

Homework

Q1) Algorithm: input $r \leftarrow P$ or Q

A

If $P(r) = Q(r)$:

| output P w.p. $1/2$

| output Q w.p. $1/2$

Else if $P(r) > Q(r)$:

| output P

Else, $P(r) < Q(r)$:

| output Q

$$P(\text{win}) = P(A(r) = P / r \leftarrow P) \cdot P(r \leftarrow P) \\ + P(A(r) = Q / r \leftarrow Q) \cdot P(r \leftarrow Q)$$

$$= \frac{1}{2} \left[P(A(r) = P / r \leftarrow P) + P(A(r) = Q / r \leftarrow Q) \right]$$

Given $r \in R$, let us define $a_r = P(A(r) = P)$
We have:

$$P(\text{win}) = \frac{1}{2} \left[\sum_{r \in R} P(r) \cdot a_r + \sum_{r \in R} Q(r) \cdot (1 - a_r) \right]$$

$$= \frac{1}{2} \left[\sum_{r \in R} a_r \cdot [P(r) - Q(r)] + \sum_{r \in R} Q(r) \right]$$

This holds
for every A

$$\left(= \frac{1}{2} \left[1 + \sum_{r \in R} a_r \cdot [P(r) - Q(r)] \right] \right)$$

Now, we have in our case $a_r = \begin{cases} 1/2 & \text{if } P(r) = Q(r) \\ 1 & \text{if } P(r) > Q(r) \\ 0 & \text{else} \end{cases}$

$$\begin{aligned} \text{So } \sum_{r \in R} a_r \cdot [P(r) - Q(r)] &= \sum_{\substack{r \in R \\ P(r) > Q(r)}} P(r) - Q(r) \\ &= \Delta(P, Q) \text{ by the proposition} \end{aligned}$$

$$\text{So } P(\text{win}) = \frac{1}{2} [1 + \Delta(P, Q)]$$

Optimality: Take an algorithm A , and as before let $a_r = P(A(r) = P)$, we also have:

$$P(\text{win}) = \frac{1}{2} \left[1 + \sum_{r \in R} a_r \cdot [P(r) - Q(r)] \right]$$

Now let us bound

delete negative terms

$$\sum_{r \in R} a_r \cdot [P(r) - Q(r)] \leq \sum_{\substack{r \in R \\ P(r) > Q(r)}} a_r [P(r) - Q(r)]$$

$$\leq \sum_{\substack{r \in R \\ P(r) > Q(r)}} P(r) - Q(r) = \Delta(P, Q)$$

$$\text{So } P(\text{win}) \leq \frac{1}{2} [1 + \Delta(P, Q)] \quad \square$$

Q2) • $e_1 = |x+1|$ is already diagonalized of eigen values 1 and 0 so

$$e_1 \log(e_1) = 1 \cdot \log(1) |x+1| = 0$$

hence $H(e_1) = 0$

$$\bullet e_2 = \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \frac{1}{2} |0 \times 0| + \frac{1}{2} |1 \times 1|$$

$$\text{So } \rho_2 \log(\rho_2) = -\frac{1}{2}(1 \times 0 + 1 \times 1)$$

$$\text{and } H(\rho_2) = 1$$

Exercise 2

Let M hermitian

$$M = \rho - \rho'$$

1. Write $M = \sum_{i=1}^n \lambda_i |\psi_i\rangle\langle\psi_i|$

With $|\psi_i\rangle$ O.N. basis, we have that

$$UMU^\dagger = \sum_{i=1}^n \lambda_i U|\psi_i\rangle\langle\psi_i|U^\dagger$$

with $U|\psi_i\rangle$ O.N.B. So M and UMU^\dagger share the same eigenvalues, hence

$$T_2(|UMU^\dagger|) = T_2(U|M|U^\dagger) = T_2(|M|) \quad \square$$

2. Let ρ_1, ρ_2, ρ_3 density matrix

$$\begin{aligned} \Delta(\rho_1, \rho_2) &= \frac{1}{2} T_2(\pi(\rho_1 - \rho_2)) \text{ for a certain } \pi \\ &= \frac{1}{2} T_2(\pi(\rho_1 - \rho_3)) + \frac{1}{2} T_2(\pi(\rho_3 - \rho_2)) \\ &\leq \Delta(\rho_1, \rho_3) + \Delta(\rho_3, \rho_2) \end{aligned}$$

