,
             OK,
             basis_canonical_emb,
             basis_roi,
             places,
              num_samples,
             num_workers,
         print_result(f, d, current_m, r, beta, out)
def fixed_conductor_mode(f, r_min, r_max, m, sigma, beta, num_samples,
     \hookrightarrow num_workers):
    Compute stats on the size of the kernel of Gaussian matrices with varying rank and fixed number field QQ(zeta_f).
    :param f:
    :param r_min:
```

```
:param r_max:
     :param m:
     :param sigma:
     :param beta:
     :param num_samples:
     :param num_workers:
     print(
         f"Fixed conductor mode, parameters: f={f}, r_min={r_min}, f_max={
                \hookrightarrow r_max},
          f"m=\{m\}, sigma=\{sigma\}, beta=\{beta\}, num\_samples=\{num\_samples\}"
    )
    K, _ = CyclotomicField(f, "zeta").objgen()
d = K.degree()
     OK = K.ring_of_integers()
basis_roi = OK.basis()
places = K.places()
     basis_canonical_emb = matrix(
          RDF, [list(nf_element_to_rd_vector(b, places)) for b in basis_roi]
     for r in range(r_min, r_max):
          current_m = int(eval(m))
          current_sigma = float(eval(sigma))
          out = run_batch(
               current_m,
               current_sigma,
               beta,
               OK,
               basis_canonical_emb,
               basis_roi,
               places,
               num_samples,
               num_workers,
          print_result(f, d, current_m, r, beta, out)
def main():
     global number_trial
     parser = argparse.ArgumentParser(
          formatter_class=argparse.RawDescriptionHelpFormatter,
          description="""Calculate statistics on the size in infinity norm of
                   the
canonical embedding of a beta-reduced basis of the kernel of some Gaussian matrix on the ROI of a cyclotomic number field. In rank mode, the number field is fixed and r_min <=r < r_max. In conductor mode, the rank is fixed and f_min <=f < f_max. In fixed mode, r and f are fixed.
Format of the output: conductor degree rank width (ZZ-dimension of the kernel
BKZ
""",
BKZ block-size) min, 1st quartile, median, 3rd quartile, maximum
     parser.add_argument(
          "-f", type=int, help="The conductor of the number field (rank or \hookrightarrow fixed mode)"
     {\tt parser.add\_argument} \; (
          "-r",
type=int,
          help="The rank of the module to consider (conductor or fixed mode)",
```

```
parser.add_argument(
      '-m".
     type=str,
     help="Formula for m in function of r, f and d (in python syntax).
          \hookrightarrow Default 2*r*log(d)",
    default="2*r*log(d)",
parser.add_argument(
      --sigma",
     type=str,
    help="Formula for sigma in function of r, f and d (in python syntax). 
 \hookrightarrow Default 2*sqrt(rd)", 
 default="2*sqrt(r*d)",
parser.add_argument(
     "--f_min",
     type=int,
    help="The min conductor of the number field (in conductor mode)",
parser.add_argument(
     "--f_max",
     type=int,
    \verb|help="The max conductor of the number field (in conductor mode)"|,
parser.add_argument(
     "--r_min", type=int, help="The min rank of the module (in rank mode)"
)
parser.add_argument(
      --r_max", type=int, help="The max rank of the module (in rank mode)"
parser.add_argument(
     type=float
    help="The block-size coefficient for BKZ (default 0.5). The final \hookrightarrow block-size will be beta*m*r",
    default=0.5,
{\tt parser.add\_argument} \, (
     "--samples",
     tvpe=int.
     default=100,
    help="Number of sample to use (default: 100)",
parser.add_argument(
      '--cores", type=int, default=0, help="Number of cores to use (default \hookrightarrow : all)"
if len(argv) == 1:
    parser.print_help()
     exit(1)
args = parser.parse_args()
num_workers = args.cores if args.cores > 0 else None
number_trial = args.samples
beta = args.beta
num_samples = args.samples
m = args.m
sigma = args.sigma
if args.r is not None and args.f is not None:
    f = args.f
```

