

2. Shannon Channel Coding Theorem:

2.1 1. $D(y^m) = \underset{(\epsilon)}{\operatorname{argmax}} \underbrace{W^m(y^m | E(o))}_{\text{Most likely o that outputs } y^m}$

2. $P_{err,o} = 1 - \sum_{\substack{y^m \in Y^m \\ D(y^m) = o}} W^m(y^m | E(o))$

$= \sum_{y^m \in Y^m} W^m(y^m | E(o)) \mathbb{1}_{D(y^m) \neq o}$

3. $\mathbb{1}_{D(y^m) \neq o} \leq \mathbb{1}_{\underbrace{W^m(y^m | E(o)) < \alpha(m, y^m)}_{\text{either } o \text{ is not in } 1^{\text{st}} \text{ case}} \vee \underbrace{(\exists o' \neq o W^m(y^m | E(o')) \geq \alpha(m, y^m))}_{o \text{ is not the only } o' \text{ in } 1^{\text{st}} \text{ case}}}$

$\leq \mathbb{1}_{W^m(y^m | E(o)) < \alpha(m, y^m)} + \mathbb{1}_{\exists o' \neq o W^m(y^m | E(o')) \geq \alpha(m, y^m)}$

$\leq \mathbb{1}_{W^m(y^m | E(o)) < \alpha(m, y^m)} + \sum_{o' \neq o} \mathbb{1}_{W^m(y^m | E(o')) \geq \alpha(m, y^m)}$

Thus taking the sum over $y^m \in Y^m$, one gets $P_{err,o} \leq P_{err,o}^1 + P_{err,o}^2$

2.2 1. If $\mathbb{E}_E [P_{err}] \xrightarrow{m \rightarrow \infty} 0$

It means that for any ϵ , there exists E^m such that $P_{err}^{E^m} \leq \mathbb{E}_E [P_{err}]$

So in particular $P_{err}^{E^m} \xrightarrow{m \rightarrow \infty} 0$ **QED**

BR: $\mathbb{P}(P_{err}^{E^m} \leq \frac{2 \mathbb{E}_E [P_{err}]}{\epsilon}) \geq \frac{1}{2}$ by Markov inequality, which gives a good enough scheme E^m .

$$\begin{aligned}
 2 \text{ @ } \mathbb{E}_E [P_{m,0}^1] &= \sum_{x^n} P_{x^n}(x^n) P_{m,0, \epsilon}^1(x^n) \quad \text{by definition of random } E(\cdot) \\
 &= \sum_{x^n y^m} P_{x^n}(x^n) \underbrace{W(y^m|x^n)}_{= P_{x^n y^m}(x^n, y^m)} \mathbb{1}_{W(y^m|x^n) \leq K(m, \epsilon)} P_{y^m}(y^m) \\
 &= P_{x^n y^m}(x^n, y^m) \quad \text{by definition}
 \end{aligned}$$

But

So with $X^m Y^m = (X_1 Y_1, X_2 Y_2, \dots, X_m Y_m)$ with $X_i, Y_i \sim P_{X,Y}$
 we have that $\mathbb{E}_E [P_{m,0}^1] = \mathbb{P} \left(\prod_{i=1}^m W(Y_i | X_i) < K(m, \epsilon) \prod_{i=1}^m P_Y(Y_i) \right)$

$$\begin{aligned}
 \text{b) } &\mathbb{P} \left(\prod_{i=1}^m W(Y_i | X_i) < K(m, \epsilon) \prod_{i=1}^m P_Y(Y_i) \right) \\
 &= \mathbb{P} \left(\sum_{i=1}^m \log(W(Y_i | X_i)) < \log K(m, \epsilon) + \sum_{i=1}^m \log P_Y(Y_i) \right) \\
 &= \mathbb{P} \left(\sum_{i=1}^m \underbrace{\log \left(\frac{P_{X,Y}(X_i, Y_i)}{P_X(X_i) P_Y(Y_i)} \right)}_{i_{X,Y}(X_i, Y_i)} < \log K(m, \epsilon) \right)
 \end{aligned}$$