3. a) $\pi = \sum_{i=1}^n |a_i\rangle\langle a_i|$ $M = \sum_{j=1}^n \lambda_j |b_j\rangle\langle b_j|$

$$\begin{aligned} \pi \cdot M &= \sum_{i=1}^n \sum_{j=1}^n \lambda_j |b_j\rangle\langle b_j| a_i\rangle\langle a_i| \\ &= \sum_{i,j} \lambda_j \cdot \langle b_j | a_i \rangle |b_j\rangle\langle a_i| \end{aligned}$$

$$\begin{aligned}
\text{Tr}(\Pi \cdot M) &= \sum_{i,j} \lambda_j \langle b_j | a_i \rangle \langle a_i | b_j \rangle \\
&= \sum_{i,j} \lambda_j |\langle a_i | b_j \rangle|^2 \\
&= \sum_j \lambda_j \sum_i |\langle a_i | b_j \rangle|^2 \leq \|b_j\|^2 = 1 \\
&\leq \sum_j \lambda_j = \text{Tr}(M)
\end{aligned}$$

b) Write $M = \sum_{i=1}^n \lambda_i |\psi_i\rangle\langle\psi_i|$

$$\begin{aligned}
&= \sum_{i=1}^a \lambda_i |\psi_i\rangle\langle\psi_i| - \sum_{i=a+1}^n (-\lambda_i) |\psi_i\rangle\langle\psi_i| \\
&= A - B
\end{aligned}$$

$$\begin{aligned}
\lambda_1, \dots, \lambda_a &\geq 0 \\
\lambda_{a+1}, \dots, \lambda_n &< 0
\end{aligned}$$

And $|M| = \sum_{i=1}^n |\lambda_i| |\psi_i\rangle\langle\psi_i|$

$$= \sum_{i \leq a} \lambda_i |\psi_i\rangle\langle\psi_i| + \sum_{i > a+1} (-\lambda_i) |\psi_i\rangle\langle\psi_i|$$

$$= A + B$$

c) Let $M = \rho - \rho' = A - B$ Note that
 $\text{Tr}(\rho) = \text{Tr}(\rho') = 1$ so $\text{Tr}(M) = 0 \Rightarrow \text{Tr}(A) = \text{Tr}(B)$

$$\Delta(\rho, \rho') = \frac{1}{2} \text{Tr}(|\rho - \rho'|) = \frac{1}{2} \text{Tr}(A + B) = \frac{1}{2} (\text{Tr}(A) + \text{Tr}(B)) = \text{Tr}(A)$$

Φ is a quantum channel, so it is trace preserving

$$\frac{1}{2} \text{Tr}(|\rho - \rho'|) = \text{Tr}(\Phi(A))$$

Now let $\Pi / \Delta(\phi(\rho), \phi(\rho')) = \text{Tr}[\Pi(\phi(\rho) - \phi(\rho'))]$

We have

$$\begin{aligned} \operatorname{Tr}(\phi(A)) &\geq \operatorname{Tr}(\pi \cdot \phi(A)) \\ &\geq \operatorname{Tr}(\pi \cdot [\phi(A) - \phi(B)]) \\ &= \operatorname{Tr}(\pi(\phi(A-B))) \\ &= \operatorname{Tr}(\pi \phi(e - e')) = \Delta(\phi(e), \phi(e')) \end{aligned}$$

□

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1 Homework 8

1. Let P, Q be two probability distributions over a set R . An element r is sampled from P with probability $1/2$ and from Q with probability $1/2$.
 - (a) Propose an algorithm which, on input r , distinguish between the case $r \leftarrow P$ and $r \leftarrow Q$ with probability $1/2 + 1/2 \cdot \Delta(P, Q)$.
 - (b) Show that this success probability is optimal.
2. Compute $H(\rho)$ for:
 - (a) $\rho = |+\rangle\langle+|$.
 - (b) $\rho = I/2$.

2 Trace distance through a quantum channel

We recall that the trace distance is defined as follows:

Definition 2.1. $\Delta(\rho, \rho') = 1/2 \cdot \text{Tr} (|\rho - \rho'|) = \min_{\pi} \text{Tr} (\pi(\rho - \rho'))$ where the minimum is taken among the orthogonal projectors.

1. Prove that $\Delta(\rho, \rho') = \Delta(U\rho U^\dagger, U\rho' U^\dagger)$ for any ρ, ρ' density operator and U unitary.
2. Show that the trace distance follows the triangular inequality.
3. Let ρ, ρ' be two density operators. Let Φ a quantum channel. We are going to show that quantum channels can only make the trace distance decrease.
 - (a) Let M a positive hermitian operator and π an orthogonal projector. Show that $\text{Tr} (\pi M) \leq \text{Tr} (M)$.
 - (b) Let M a hermitian operator. Show that $M = A - B$ with A, B positive hermitian. Show that in this case, $|M| = A + B$.
 - (c) Prove that $\Delta(\rho, \rho') \geq \Delta(\Phi(\rho), \Phi(\rho'))$.