```
r = args.r
       fixed_mode(f, r, m, sigma, beta, num_samples, num_workers)
    elif args.r is not None:
       r = args.r
       f_min = args.f_min
       f_max = args.f_max
       fixed_rank_mode(r, f_min, f_max, m, sigma, beta, num_samples,
           elif args.f is not None:
       f = args.f
       r_min = args.r_min
       r_max = args.r_max
       fixed_conductor_mode(f, r_min, r_max, m, sigma, beta, num_samples,

→ num workers)

if __name__ == "__main__":
    main()
```

#### C.3 Parametrising singular value inequalities

```
from sage.symbolic.relation import solve
     This script simplifies the tail bounds on Chi random variables defined in \hookrightarrow Section 6, Lemma 6.1 of the paper.
      We first write down the expressions obtain from the Gershgorin circle
             \hookrightarrow theorem that dominate other values in the maximum (resp. minimum).
            \hookrightarrow\, They are stated as a comment. Then we manually write down the
            \hookrightarrow corresponding tail bounds defined in variables f1 to f7.
      Next we simplify the tail bounds only making them larger (resp. smaller).
            \hookrightarrow Lastly, we plot the difference between our values and the desired \hookrightarrow upper (resp. lower) bound 1 + 1/x (resp. 1 - 1/x).
     We check that the inequalities work by plotting them. The larger we make \hookrightarrow m the closer the values are to the bound, hence we parametrise it \hookrightarrow with c and give the user the option to plot for different values \hookrightarrow of c. In the paper we use c = 7 and x \geq 4 since it satisfies
            \hookrightarrow all inequalities.
      This script has a range of different plots we draw via show(plot(...)).
            \hookrightarrow To change the plot displayed please comment out the current plot \hookrightarrow and uncomment the one required. In all plots the bound is
            \hookrightarrow satisfied whenever the line is above zero.
var("m,k,r,c,x")
m = c^2 * x^2 * k
# We only need the inequality plots for positive parameters.
assume(c>1)
assume(x>1)
# THE UPPER TAIL
# The random variables we analyse:
 * x_{(m-1)^2} + y_{(r-1)^2} + x_{(m)} * y_{(r-1)} + x_{(m-1)} * y_{(r-2)} 
# Their norm divided by m is smaller than:
f1 = (m - 1)/m + 2 * sqrt((m-1)*k/m^2) + 2*k/m + (r - 1)/m + 2 * sqrt((r-1)*k \hookrightarrow /m^2) + 2*k/m + sqrt(1 + 2* sqrt(k/m) + 2*k/m) * sqrt((r - 1)/m + 2 *
```

```
 \hookrightarrow \  \, \operatorname{sqrt}((r-1)*k/m^2) \ + \  \, 2*k/m) \ + \  \, \operatorname{sqrt}((m-1)/m \ + \  \, 2 \ * \  \, \operatorname{sqrt}((m-1)*k/m^2) \ + \  \, 2* \\ \hookrightarrow \  \, k/m) \ * \  \, \operatorname{sqrt}((r-2)/m \ + \  \, 2 \ * \  \, \operatorname{sqrt}((r-2)*k/m^2) \ + \  \, 2*k/m) 
 # We simplify the expression making the bound looser
f1 = f1.subs({(m-1):m})
f1 = f1.subs({(k-1):k})
f1 = f1.subs({(k-2):k})
f1 = f1.canonicalize_radical()
# Next random variables we analyse:
\# x_{m}(m)^{2} / m + x_{m}(m) * y_{m}(r-1) / m <=
f2 = 1 + 2 * sqrt(k/m) + 2*k/m + sqrt(1 + 2* sqrt(k/m) + 2*k/m) * sqrt((r - 2* sqrt(k/m) + 2* sqrt(k/m) + 2*k/m) * sqrt((r - 2* sqrt(k/m) + 2* sqrt(k/m) + 2* sqrt(k/m) * sqrt(k
             \hookrightarrow 1)/m + 2 * sqrt((r-1)*k/m^2) + 2*k/m)
# We simplify the expression
f2 = f2.subs({(k-1):k})
f2 = f2.canonicalize_radical()
 # The plot of f1 and f2 compared with 1 + 1/x.
show(plot(((1 + 1/x) - f2).subs(c=7), 1, 10) + plot(((1 + 1/x) - f1).subs(c \hookrightarrow =7), 1, 10, color = "green"))
# THE LOWER TAIL
           Here 0 \cdot Z \cdot Z^\transpose \cdot 0^\transpose has different form
                       \hookrightarrow depending on r so we consider cases r >= 4, r = 3 and r = 2. \hookrightarrow Otherwise the lower tail is analysed in the same way.
# r >= 4
# Random variable we analyse:
\# x_{m}(m)^{2} / m - x_{m}(m) * y_{m}(r-1) / m >=
f3 = 1 - 2 * sqrt(k/m) - sqrt(1 + 2 * sqrt(k/m) + 2 * k/m) * sqrt((r - 1)/m + \hookrightarrow 2 * sqrt((r-1)*k/m^2) + 2 * k/m)
# We simplify the expression
f3 = f3.subs({(k-1):k})
f3 = f3.canonicalize_radical()
# Random variable we analyse:
# x_(m - r + 2)^2 / m + y_(2)^2 / m - x_(m - r + 2) * y_(1) / m - x_(m - r + \hookrightarrow 3) * y_(2) / m >=
# we expect this one to be the smallest
\hookrightarrow sqrt(2/m + 2 * sqrt(2*k/m^2) + 2 * k/m)
 # We simplify the expression
f4 = f4.subs({(k-4):k})
f4 = f4.subs({(m - r + 4):m})
f4 = f4.subs({(m - r + 2):m})
f4 = f4.subs({(m - r + 3):m})
f4 = f4.subs(k = 4)
f4 = f4.canonicalize_radical()
```