$$\text{c) } \mathbb{E} [i_{X,Y}(X_i, Y_i)] = \sum_{x,y} P_{X,Y}(x,y) i_{X,Y}(x,y) = I(X:Y) = C(W)$$

since P_X has been chosen to make $I(X:Y)$ achieve $C(W)$

$$\text{d) } \lim_{m \rightarrow +\infty} \mathbb{P} \left(\left| \frac{1}{m} \sum_{i=1}^m i_{X,Y}(X_i, Y_i) - C(W) \right| < \epsilon \right) = 1$$

$$\begin{aligned}
 \text{So } \mathbb{E}_E [P_{m,0}^1] &= \mathbb{P} \left(\frac{1}{m} \sum_{i=1}^m i_{X,Y}(X_i, Y_i) - C(W) < \frac{1}{m} \log K(m, \epsilon) - C(W) \right) \\
 &\xrightarrow{m \rightarrow +\infty} 0 \quad \text{if } \frac{1}{m} \log K(m, \epsilon) - C(W) \leq -\epsilon \\
 &\text{ie. } \boxed{K(m, \epsilon) \leq 2^m (C(W) - \epsilon)} \quad (\text{for some } \epsilon > 0)
 \end{aligned}$$

$$3. \mathbb{E}_E [P_{m, \epsilon}] = \sum_{y^m \neq 0} \mathbb{E}_E \left[\sum_{x^m} W^m(y^m/x^m) \right] W^m(y^m/x^m)$$

$$\leq 2^{R_m} \sum_{x^m, y^m} P_{x^m}(x^m) P_{x^m}(x^m) W^m(y^m/x^m) \cdot \underbrace{W^m(y^m/x^m)}_{\text{does not depend on } x^m}$$

one term

no. of ops

$$\leq 2^{R_m} \sum_{x^m, y^m} P_{y^m}(y^m) P_{x^m}(x^m) \cdot \underbrace{W^m(y^m/x^m)}_{\text{are also equal by}}$$

$$\leq 2^{R_m} \sum_{x^m, y^m} \frac{W(y^m/x^m)}{K(m, \epsilon)} P_{x^m}(x^m) = \frac{2^{R_m}}{K(m, \epsilon)}$$

So it is sufficient that $\boxed{\frac{2^{R_m}}{K(m, \epsilon)} \rightarrow 0 \text{ as } m \rightarrow +\infty}$ to conclude

4. $R < C(W)$, take $\epsilon = \frac{C(W) - R}{2}$ and $K(m, \epsilon) = 2^{(R + \epsilon)m}$

Then $K(m, \epsilon)$ fulfills both conditions

$$\begin{aligned} \text{Thus } \mathbb{E}_E [P_m] &= \frac{1}{2^{\epsilon m}} \sum_0 \mathbb{E} [P_{m, \epsilon}] \\ &\leq \frac{1}{2^{\epsilon m}} \sum_0 \underbrace{\mathbb{E} [P_{m, \epsilon}^1]}_{\rightarrow 0 \text{ as } m \rightarrow +\infty} + \frac{1}{2^{\epsilon m}} \sum_0 \underbrace{\mathbb{E} [P_{m, \epsilon}^2]}_{\rightarrow 0 \text{ as } m \rightarrow +\infty} \end{aligned}$$

So $\mathbb{E}_E [P_m] \rightarrow 0$ as $m \rightarrow +\infty$, and we conclude with $|R|$

2.3 1 $P_x = (p, 1-p)$ for some $p \in (0, 1)$

Then $P_Y(0) = (1-p)(1-\beta) + p\beta$

$$P_Y(1) = p(1-\beta) + (1-p)\beta$$

So $P_{Y|X=0} = (1-\beta, \beta)$, $P_{Y|X=1} = (\beta, 1-\beta)$: same entropy in both cases, so in particular true for their average $H(Y|X)$.

Thus $I(X; Y) = H(Y) - H(Y|X)$

$$= h_2((1-p)(1-\beta) + p\beta) - h(\beta)$$

where $h_2(x) = -x \log_2 x - (1-x) \log_2 (1-x)$

BUT $h_2((1-p)(1-\beta) + p\beta) \leq 1$ with equality if $p = \frac{1}{2}$

Thus $C(W) = 1 - h(\beta)$ QED.