3 Information Theory Quantities

3.1 Conditional entropy

Definition 3.1. The conditional entropy $H(X|Y)$ is defined by

$$H(X|Y) = \sum_{y \in \mathcal{Y}} P_Y(y) H(X|Y = y),$$

where $H(X|Y = y)$ is the entropy of the conditional distribution $P_{X|Y=y}$. Note that elements $y \in \mathcal{Y}$ with $P_Y(y) = 0$ do not participate to the sum.

1. What is the value of $H(X|X)$?

A: $H(X|X = x) = 0$ since $P_{X|X=x}$ is a dirac. Thus, $H(X|X) = 0$.

2. If X and Y are independent random variables, what is the value of $H(X|Y)$?

A: $H(X|Y) = \sum_{y \in \mathcal{Y}} P_Y(y) H(P_{X|Y=y}) = \sum_{y \in \mathcal{Y}} P_Y(y) H(P_X) = H(X)$.

3. Prove that $0 \leq H(X|Y) \leq \log_2(|\mathcal{X}|)$.

A: For all y , $0 \leq H(X|Y = y) \leq \log_2(|\mathcal{X}|)$, and since $H(X|Y)$ is a convex combination of the $H(X|Y = y)$, we get the claimed result.

4. Prove that $H(X|Y) = H(XY) - H(Y)$, where $H(XY) := H((X, Y)) = -\sum_{x,y} P_{XY}(x, y) \log_2 P_{XY}(x, y)$.

A: This is just writing $P_{X|Y=y}(x) = \frac{P_{XY}(x,y)}{P_Y(y)}$ in terms of entropies. We have

$$\begin{aligned} H(XY) &= - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2(P_{XY}(x, y)) \\ &= - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2(P_Y(y) P_{X|Y=y}(x)) \\ &= - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2 P_Y(y) - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2 P_{X|Y=y}(x) \\ &= - \sum_{y \in \mathcal{Y}} P_Y(y) \log_2 P_Y(y) - \sum_{y \in \mathcal{Y}} P_Y(y) \sum_{x \in \mathcal{X}} P_{X|Y=y}(x) \log_2 P_{X|Y=y}(x) \\ &= H(Y) + H(X|Y). \end{aligned}$$

3.2 Mutual Information

Definition 3.2. The mutual information is defined by

$$\begin{aligned} I(X : Y) &= H(X) - H(X|Y) \\ &= H(X) + H(Y) - H(XY). \end{aligned}$$

Writing out the definitions, we get $I(X : Y) = \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2 \frac{P_{XY}(x,y)}{P_X(x)P_Y(y)}$.

1. What is the value of $I(X : X)$?

A: $I(X : X) = H(X) - H(X|X) = H(X)$.

2. If X and Y are independent, what is the value of $I(X : Y)$?

A: $I(X : Y) = H(X) - H(X|Y) = H(X) - H(X) = 0$.

3. For any pair of random variables, prove that $I(X : Y) \geq 0$.

Hint: Use Jensen's inequality on $-I(X : Y)$.

A:

$$\begin{aligned} -I(X : Y) &= - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2 \frac{P_{XY}(x, y)}{P_X(x)P_Y(y)} \\ &= \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \log_2 \frac{P_X(x)P_Y(y)}{P_{XY}(x, y)} \\ &\leq \log_2 \left(\sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_{XY}(x, y) \frac{P_X(x)P_Y(y)}{P_{XY}(x, y)} \right) \\ &= \log_2 \left(\sum_{x \in \mathcal{X}, y \in \mathcal{Y}} P_X(x)P_Y(y) \right) = \log_2(1) = 0 \end{aligned}$$

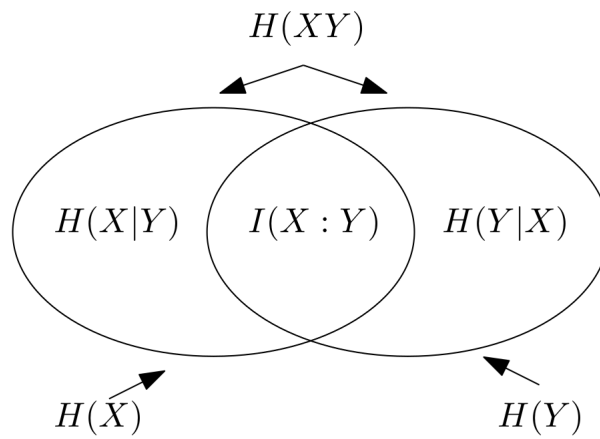


Figure 1: Relation between the entropic measures we have introduced