```
# Random variable we analyse:
f5 = 1 - (r - 2)/m - 2 * sqrt((m - r + 2)*k/m^2) - sqrt((m - r + 2)/m + 2 * \hookrightarrow sqrt((m - r + 2)*k/m^2) + 2 * k/m) * sqrt(1/m + 2 * sqrt(k/m^2) + 2 *
# We simplify the expression
f5 = f5.subs({(k-2):k})
f5 = f5.subs({(m - r + 2):m})
f5 = f5.subs(k = 4)
f5 = f5.canonicalize_radical()
\# c = 6 works for x > 4
# Uncomment to show the plot one at a time.
# special case for r = 3 (since we have m - r + 4) f6 = f4 but the bounds \hookrightarrow after are different
f6 = (m - r + 4)/m - 2 * sqrt((m - r + 4)*k/m^2) - sqrt((m - r + 2)/m + 2 *  \hookrightarrow \text{sqrt}((m - r + 2)*k/m^2) + 2 * k/m) * \text{sqrt}(1/m + 2 * \text{sqrt}(k/m^2) + 2 * \\ \hookrightarrow k/m) - \text{sqrt}((m - r + 3)/m + 2 * \text{sqrt}((m - r + 3)*k/m^2) + 2 * k/m) * \text{sqrt} \\ \hookrightarrow (2/m + 2 * \text{sqrt}(2*k/m^2) + 2 * k/m)
# We simplify the expression
f6 = f6.subs({(m - r + 4):(m+1)})
f6 = f6.subs({(m - r + 2):m})
f6 = f6.subs(\{(m - r + 3):m\})
f6 = f6.subs(k = 3)
f6 = f6.canonicalize_radical()
\# since k = 3 now the combined effect makes c = 6 require x > 8 which is a
              \hookrightarrow bit worse than for r > 4
\# Uncomment to show the plot one at a time.
#show(plot((f3 - (1 - 1/x)).subs(c=7), 1, 10) + plot((f6 - (1 - 1/x)).subs(c \hookrightarrow =7), 1, 10, color = "green") + plot((f5 - (1 - 1/x)).subs(c=7), 1, 10, \hookrightarrow color = "red"))
# for r = 2 the matrix looks simpler and we only get x_{m} = 2 / m - x_{m} = 2 / m - x_{m} = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m = 2 / m =
              \hookrightarrow (r-1)) / m >=
f7 = 1 - 2 * sqrt(k/m) - sqrt( 1 + 2 * sqrt(k/m) + 2 * k/m) * sqrt((r - 1)/m \hookrightarrow + 2 * sqrt((r-1)*k/m^2) + 2 * k/m)
# We simplify the expression
f7 = f7.subs({(k-1):k})
f7 = f7.canonicalize_radical()
\# Uncomment to show the plot one at a time.
\#show(plot((f7 - (1 - 1/x)).subs(c=5), 1, 10))
```