

## TUTORIAL 9

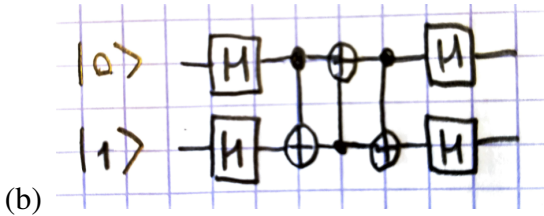
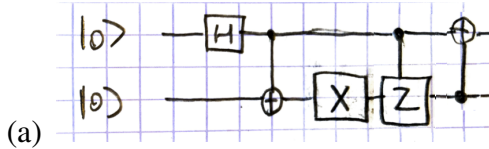
### 1 Some Small Calculations

1. Give the probability of each outcome when you measure  $|\varphi\rangle$  in basis  $\mathcal{B}$ :

(a)  $|\varphi\rangle = \frac{|0\rangle+i|1\rangle}{\sqrt{2}}$  and  $\mathcal{B} = \left( \frac{|0\rangle+|1\rangle}{\sqrt{2}}, \frac{|0\rangle-|1\rangle}{\sqrt{2}} \right)$ .

(b)  $|\varphi\rangle = \frac{|00\rangle+(1-i)|10\rangle-|11\rangle}{2}$  and  $\mathcal{B} = \left( |00\rangle, |01\rangle, \frac{|10\rangle+|11\rangle}{\sqrt{2}}, \frac{|10\rangle-|11\rangle}{\sqrt{2}} \right)$ .

2. Compute the probability of measuring 1 on the first qubit of the following circuits:



3. Using a circuit computing  $U_f^\oplus : |x\rangle|a\rangle|0\rangle \mapsto |x\rangle|a \oplus f(x)\rangle|0\rangle$ , build a circuit that computes  $Z_f : |x\rangle|0\rangle \mapsto (-1)^{f(x)}|x\rangle|0\rangle$ .

### 2 Simon’s Problem Generalized

Consider a function  $f : \{0, 1\}^n \rightarrow \{0, 1\}^n$  with the promise that there exists a vector subspace  $V$  of  $\{0, 1\}^n$  (seen as a vector space over  $\mathbb{F}_2$ ) such that:

$$\forall x, y \in \{0, 1\}^n, f(y) = f(x) \Leftrightarrow \exists v \in V, x = y + v .$$

1. Recall Simon’s algorithm circuit and compute its final state.
2. With  $V = \text{span}(b_1, \dots, b_c)$ , prove the following equality:

$$\sum_{v \in V} (-1)^{v \cdot y} = \prod_{i=1}^c (1 + (-1)^{b_i \cdot y}) .$$

3. Show that one run of Simon’s algorithm outputs  $x \in \{0, 1\}^n$  such that  $x$  is orthogonal to  $V$  ( $\forall y \in V, x \cdot y = 0 \pmod 2$ ).

*Hint: Measure first the  $n$  first qubits: it won’t change the result of the measurement on the  $n$  last qubits, but the analysis will be easier.*

*Remark.* The usual version of Simon’s Problem is when  $V = \{0, a\}$  for some  $a \in \{0, 1\}^n$ . □

### 3 Superdense Coding

In this exercise, we will have two actors, Alice and Bob. Alice has two classical bits of information  $(b_0, b_1) \in \{0, 1\}^2$  and she want to send them to Bob.

- (informal) What is the minimum number of classical bits that Alice has to send to Bob in order to communicate him  $(b_0, b_1)$ ?
- Now suppose that Alice and Bob share an entangled **EPR pair** (or **Bell pair**), that is to say there is a quantum state  $|\phi\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$  such that the first qubit is owned by Alice (she can only perform operations of the form  $U \otimes I$  on  $|\phi\rangle$ ) and the second qubit is owned by Bob.  
 Alice is going to perform operations on her qubit and send it to Bob. If  $b_1 = 1$ , she applies  $X$  (bit flip) and then if  $b_0 = 1$  she applies  $Z$  (phase flip), she send her part of the qubit back to Bob.  
 Explain how Bob can recover the values of  $b_0$  and  $b_1$  using  $|\phi\rangle$ .

